# PUBLICATIONS

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# MATLOG Program in FORTRAN

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# Preface

This series of reports, entitled MATLOG Program in FORTRAN, summarizes the formulation of BVMF in a practical manner in the form of FORTRAN statements of our BVMF techniques. They closely follow the theoretical approaches developed in Lecture Series 2 and 3. The first three parts will deal with the subprograms that are required for working out atomic statements in SNS logic, quantifier logic, and general logic dealing with the theory of relations.

Part I contains two subprograms BVMF and SNSLOG dealing respectively with the basic notation and algorithms of BVMF in general, and their application to BA-2 algebra as adopted for SNS logic. Part II will contain the extension of this to BA-3 algebra required for quantifier logic, and Part III to general m x n matrices required for multivalued logic adopted for the general theory of relations.

It is believed that the subroutines included in these reports are fairly exhaustive and are sufficient to work out problems in propositional calculus and in predicate calculus. A number of examples are worked out, using the program. In fact, the conclusions derived from these examples are interesting and/ the computer-calculated outputs even have/led to further refinements of the basic Boolean algebraic theory, via the vector-matrix formalism, of logic.

# CONTENTS

# Summary of FORTRAN formulae for the program MATLOG

The program will consist initially of four parts,
namely BVMF, SNSLOG, QLOGIC and GLOGIC.

## 1. BVMF

### (a) BA-1 algebra

In MATLOG we use only Boolean variables and constants
having the two values 1, 0 denoted by the symbols B1, B2.
Variables in BA-1 algebra will be denoted by BA, BB, etc.,
Four basic functions are employed, namely BSUM, BPDT, BEQU,
BCMP corresponding to the binary relations $a \oplus b = c$,
$a \otimes b = c$, $a \Leftrightarrow b = c$, / $a^c = b$ respectively. The essence
and the unary relation
of the FORTRAN formulae for these is as follows:

Boolean sum

$$BSUM \ (BA, \ BB) \ = \ BA \ + \ BB$$

$$IF \ BA \cdot AND \cdot BB \ = \ 1 \ , \ BC \ = \ 1$$

Boolean product

$$BPDT \ (BA, \ BB) \ = \ BA * BB$$

Boolean equivalence $\quad$ BEQU(BA,BB) $=$ BC

$$IF \ BA \cdot EQ \cdot BB \ , \ BC \ = \ 1$$

$$ELSE \ BC \ = \ 0$$

Boolean complement $\quad$ BCMP(BA) $=$ BB

$$BB \ = \ 1 - BA$$

## (b) Vectors

Vectors are indicated in general by GVA,M , GVB,M , etc

where G stands for "general", V stands for "vector" and

M stands for the number of components in the vector . Thus,

GVA(I) , I = 1, MI are the components of GVA,MI .

For the particular cases of MI = 2 and 3 (namely SNS

algebra and quantifier algebra), the vectors are denoted by

SVA and QVA for the two cases of MI = 2, 3 and MI is not

explicitly mentioned, but is indicated by the letters S and Q

at the beginning of the name of the variable or function as the case may be.

## (c) Matrices

These are indicated by GMZ,M,N etc., where G denotes

a general M x N matrix. the letter M stands for "matrix"

and Z is the name of the matrix. The parameters M, N are

the dimensions of the matrix so that

$$GMZ(I,J) ., \quad I = 1, M, \quad J = 1, N$$

represent the components of a general matrix GMZ,M,N

In this case also, if G is replaced by S or Q, then

M = N = 2 or 3 as the case may be, and the matrix refers to

SNS algebra and quantifier algebra respectively.

## (d) Constants in SNS and QL algebra

There are four constants (truth values) in SNS indicated by S1, S2, S3, S4 standing for the vectors (1 0), (0 1), (1 1), (0 0) respectively, for the truth values T, F, D, X respectively.

Similarly, in quantifier algebra, Q1 to Q8 will represent the eight possible quantifier states (constants) in BA-3 algebra. The convention we shall follow is what has been adopted in our BVMF as given below.

| Logical symbol | FORTRAN symbol | Name | BA-3 vector | Logical symbol | FORTRAN symbol | Name | BA-3 vector |
|---|---|---|---|---|---|---|---|
| ∀ | Q1 | 'ALL' | (1 0 0) | ⌽ | Q5 | 'NEX' | (0 0 1) |
| ∧ | Q2 | 'NAL' | (0 1 1) | ∃ | Q6 | 'EXS' | (1 1 0) |
| ⋨ | Q3 | 'SOM' | (0 1 0) | △ | Q7 | 'IND' | (1 1 1) |
| ⊖ | Q4 | 'AON' | (1 0 1) | ⌀ | Q8 | 'IMP'* | (0 0 0) |

*This will be changed to 'XXX' in Part IV .

## (e) Logical connectives in SNS and QL

The ten possible logical connectives in SNS are denoted by SMA,SK,SL; SMO,SK,SL; SME,SK,SL; standing respectively for the matrices $\underline{A}(k, \ell)$, $\underline{O}(k, \ell)$, $\underline{E}(k, \ell)$, $k, \ell = 1, 4$. A similar notation is adopted in QLOGIC, namely

QMA,QK,QL; QMO,QK,QL; QME,QK,QL; QM2,QK,QL standing for $\underline{Z}(k, \ell)$,
$$k, \ell = 1, 8$$

The subprograms in SNS for developing these matrices are given in the appropriate sections below. There are four different A's, and four different O's, while only two different E's are present, making in all the ten possible logical connectives.

## 2. SNSLOG

### (a) Algebraic operators

In this section, we shall write the various functions which are of value in SNS logic. They are given names that are suitable for both logical as well as algebraic purposes, and a general notation **extendable to QLOGIC and GLOGIC is followed. However, the idea is that for** these values as such $M = N = 2$, or $3$, the subroutines are written **for** / without generalizing them / to all values of M and N. These simple subroutines define the functions in these two subprograms and they will be generalized later for GLOGIC. Those required in SNSLOG are as follows.

(i) Boolean sum of two vectors $\underline{a} \oplus \underline{b} = \underline{c}$ ($\oplus = \underline{U} =$ union)

SUNION (SVA,SVB) = SVC

BSUM (SVA(1),SVB(1)) = SVC(1)

BSUM (SVA(2),SVB(2)) = SVC(2)

(ii) Boolean product of two vectors $\underline{a} \otimes \underline{b} = \underline{c}$ ($\otimes = \underline{V} =$ vidya)

SVIDYA(SVA,SVB) = SVC

BPDT(SVA(1),SVB(1)) = SVC(1)

BPDT(SVA(2),SVB(2)) = SVC(2)

(iii) **Boolean complement of SNS vector** $\underline{a}^c = \underline{b}$

$$SCOMP(SVA) = SVB$$
$$BCMP(SVA(1)) = SVB(1)$$
$$BCMP(SVA(2)) = SVB(2)$$

(iv) **Boolean sum of two matrices** $\underline{P} \oplus \underline{Q} = \underline{R}$

$$SMXSUM(SMP,SMQ) = SMR$$
$$BSUM(SMP(I,J), SMQ(I,J)) = SMR(I,J), \quad I, J = 1, 2$$

(v) **Boolean product of two matrices** $\underline{P} \otimes \underline{Q} = \underline{R}$

Replace SMXSUM in (iv) by SMXPDT, and BSUM by BPDT

(vi) **Boolean complement of a matrix operator** $\underline{P}^c = \underline{Q}$

$$SMXCMP(SMP) = SMQ$$
$$BCMP(SMP(I)) = SMQ(I), \quad I = 1, 2$$

(vii) **Scalar product of two vectors** $\langle a \mid b \rangle = c = a_1 \otimes b_1 \oplus a_2 \emptyset$

$$SSCPDT(SVA,SVB) = BC \qquad \text{as above}$$
$$BC = BPDT(SVA(1),SVB(1))$$
$$BC = BSUM(BC,BPDT(SVA(2),SVB(2)))$$

(viii) **Unary product of a vector with a matrix**

$$\langle a \mid z \mid = \langle b \mid , \quad \bigoplus_i a_i \otimes z_{ij} = b_j$$

$$SUNPDT(SVA,SMZ) = SVB$$
$$SVB(J) = 0$$
$$Do \quad I,J = 1, 2$$
$$SVB(J) = BSUM(SVB(J),BPDT(SVA(I),SMZ(I,J)))$$

(This is used for finding the output of a unary relation in SNS. In BVMF, in general, it is the matrix product of a 2-vector with a 2x2 matrix.)

### (ix) Binary product of a matrix with two vectors $\langle a|Z|b\rangle = c$

$$SBINPT(SVA,SMZ,SVB) = BC$$

$$SVAP = SUNPDT(SVA,SMZ)$$

$$BC = SSCPDT(SVAP,SVB)$$

### (x) Matrix product of two matrices : $\bigoplus_j P_{ij} \otimes Q_{jk} = R_{ik}$

$$SMATPT(SMP,SMQ) = SMR$$

$$SMR(I,K) = 0$$

$$DO \quad I,J,K = 1, 2$$

$$SMR(I,K) = BSUM(SMR(I,K), BPDT(SMP(I,J), SMQ(J,K)))$$

(This is used for the successive application of implications in particular and of matrix operations in general).

### (xi) Direct product of two vectors: $a \times b = Z$

$$SDIRPT(SVA,SVB) = SMZ$$

$$SMZ(I,J) = BPDT(SVA(I),SVB(J)), \quad I,J = 1,2$$

### (xii) Direct sum of two vectors: $a + b = Z$

$$SDIRSM(SVA,SVB) = SMZ$$

$$SMZ(I,J) = BSUM(SVA(I),SVB(J)), \quad I,J = 1,2$$

### (xiii) Direct equivalence of two vectors: $(a \equiv b) = Z$

$$SDIREQ(SVA,SVB) = SMZ$$

$$SMZ(I,J) = BEQU(SVA(I),SVB(J)), \quad I,J = 1,2$$

### (xiv) Scalar vector direct product $a \times b = c$

$$SSVDPT(BA,SVB) = SVC$$

$$SVC(I) = BPDT(BA,SVB(I)), \quad I = 1, 2$$

## (d) Constants in SNS and QL algebra

There are four constants (truth values) in SNS indicated

by S1, S2, S3, S4 standing for the vectors (1 0), (0 1), (1 1),

(0 0) respectively, for the truth values T, F, D, X respectively.

Similarly, in quantifier algebra, Q1 to Q8 will represent

the eight possible quantifier states (constants) in BA-3 algebra.

The convention we shall follow is what has been adopted in our

BVMF as given below.

| Logical symbol | FORTRAN symbol | Name | BA-3 vector | Logical symbol | FORTRAN symbol | Name | BA-3 vector |
|---|---|---|---|---|---|---|---|
| ∀ | Q1 | 'ALL' | (1 0 0) | ∯ | Q5 | 'NEX' | (0 0 1) |
| ∧ | Q2 | 'NAL' | (0 1 1) | ∃ | Q6 | 'EXS' | (1 1 0) |
| ⨸ | Q3 | 'SOM' | (0 1 0) | △ | Q7 | 'IND' | (1 1 1) |
| ⊖ | Q4 | 'AON' | (1 0 1) | ⌀ | Q8 | 'IMP'* | (0 0 0) |

*This will be changed to 'XXX' in Part IV .

## (e) Logical connectives in SNS and QL

The ten possible logical connectives in SNS are denoted by

SMA,SK,SL; SMO,SK,SL; SME,SK,SL; standing respectively for the

matrices $\underline{A}(k, \ell)$, $\underline{O}(k, \ell)$, $\underline{E}(k, \ell)$, $k, \ell = 1, 4$. A similar

notation is adopted in QLOGIC, namely

QMA,QK,QL; QMO,QK,QL; QME,QK,QL; QM2,QK,QL standing for $\underline{z}(k, \ell)$,
$$k, \ell = 1, 8$$

The subprograms in SNS for developing these matrices are given

in the appropriate sections below. There are four different A's,

and four different O's, while only two different E's are present,

making in all the ten possible logical connectives.

"general" logic. However, in GLOGIC, we have to specify

also the values of M and N for matrices and vectors. The

list of these Boolean algebraic formulae will be given in

Parts 2 and 3, and they follow very similar patterns in

QLOGIC and GLOGIC as those stated above for SNSLOG. However,

in the application of these formulae to obtain various results

of significance in logic per se, as distinct from algebra in

BVMF, the formulae are slightly different for the three cases

of SNS (BA-2), QL(BA-3) and general logic (BA-n). This is

particularly so for the functions and definitions given below

in subsection (b).


(b) Logical operators

(xvii) Logical connectives $\underline{A}(k, \ell)$, $\underline{O}(k, \ell)$, $\underline{E}(k, \ell)$, $\underline{I}(k, \ell)$

In an obvious notation these are representable in FORTRAN

as follows.

$$SMA,SK,SL = SDIRPT(SK, SL)$$
$$SMO,SK,SL = SDIRSM(SK, SL)$$
$$SME,SK,SL = SDIREQ(SK, SL)$$
$$SMI,SK,SL = SDIRSM(SCOMP(SK), SL)$$

These formulae for the standard logical connectives

for conjunction, disjunction, equivalence or negation, and

implication, require only K and L equal to 1 or 2. However,

they can be analytically continued, for K,L = 3,4 also, for

$\underline{A}(k, \ell)$ and $\underline{O}(k, \ell)$, and hence to $\underline{E}(k, \ell)$ and $\underline{I}(k, \ell)$, which

covers all sixteen possible 2x2 Boolean matrices, as discussed

in Lecture-2 Series-3, (see Tables 2 and 3.)

<u>(xviii) Relative truth value of one term for another</u> :
$$\underline{t}(\underline{a} \mid \underline{b}) = (\langle a \mid b \rangle , \langle a \mid b^c \rangle) = (c_1, c_2)$$

$$SRELTV(SVA, SVB) = SVC$$

$$SVC(1) = SSCPDT(SVA, SVB)$$

$$SVC(2) = SSCPDT(SVA, SCOMP(SVB))$$

The four truth values for $\underline{c}$, namely T, F, D, X, correspond

to the following inclusion properties of the vectors $\underline{a}$ and $\underline{b}$.

T $\longleftrightarrow$ $\underline{a}$ is fully included in $\underline{b}$, and has no
intersection with $\underline{b}^c$.

F $\longleftrightarrow$ $\underline{a}$ is fully included in $\underline{b}^c$ and is not
included in $\underline{b}$.

D $\longleftrightarrow$ $\underline{a}$ is partially included in $\underline{b}$ and is
also partially included in $\underline{b}^c$.

X $\longleftrightarrow$ $\underline{a}$ is the null vector X = (0 0), which
is not included in either $\underline{b}$ or $\underline{b}^c$

binary
**(xix )** SNS truth value of the/relation $Z$ for inputs $\underline{a}$, $\underline{b}$

(formula via contracted product) for a general 2x2 matrix $\underline{Z}$

$$\underline{t}(\underline{a} \ \underline{Z} \ \underline{b}) = (\underline{a}|\underline{Z}|\underline{b}) = \underline{c} = (c_1, c_2) \ ,$$

where $c_1 = \langle \underline{a}|\underline{Z}|\underline{b}\rangle$, $c_2 = \langle \underline{a}|\underline{Z}^c|\underline{b}\rangle$

$$SSTV(SVA, \ SMZ, \ SVB) = SVC$$

$$SVC(1) = SBINPT(SVA, \ SMZ, \ SVB)$$

$$SVC(2) = SBINPT(SVA, \ SCOMP(SMZ), \ SVB)$$

This function SSTV is the general formula for the truth

value of a SNS logical relation when we are given the four

components of the 2x2 matrix $\underline{Z}$ corresponding to the connective,

without specifying its logical nature as being 'and', 'or',

'equ', 'imp' etc. This has been used in the previous programs,

and is generally valid also for QLOGIC and GLOGIC for all

values of M, N.

However, a variation of this function is given below

as SSTV2, which can be used when the logical nature of the

connective is specified as $\underline{Z}(k, \ell)$, with $\underline{Z} = \underline{A}$, $\underline{O}$, $\underline{E}$, or $\underline{I}$,

as the case may be. The corresponding 2x2 matrices are all

direct sums, direct products, or direct equivalences, of two

2-vectors, and therefore the formulae are simpler to manipulate,

**particularly in GLOGIC (see later).**

(xx) SNS truth value of a relation involving the connective

$\underline{z}(k, \ell)$ for inputs $\underline{a}$, $\underline{b}$ (See Lecture 4, Series 2, p.1 of MR-

SSTV2(SVA, SMZ, SVB) = SVC

IF   SMZ = SMA,SK,SL

SVC(1) = BPDT(SSCPDT(SVA, SK),SSCPDT(SVB, SL))

SVC(2) = BSUM(SSCPDT(SVA, SCOMP(SK)),SSCPDT(SVB, SCOMP(SL))

IF   SMZ = SMO,SK,SL

SVC(1) = BSUM(SSCPDT(SVA, SK),SSCPDT(SVB, SL))

SVC(2) = BPDT(SSCPDT(SVA, SCOMP(SK)),SSCPDT(SVB, SCOMP(SL))

IF   SMZ = SME,SK,SL

SKP = SCOMP(SK),  SLP = SCOMP(SL)

SVCA = SSTV2(SVA, SMA,SK,SL, SVB)

SVCB = SSTV2(SVA, SMA,SKP,SLP, SVB)

SVC(1) = BSUM(SVCA(1),SVCB(1))

SVC(2) = BPDT(SVCA(2),SVCB(2))

IF   SMZ = SMI,SK,SL

SKP = SCOMP(SK)

SMZ = SMO,SKP,SL

This subroutine is very much more in the spirit of orthodox logic, and should preferably be used for straightforward logical connectives of the type mentioned in the subroutine. However, both in SNS as well as in QL, general 2x2 matrices and 3x3 matrices can occur for relations and the previous subrout: SSTV in (xix) will have to be invoked in such cases.

**(xxi) Output of unary relation for** $\underline{Z}(k, \ell)$ **via relative**
**truth value:** $\langle \underline{a} | s(k) \rangle \ Z \ \langle s(\ell) \rangle | = \underline{b}$ , $Z = \otimes$ or $\oplus$

$$\text{SUNPT2(SVA,SMZ)} = \text{SVB}$$

IF  SMZ = SMA,SK,SL

$$\text{BAP} = \text{SSCPDT(SVA,SK)}$$

$$\text{SVB} = \text{SSVDPT(BAP,SL)}$$


IF  SMZ = SMO,SK,SL

$$\text{BAP} = \text{SSCPDT(SVA,SK)}$$

$$\text{SVB} = \text{SSVDSM(BAP,SL)}$$


IF  SMZ = SMI,SK,SL

$$\text{BAP} = \text{SSCPDT(SVA,SCOMP(SK))}$$

$$\text{SVB} = \text{SSVDSM(BAP,SL)}$$


IF  SMZ = SME,SK,SL

$$\text{SKP} = \text{SCOMP(SK)}, \quad \text{SLP} = \text{SCOMP(SL)}$$

$$\text{SVBA} = \text{SUNPT2(SVA, SMA,SK,SL)}$$

$$\text{SVBB} = \text{SUNPT2(SVA, SMA,SKP,SLP)}$$

$$\text{SVB} = \text{SUNION(SVBA,SVBB)}$$


**(xxii)** **Matrix-Boolean**
**/binary operator "agree"** $\underline{G}$ , **for checking agreement**
**of two SNS terms :** $(\underline{a} | \underline{G} | \underline{b}) = \underline{c}$

$$\text{SMBG(SVA,SVB)} = \text{SVC}$$

$$\text{BCA} = \text{BEQU(SVA(1),SVB(1))}$$

$$\text{BCB} = \text{BEQU(SVA(2),SVB(2))}$$

$$\text{SVG(1)} = \text{BPDT(BCA,BCB)}$$

$$\text{SVG(2)} = \text{BCMP(SVG(1))}$$

(xxiii) Matrix-Boolean unary operators  E, N, M, L

These are simple operators, defined analogously to quantifier logic, which have been found to be useful in SNS algebra also. In an obvious notation these are

$$SEQU(SVA) = SVB$$
$$SVB(1) = SVA(1), \quad SVB(2) = SVA(2)$$

$$SNOT(SVA) = SVB$$
$$SVB(1) = SVA(2), \quad SVB(2) = SVB(1)$$

$$SCOMP(SVA) = SVB$$
$$SVB(1) = BCMP(SVA(1)), \quad SVB(2) = BCMP(SVA(2))$$

$$SELL(SVA) = SVB$$
$$SVB(1) = BCMP(SVA(2)), \quad SVB(2) = BCMP(SVA(1))$$

Three of these operators are equivalent to those already defined earlier, namely

SEQU $\Longleftrightarrow$ SME,S1,S1,   SNOT $\Longleftrightarrow$ SME,S1,S2,   SCOMP $\Longleftrightarrow$ (ii)

However, the adoption of the names given here, employing simpler definitions given above, makes the implementation of many problems simpler. In fact, even the analogous functions QEQU and QNOT are not definable in terms of QMZ,QK,QL in quantifier logic.

3. Checking of the subroutine in practical problems

The above list of subroutines in SNSLOGIC appear to be sufficient for all applications. A few problems are given below, including routine testing of subroutines, which will both check these, as well as illustrate their applications. In particular, the subroutines (i) to (xvi) for BVMF algebra are not checked as such, but their application in subroutines (xvii) to (xxiii) for SNSLOGIC is tested by working out truth tables and so on for logical derivations.

Problem 1A

Work out the 4x4 truth tables for the SNS truth values
for $a$ and $b$,
T, F, D, $X_{ij}$ for the 16 possible matrix connectives each, of the types $\underline{A}(k, \ell)$, $\underline{O}(k, \ell)$, $\underline{E}(k, \ell)$ and $\underline{I}(k, \ell)$. Do this by using

(a) subroutine (xix) — SSTV

(b) subroutine (xx) — SSTV2

and check the equivalence of each table obtained by employing the two functions SSTV and SSTV2.

Print out the 4x4 truth values for the latter (SSTV2) for all of them and note the differences if any between SSTV and SSTV2.

## Conclusions made from the output of Problem 1A

On checking for discrepancies between the 4x4 truth tables,
it was found that the 3x3 tables for the states  T, F and D
were completely in agreement, while there could be disagreement
if either of the inputs  s($\underline{a}$) or s($\underline{b}$) is  X.  (See Table 1
in the next page 16 for typical examples.)  It will be seen from
the table that, when either $\underline{a}$ or $\underline{b}$ is X , the output truth
value is also X , uniformly for all connectives with SSTV.
On the other hand, it could also be either T or F with SSTV2.

These two correspond to two different interpretations of the input
state  X = (0  0), namely, as corresponding to "absent information"
for SSTV2, and "contradictory information" for SSTV.  (This
will become still clearer when we discuss QLOGIC, where the
differences between BA-3 states being used for describing the
eight quantifier truth values , and for describing multi-
component sets of upto three terms by 3-vectors ($a_1$ $a_2$ $a_3$),
becomes prominent.  In the former case, $\emptyset$ indicates an impossible
truth value, while in the latter case, it only indicates the
non-presence of all the three terms in the set under consideration

## Table 1A Differences between SSTV and SSTV2

Only the tables for $\underline{A}(1, 1)$, $\underline{O}(1, 1)$, $\underline{E}(1, 1)$ and $\underline{I}(1, 1)$ are given below, but the trend for all $\underline{Z}(k, \ell)$ is reasonably clear. (See Part 2)

### SSTV / SSTV2 — A(1, 1)

| a \ b | SSTV T | SSTV F | SSTV D | SSTV X | | SSTV2 T | SSTV2 F | SSTV2 D | SSTV2 X |
|---|---|---|---|---|---|---|---|---|---|
| T | T | F | D | X | | T | F | D | X |
| F | F | F | F | X | | F | F | F | (F) |
| D | D | F | D | X | | D | F | D | (F) |
| X | X | X | X | X | | X | (F) | (F) | X |

$\underline{A}(1, 1) = \underline{a} \wedge \underline{b}$

(F) if $\underline{a} \supseteq s(2), \underline{b}=X$
or $\underline{b} \supseteq s(2), \underline{a}=X$

### SSTV / SSTV2 — O(1, 1)

| a \ b | SSTV T | SSTV F | SSTV D | SSTV X | | SSTV2 T | SSTV2 F | SSTV2 D | SSTV2 X |
|---|---|---|---|---|---|---|---|---|---|
| T | T | T | T | X | | T | T | T | (T) |
| F | T | F | D | X | | T | F | D | X |
| D | T | D | D | X | | T | D | D | (T) |
| X | X | X | X | X | | (T) | X | (T) | X |

$\underline{O}(1, 1) = \underline{a} \vee \underline{b}$

(T) if $\underline{a} \supseteq s(1), \underline{b}=X$
or $\underline{b} \supseteq s(1), \underline{a}=X$

### SSTV / SSTV2 — E(1, 1)

| a \ b | SSTV T | SSTV F | SSTV D | SSTV X | | SSTV2 T | SSTV2 F | SSTV2 D | SSTV2 X |
|---|---|---|---|---|---|---|---|---|---|
| T | T | F | D | X | | T | F | D | X |
| F | F | T | D | X | | F | T | D | X |
| D | D | D | D | X | | D | D | D | (F) |
| X | X | X | X | X | | X | X | (F) | X |

$\underline{E}(1, 1) = \underline{a} \Longleftrightarrow \underline{b}$

This has to be checked further after the two ways of defining $\underline{E}(k, \ell$ in Problem 4B, c are worked out.

### SSTV / SSTV2 — I(1, 1)

| a \ b | SSTV T | SSTV F | SSTV D | SSTV X | | SSTV2 T | SSTV2 F | SSTV2 D | SSTV2 X |
|---|---|---|---|---|---|---|---|---|---|
| T | T | F | D | X | | T | F | D | X |
| F | T | T | T | X | | T | T | T | (T) |
| D | T | D | D | X | | T | D | D | (T) |
| X | X | X | X | X | | (T) | X | (T) | X |

$\underline{I}(1, 1) = \underline{a} \Longrightarrow \underline{b}$
$= \neg\underline{a} \vee \underline{b}$

(T) if $\underline{a} \supseteq s(2), \underline{b}=X$
or $\underline{b} \supseteq s(1), \underline{a}=X$

Problem 1B

Similarly, compare the outputs of SUNPDT and SUNPT2 for the unary relation $a \underline{Z} = b$ for all the connectives $\underline{Z}(k, \ell)$ $k, \ell = 1$ to 4, for $\underline{Z} = \underline{A}, \underline{O}, \underline{I}, \underline{E}$ .

Comment on the output of Problem 1B

Just as for the comparison between SSTV and SSTV2, in this case also, the/two subroutines SUNPDT and SUNPT2 agree, $\overset{output\ b\ of\ the}{}$ for the three inputs $\overset{a=}{}$/T, F, D, for all connectives $\underline{Z}(k, \ell)$, but for X as input, they do not agree sometimes. Some typical examples are given in Table 1B.

It will be seen that the two agree for $\underline{A}(k, \ell)$ and $\underline{E}(k, \ell)$ and this is true quite generally. However, for $\underline{Q}(k, \ell)$ and $\underline{I}(k, \ell)$, for X as input, SUNPDT gives X as output, while SUNPT2 always gives s( $\ell$ ) as output. This is also explicable by the same considerations as mentioned above for Problem 1A.

We shall give in problems 2 and 3 the way in which a set of sequentially implementable BVMF equations can be worked out in FORTRAN for MATLOG. The examples are taken from previous lectures.

## Table 1B: Differences between SUNPDT and SUNPT2

Chosen examples are given below and a summary of the observations is given in the text. The first columns correspond to SUNPDT and the second to SUNPT2

|   | A(1, 1) |   | A(1, 2) |   | A(1, 3) |   | A(1, 4) |   |
|---|---|---|---|---|---|---|---|---|
| T | T | T | F | F | D | D | X | X |
| F | X | X | X | X | X | X | X | X |
| D | T | T | F | F | D | D | X | X |
| X | X | X | X | X | X | X | X | X |

|   | O(1, 1) |   | O(1, 2) |   | O(1, 3) |   | O(1, 4) |   |
|---|---|---|---|---|---|---|---|---|
| T | D | D | D | D | D | D | D | D |
| F | T | T | F | F | D | D | X | X |
| D | D | D | D | D | D | D | D | D |
| X | X | Ⓣ | X | Ⓕ | X | Ⓓ | X | X |

|   | O(2, 1) |   | O(2, 2) |   | O(2, 3) |   | O(2, 4) |   |
|---|---|---|---|---|---|---|---|---|
| T | T | T | F | F | D | D | X | X |
| F | D | D | D | D | D | D | D | D |
| D | D | D | D | D | D | D | D | D |
| X | X | Ⓣ | X | Ⓕ | X | Ⓓ | X | X |

|   | I(1, 1) |   | I(1, 2) |   | I(1, 3) |   | I(1, 4) |   |
|---|---|---|---|---|---|---|---|---|
| T | T | T | F | F | D | D | X | X |
| F | D | D | D | D | D | D | D | D |
| D | D | D | D | D | D | D | D | D |
| X | X | Ⓣ | X | Ⓕ | X | Ⓓ | X | X |

|   | E(1, 1) |   | E(1, 2) |   | E(1, 3) |   | E(1, 4) |   |
|---|---|---|---|---|---|---|---|---|
| T | T | T | F | F | D | D | X | X |
| F | F | F | T | T | X | X | D | D |
| D | D | D | D | D | D | D | D | D |
| X | X | X | X | X | X | X | X | X |

## Problem 2

Take the argument given in Lecture-3, Series-3, page 20.

The formulae in BVMF and in FORTRAN for MATLOG are given below.

The output data are to be printed out in the format given.

<div align="center">

**T, F for**
INPUT/SVA, SVB, SVC

</div>

1. $\underline{a}\ \underline{A}\ \underline{b} = \underline{g}$      SVG = SSTV2(SVA, SMA,S1,S1, SVB)

2. $\underline{c}\ \underline{O}\ \underline{b} = \underline{h}$      SVH = SSTV2(SVC, SMO,S1,S1, SVB)

3.       $T = \underline{x}''$      SVXPP = S1

4. $\underline{g}\ \underline{I} = \underline{k}$      SVK = SUNPLT(SVG, SMI,S1,S1)

5. $\underline{h}\ \underline{I}(1,2) = \underline{y}$      SVY = SUNPDT(SVH, SMI,S1,S2)

6. $\underline{h}\ \underline{A}\ \underline{k} = \underline{x}'$      SVXP = SSTV2(SVH, SMA,S1,S1, SVK)

7. $\underline{x}'\ \underline{\underline{V}}\ \underline{x}'' = \underline{x}$      SVX = SVIDYA(SVXP,SVXPP)

                     PRINT SVG, SVH, SVXPP, SVK, SVY, SVXP, SVX

| Sl. No. | A | B | C | G | H | XPP | K | Y | XP | X |
|---|---|---|---|---|---|---|---|---|---|---|
| 1. | T | T | T | | | | | | | |
| | | — | | | | | | | | |
| | | — | | | | | | | | |
| 8. | F | F | F | | | | | | | |

## Solution of Problem 2

The output as given by the program for this problem is

given below. It can be verified that it agrees completely

with the table given in Lecture-3, Series-3 for this problem.

| A | B | C | G | H | XPP | K | Y | XP | X |
|---|---|---|---|---|-----|---|---|----|---|
| T | T | T | T | T | T | T | F | T | T |
| T | T | F | T | T | T | T | F | T | T |
| T | F | T | F | T | T | D | F | D | T |
| T | F | F | F | F | T | D | D | F | X |
| F | T | T | F | T | T | D | F | D | T |
| F | T | F | F | T | T | D | F | D | T |
| F | F | T | F | T | T | D | F | D | T |
| F | F | F | F | F | T | D | D | F | X |

It so happens that the only output, SVX, is either T,or X,

but this is not generally true. Also, note that the state D

occurs as an intermediate output, e.g. for SVK,SVY and SVXP,

and F occurs for all other intermediate outputs (SVXPP is an

input). See the text of Lecture 3, Series 3 for further discussion.

Problem 3

Consider the problem given in pages 34, 36 of Lecture-3,Series
The FORTRAN formulae are given/and they have already been
below
arranged for sequential implementation,with the six stages as
marked.

INPUTS    SVA1, SVA2, SVA3, SVA4

SVB1  =  SUNPDT(SVA1, SMI,S1,S1)                    (1.1)

SVC1  =  SSTV(SVA1, SMO,S1,S1, SVA2)                (1.2)

SVB2  =  SUNPDT(SVA3, SMI,S1,S2)                    (1.3)

SVD2P =  SSTV(SVA3, SMO,S1,S1, SVA4)                (1.4)


SVC2  =  SSTV(SVA2, SMA,S1,S1, SVB1)                (2.1)

SVD1  =  SSTV(SVA3, SMO,S2,S2, SVC1)                (2.2)


SVD2PP =  SSTV(SVD1, SMA,S1,S1, SVB2)               (3.1)

SVE1P =  SUNPDT(SVC2, SMI,S1,S1)                    (3.2)


SVD2  =  SVIDYA(SVD2P, SVD2PP)                      (4.1)
PRINT * * * IF SVIDYA = X,  STOP

SVE1PP = SUNPDT(SVD2, SMI,S1,S1)                    (4.2)


SVE2  =  SSTV(SVA4, SMO,S1,S1, SVD2)                (5.1)
SVE1  =  SVIDYA(SVE1P, SVE1PP)                      (5.2)
PRINT * * * IF SVIDYA = X,  STOP


SVE   =  SVIDYA(SVE1, SVE2)                         (6.1)
PRINT * * * IF SVIDYA = X,  STOP

The output is to be printed out in the same manner as Problem 2
with spaces between the different stages.

For ready reference, the classical logic equivalents

of the MATLOG equations (1.1) to (6.1) are given in CL-1 to

CL-11 below.

$$a1 \implies b1 \tag{CL-1}$$

$$a1 \lor a2 \iff c1 \tag{CL-2}$$

$$a3 \implies \neg g1 \tag{CL-3}$$

$$a2 \land b1 \iff c2 \tag{CL-4}$$

$$\neg a3 \lor \neg c1 \iff d1 \tag{CL-5}$$

$$\neg g1 \implies \neg b2 \tag{CL-6}$$

$$d1 \land b2 \iff d2 \tag{CL-7}$$

$$a3 \lor a4 \iff d2 \tag{CL-8}$$

$$c2 \implies x \tag{CL-9}$$

$$d2 \implies x \tag{CL-10}$$

$$a4 \lor d2 \iff x \tag{CL-11}$$

It should be noticed that CL-7 and CL-8 both have the same

output $d2$, and should therefore be checked for consistency,

and similarly, CL-9, 10, 11, all of which have the same output

X, should also be tested for consistency.

The logical graph of the CL statements, and their modification

in BVMF by including vidya checks, is contained in Fig.1 on page

21b. For further details Lecture-3, Series-3 may be referred to.

Fig.1. Logical graph corresponding to the set of statements
listed above.

Solution of Problem 3

Just as in the case of Problem 2, for this problem also,

the computer output fully substantiates the results worked out

manually in Lecture-3, Series-3. The output is given in the

next page. There is a small difference in the logical graph

in that $\underline{d2}$ is the input in $(4.2)_1$ instead of $\underline{d2"}_1$ and there

are minor differences as a consequence, but no changes in the

occurrence of contradictions. The predictions made in

Lecture-3, as to which of the 16 possible inputs will lead to

contradictions, and which will not, are fully verified.

for Problem 1,
In view of the discussion given above, some additional

problems in SNSLOG were worked out and are described below.


4.   Additional problems in SNSLOG

A fuller analysis of the differences between SSTV and

SSTV2, and between SUNPDT and SUNPT2, will be given in Part-3

while dealing with GLOGIC, after considering also similar

differences found in QLOGIC and GLOGIC. In essence, the two

functions are not logically equivalent and therefore, not

identical, but they do not lead to any differences in the

predictions for classical logic using BA-1 truth values.

However, their difference becomes manifest in SNS and QLOGIC

employing BA-2 and BA-3 algebra, and particularly so with

m x n matrices in GLOGIC. A general discussion is given in

Part 3, Section 5 , but we shall consider below two or three

examples in the form of problems  to illustrate the merits,

and consistency, of the matrix formulation in SNS logic. These

will be followed up for QLOGIC and later for GLOGIC.


Problem 4A

The following two equations (1) and (2) have been

implicitly used in our studies with BVMF so far. A discussion

of these will be given in Part-3, which will also indicate

their limitations, as well as the conditions under which they

are valid.

$$(a \underset{=}{P} b) \vee (a \underset{=}{Q} b) = a (P \oplus Q) b \tag{1}$$

$$(a \underset{=}{P} b) \wedge (a \underset{=}{Q} b) = a (P \otimes Q) b \tag{2}$$

The l.h.s of Eqs (1) and (2) are the disjunction and conjunction

respectively of the truth values of the relations $a \underset{=}{P} b$

and $a \underset{=}{Q} b$, connecting two terms $a$ and $b$ via  two different

relations P and Q. However, in BVMF, these can be converted

into single relations a R b, and a R' b , respectively, with

R = P $\oplus$ Q   and   R' = P $\otimes$ Q , and this can be

verified, for the inputs T and F for a and b , in all cases.

Obviously, these are the only states that are valid in

classital logic, and it would appear as if the extension to

SNSLOGIC in Eqs.(1) and (2) would be valid generally. However,

a complete test of this, made by employing the subroutines-in
 this is not always true.
MATLOG, indicate that / This was done by using the following

MATLOG equations for checking the validity of Eqs.(1) and (2).

$$SSTV(SVA,SMP,SMB) = SVVP$$

$$SSTV(SVA,SMQ,SMB) = SVVQ$$

$$SSTV(SVVP,SMO,S1,S1,SVVQ) = SVR1$$

$$SMXSUM(SMP,SMQ) = SMR$$

$$SSTV(SVA,SMR,SVB) = SVR2$$

$$SMBG(SVR1,SVR2) = SVG1$$

$$SSTV(SVVP,SMA,S1,S1,SVVQ) = SVR3$$

$$SMXPDT(SMP,SMQ) = SMRP$$

$$SSTV(SVA,SMRP,SVB) = SVR4$$

$$SMBG(SVR3,SVR4) = SVG2$$

The checks to be made can be given as follows.

Compare the 4x4 truth tables of SVR1 and SVR2, and of SVR3

and SVR4, for T, F, D, X /for the truth values for SVA and
                          as inputs

SVB, and print these, as well as the truth values SVG1 and SVG2

of the agreement between these, for the following pairs $P,Q$ :


$A(1,1),A(1,2)$;  $A(1,1),A(2,1)$; $A(1,1),A(2,2)$;  $A(1,1),A(1,3)$;
                                                    $A(3,3),A(4,4)$;

$Q(2,1),Q(1,2)$;  $Q(1,1),Q(2,2)$;  $Q(1,1),Q(1,2)$;  $Q(3,1),Q(3,2)$;
                                                      $Q(1,3),Q(1,4)$;

$E(1,1),E(1,2)$;  $E(1,3),E(3,2)$;  $E(1,4),E(2,1)$;

$I(1,1),I(2,2)$;  $I(1,2),I(2,1)$;  $A(1,1),Q(1,1)$;  $A(2,1),Q(1,2)$

$Q(1,2),A(1,1)$;  $Q(2,2),A(1,1)$;  $Q(3,4),A(1,2)$;  $Q(3,3),A(2,1)$


## Comments on the output of Problem 4A

Some typical examples of the outputs obtained in the

above tests are given below in Table 3.    In all the outputs,

wherever G1 or G2 is F, showing disagreement between the l.h.s

and r.h.s of Eq.(1) or (2), a ring has been put round this

entry.  It will be noticed that the rings are very few and

far between, but it would be highly desirable to obtain

definite rules for identifying their occurrence.  Therefore,

A(1.1)    A(1,2)

| A | B | R1 | R2 | G1 | R3 | R4 | G2 |
|---|---|----|----|----|----|----|----|
| 1 | T | T | T | T | F | F | T |
| 1 | F | T | T | T | F | F | T |
| 1 | D | D | T | (F) | D | F | (F) |
| 1 | X | X | X | T | X | X | T |
| F | T | F | F | T | F | F | T |
| F | F | F | F | T | F | F | T |
| F | D | F | F | T | F | F | T |
| F | X | X | X | T | X | X | T |
| C | T | D | D | T | F | F | T |
| C | F | D | D | T | F | F | T |
| C | D | D | D | T | D | F | (F) |
| C | X | X | X | T | X | X | T |
| X | T | X | X | T | X | X | T |
| X | F | X | X | T | X | X | T |
| X | D | X | X | T | X | X | T |
| X | X | X | X | T | X | X | T |

**DISAGREE    DISAGREE**

$$|P| = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$|Q| = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$$

$$|R| = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix}$$

$$|R'| = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}$$

Table 3(a)

A(1.1) . A(2,2)

| A | B | R1 | R2 | G1 | R3 | R4 | G2 |
|---|---|----|----|----|----|----|----|
| 1 | T | T | T | T | F | F | T |
| 1 | F | F | F | T | F | F | T |
| 1 | D | D | D | T | F | F | T |
| 1 | X | X | X | T | X | X | T |
| F | T | F | F | T | F | F | T |
| F | F | T | T | T | F | F | T |
| F | D | D | D | T | F | F | T |
| F | X | X | X | T | X | X | T |
| C | T | D | D | T | F | F | T |
| C | F | D | D | T | F | F | T |
| C | D | D | D | T | C | F | (F) |
| L | X | X | X | T | X | X | T |
| X | T | X | X | T | X | X | T |
| X | F | X | X | T | X | X | T |
| X | D | X | X | T | X | X | T |
| X | X | X | X | T | X | X | T |
| | | AGREE | | | DISAGREE | | |

$$|P| = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$|Q| = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$

$$|R| = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = |E(1, 1)|$$

$$|R'| = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}$$

$$R = E(1, 1)$$

Table 3(b)

A(1.1) . A(1.3)

| A | B | R1 | R2 | G1 | R3 | R4 | G2 |
|---|---|----|----|----|----|----|----|
| T | T | T | T | T | T | T | T |
| T | F | T | T | T | F | F | T |
| T | D | T | T | T | D | D | T |
| T | X | X | X | T | X | X | T |
| F | T | F | F | T | F | F | T |
| F | F | F | F | T | F | F | T |
| F | D | F | F | T | F | F | T |
| F | X | X | X | T | X | X | T |
| D | T | D | D | T | D | D | T |
| D | F | D | D | T | F | F | T |
| D | D | D | D | T | D | D | T |
| D | X | X | X | T | X | X | T |
| X | T | X | X | T | X | X | T |
| X | F | X | X | T | X | X | T |
| X | D | X | X | T | X | X | T |
| X | X | X | X | T | X | X | T |

$$|P| = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$|Q| = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix}$$

$$|R| = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$|R'| = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix}$$

AGREE          AGREE

$$\underline{P} \subset \underline{Q}$$

Table 3(c)

C(2.1) . O(1,2)

| A | B | R1 | R2 | G1 | R3 | R4 | G2 |
|---|---|----|----|----|----|----|----|
| 1 | T | T | T | T | T | T | T |
| 1 | F | T | T | T | F | F | T |
| 1 | D | T | T | T | D | C | T |
| 1 | X | X | X | T | X | X | T |
| F | T | T | T | T | F | F | T |
| F | F | T | T | T | T | T | T |
| F | D | T | T | T | C | C | T |
| F | X | X | X | T | X | X | T |
| C | T | T | T | T | D | C | T |
| C | F | T | T | T | C | C | T |
| C | D | D | T | (F) | D | C | T |
| C | X | X | X | T | X | X | T |
| X | T | X | X | T | X | X | T |
| X | F | X | X | T | X | X | T |
| X | D | X | X | T | X | X | T |
| X | X | X | X | T | X | X | T |

DISAGREE        AGREE

$$R' = E(1, 1)$$

Table 3(d)

$$|P| = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$$

$$|Q| = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$$

$$|R| = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$

$$|R'| = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = E(1, 1)$$

C(1.1) . O(1.2)

| A | B | R1 | R2 | G1 | R3 | R4 | G2 |
|---|---|----|----|----|----|----|----|
| 1 | T | T | T | T | T | T | T |
| 1 | F | T | T | T | T | T | T |
| 1 | D | T | T | T | T | T | T |
| 1 | X | X | X | T | X | X | T |
| F | T | T | T | T | F | F | T |
| F | F | T | T | T | F | F | T |
| F | D | D | T | (F) | D | F | (F) |
| F | X | X | X | T | X | X | T |
| D | T | T | T | T | D | D | T |
| D | F | T | T | T | D | D | T |
| D | D | D | T | (F) | D | D | T |
| D | X | X | X | T | X | X | T |
| X | T | X | X | T | X | X | T |
| X | F | X | X | T | X | X | T |
| X | D | X | X | T | X | X | T |
| X | X | X | X | T | X | X | T |

$$|P| = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$$

$$|Q| = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$$

$$|R| = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$

$$|R'| = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix}$$

Table 3(e)

| A | B | R1 | R2 | G1 | R3 | R4 | G2 |
|---|---|----|----|----|----|----|----|
| T | T | T | F | T | F | F | T |
| T | F | T | F | F | F | T | T |
| T | D | T | F | T | D | D | T |
| T | X | X | X | T | X | X | T |
| F | T | F | F | T | F | F | T |
| F | F | T | F | T | F | F | T |
| F | D | D | D | T | F | F | T |
| F | X | X | X | T | X | X | T |
| D | T | D | D | T | F | F | T |
| D | F | T | F | T | F | F | T |
| D | D | D | D | T | D | D | T |
| D | X | X | X | T | X | X | T |
| X | T | X | X | T | X | X | T |
| X | F | X | X | T | X | X | T |
| X | D | X | X | T | X | X | T |
| X | X | X | X | T | X | X | T |

$$|P| = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix}$$

$$|Q| = \begin{pmatrix} 0 & 1 \\ 0 & 1 \end{pmatrix}$$

$$|R| = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$$

$$|R'| = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$$

    AGREE          AGREE

$P^c = Q$    P and Q have two 1's in a row or column

Table 3(f)

I(1,2) . I(2,1)

| A | B | R1 | R2 | G1 | R3 | R4 | G2 |
|---|---|----|----|----|----|----|----|
| T | T | T | T | T | F | F | T |
| T | F | T | T | T | T | T | T |
| T | D | T | T | T | D | D | T |
| T | X | X | X | T | X | X | T |
| F | T | T | T | T | T | T | T |
| F | F | T | T | T | F | F | T |
| F | D | T | T | T | D | D | T |
| F | X | X | X | T | X | X | T |
| D | T | T | T | T | D | D | T |
| D | F | T | T | T | D | D | T |
| D | D | D | T | Ⓕ | D | D | T |
| D | X | X | X | T | X | X | T |
| X | T | X | X | T | X | X | T |
| X | F | X | X | T | X | X | T |
| X | D | X | X | T | X | X | T |
| X | X | X | X | T | X | X | T |

$$|P| = \begin{pmatrix} 0 & \cdot 1 \\ 1 & 1 \end{pmatrix}$$

$$|Q| = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$$

$$|R| = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$

$$|R'| = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} =$$

AGREE

$$\underline{R}' = \underline{E}(1, 2)$$

Table 3(g)

C(1,2) . A(1,1)

| A | B | R1 | R2 | G1 | R3 | R4 | G2 |
|---|---|----|----|----|----|----|----|
| 1 | T | T | T | T | T | T | T |
| 1 | F | T | T | T | F | F | T |
| 1 | D | T | T | T | D | C | T |
| 1 | X | X | X | T | X | X | T |
| F | T | F | F | T | F | F | T |
| F | F | T | T | T | F | F | T |
| F | D | D | D | T | F | F | T |
| F | X | X | X | T | X | X | T |
| C | T | D | D | T | C | D | T |
| C | F | T | T | T | F | F | T |
| C | D | D | D | T | D | C | T |
| C | X | X | X | T | X | X | T |
| X | T | X | X | T | X | X | T |
| X | F | X | X | T | X | X | T |
| X | D | X | X | T | X | X | T |
| X | X | X | X | T | X | X | T |

AGREE     AGREE

$$|P| = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$$

$$|Q| = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$|R| = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$$

$$|R'| = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$P \supset Q$

Table 3(h)

we shall first focus attention on those cases where there is complete agreement between the left- and right-hand sides of Eq.(1), or (2). These are marked by the entry "Agree" below the corresponding table. On looking at these, it is found that the agreements occur only if one of the following three conditions occur :

(a) The condition $|\underline{P}| \supseteq |\underline{Q}|$, $|\underline{P}| \subseteq |\underline{Q}|$, is satisfied, when both sum and product agree.

(b) The condition $|\underline{P}| = |\underline{Q}^c|$ is satisfied, when also $|\underline{P}|$ and $|\underline{Q}|$ have two 1's when also both sum and product agree.

(c) The matrix of the resultant operator $\underline{R}$, or $\underline{R}'$, is $|\underline{E}|$ or $|\underline{N}|$, in which case either sum or product agrees, whichever leads to the "equivalence" operator $\underline{E}(k,\ell)$, $k, \ell = 1,2$.

Very similar conditions are found to be true also in QLOGIC and still more generally in GLOGIC, for $\underline{Z}(k, \ell)$, and a theoretical justification of these is given for the general case in Part 3.

Thus, Eqs.(1) and (2) are not valid for SNS truth values as a general rule, but only under the restricted conditions given in (a), (b), (c) above. However, as will be shown in Part 3, for "Boolean truth values" defined by the function BTV(SVA,SMZ,SVB)$\Longleftrightarrow$SBINPT(SVA,SMZ,SVB) = SVC, which corresponds to the Boolean algebraic equation $\langle \underline{a}|\underline{Z}|\underline{b}\rangle = c = $ BTV, the analogue of Eq.(1) using the Boolean sum is always true,while that of Eq.(2) employing the Boolean product is not necessarily true.

However, for the standard logical connectives $\underline{E}(= \underline{E}(1, 1))$

and $\underline{N} = (= \underline{E}(1, 2))$, which can be defined both as conjunction

and disjunction of relations involving $\underline{A}(k, \ell)$ and $\underline{O}(k, \ell)$,

**both** Eqs.(1) and (2) are valid. Thus, three different, but

equivalent, definitions of $\underline{E}(1, 1)$ can be given, as in (3a),(3b)

and (3c) below:

$$\underline{E}(1, 1) = |s(1)\rangle \iff \langle s(1)|$$ (3a)

$$= \underline{A}(1,1) \oplus \underline{A}(2,2)$$ (3b)

$$= O(2,1) \otimes O(1,2)$$ (3c)

All of them lead to the same matrix as shown in 4(a,b,c).

$$|E| = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \equiv (1 \ 0) = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$ (4a)

$$\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \oplus \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$ (4b)

$$\begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \otimes \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$ (4c)

In fact, the general definition of $\underline{E}(k, \ell)$ as SME,SK,SL holds

in a manner very similar to Eqs(4a,b,c), as was verified by

detailed computer checking **via** MATLOG using Eqs.(5a,b,c) below

and the subroutine SSTV employing the BVMF binary product

$(\underline{a}|\underline{Z}|\underline{b}) = \underline{c}$ . An algebraic checking of this, following the

pattern in Eqs.(4a,b,c), is given in Table 4 below in the section

dealing with comments on Problem 4B.

This means that the ten standard connectives, $\underline{A}(k,\ell)$, $\underline{Q}(k,\ell)$, &

$\underline{E}(k,\ell)$ (namely $\underline{E}(1, 1)$ and $\underline{E}(1, 2)$) can be treated freely

using the SSTV formalism without any difficulty.

Probmem 4B
In order to check this for all $\underline{E}(k, \ell)$ with $k, \ell$ = 1 to 4,

the results given by three ways of defining the equivalence

operator as given in Eq.(5) below, can be tested _via_ the MATLOG

statements following them.

$$\underline{E}(k, \ell) \; = \; \left|s(k)\right\rangle \Longleftrightarrow \left\langle s(\ell)\right| \tag{5a}$$

$$\underline{E}(k, \ell) \; = \; \underline{A}(k, \ell) \; \oplus \; \underline{A}(k^c, \ell^c) \tag{5b}$$

$$\underline{E}(k, \ell) \; = \; \underline{Q}(k^c, \ell) \; \otimes \; \underline{Q}(k, \ell^c) \tag{5c}$$

SVC1 = SSTV(SVA,SME,SK,SL,SVB) $\Longleftrightarrow$ (4a)

SVC2 = SSTV(SVA1,SMO,S1,S1,SVA2) $\Longleftrightarrow$ (4b)

where
SVA1 = SSTV(SVA,SMA,SK,SL,SVB)

SVA2 = SSTV(SVA,SMA,SCOMP(SK),SCOMP(SL),SVB)

SVC3 = SSTV(SVI1,SMA,S1,S1,SVI2) $\Longleftrightarrow$ (4c)
where
SVI1 = SSTV(SVA,SMO,SCOMP(SK),SL,SVB)

SVI2 = SSTV(SVA,SMO,SK,SCOMP(SL),SVB)

Make the check for all $\underline{E}(k,\ell)$, $k, \ell$ = 1 to 4 and print

them. Verify that the three values SVC1, SVC2 and SVC3 agree

in each case.

## Comments on Problem 4B

It is found that the three different ways of defining

$E(k, \ell)$ , as in (5a), (5b) and (5c) , lead to identically the

same truth values for SVC1, SVC2 and SVC3 , for all k, $\ell$ , and for

all combinations of SVA and SVB. This can be seen even more

clearly, from the matrix formalism viewpoint , in Table 4, where

the extension of the definition of $E(k, \ell)$, from $E(1, 1)$ and

$E(1, 2)$, to all possible combinations of k, $\ell$ is illustrated.

The three ways of defining E are given respectively in columns

3, 2 and 4. It is very int¢resting to note that, formally, all the

possible matrices of the type $E(k, \ell)$ cover the eight 2x2

matrices that are left, after considering $A(k, \ell)$ and $O(k, \ell)$

for k, $\ell$ = 1, 2. It is also reassuring to note that the formal

Boolean algebraic definition of the logical operators A, O and E

as Boolean direct product, direct sum, and direct equivalence,

of SK and SL is completely valid in the extended BA-2 algebra.

## Problem 4C
In a similar manner, the problem in 4B can be checked

using SSTV2 formalism. However, in this case, Eq.(5a) has no

description via SSTV2, and therefore only the possibilities (5b)

and (5c.) need be tested. This can be done by using analogous

MATLOG statements as those given for Problem 4B, but replacing

symbol SSTV2 by SSTV in all places. Calling these outputs

for (5b) and (5c) as SVC2P and SVC3P respectively, Problem 4C

can be stated as follows.

Make the calculation for all $\underline{\underline{E}}(k, \ell)$, $k, \ell = 1$ to 4

with SSTV2 and print the results for SVC2P and SVC3P in each

case. Check them for agreement.

## Comments on Problem 4C

The full computer output has not been obtained yet[*], but

the essential features are evident from a consideration of

the simple example of $\underline{\underline{E}}(1, 1)$. It is found that the 4 x 4

truth tables as given by SSTV and SSTV2 do not/agree, and
<br>fully

what is more, the two different definitions in Eqs. (5b) and

(5c) which can be programed in SSTV2 do not themselves fully agree

with one another. Table 5 below contains the full 4x4 truth

tables corresponding to Eq.(5a), using SSTV, and for Eqs.(5b),

(5c), using SSTV2. Although for T, F and D there is complete

agreement between all three tables, the two entries for

(D, X) and (X, D) are completely different in the three cases —

namely, equal to X for SVC1, F for SVC2P and T for SVC3P. Thus,

it appears that the extension from classical logic into the

*This has been obtained and will be discussed
in the appendix to Part III, MR-62. (Added 10.9.87).

Table 5. Comparison of 4x4 truth tables for $\underline{E}(1, 1)$ <u>calculated via</u> the subroutine SSTV2 for SVC2P and SVC3P with that for SSTV

| $\underline{a}$ | $\underline{b}$ | SVC1* (in SSTV) | | | | SVC2P (in SSTV2) | | | | SVC3P (in SSTV2) | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| | | T | F | D | X | T | F | D | X | T | F | D | X |
| T | | T | F | D | X | T | F | D | X | T | F | D | X |
| F | | F | T | D | X | F | T | D | X | F | T | D | X |
| D | | D | D | D | X | D | D | D | Ⓕ | D | D | D | Ⓣ |
| X | | X | X | X | X | X | X | Ⓕ | X | X | X | Ⓣ | X |

*SVC2 and SVC3 completely agree with the table for SVC1.

---

BA-2 algebra of SNSLOGIC is consistent only for the matrix formalism and leads to inconsistencies if it is extended <u>via</u> truth values. Although the discrepancies occur only for the extremely artificial situation of the input and output being tautology (D) and contradiction (X) respectively, still, for mathematical consistency, one should expect the algebra to give identical results. It is not yet clear* whether any logical redefinition of the functions SSTV2 could lead to complete consistency between different ways of defining the equivalence operator. This should be discussed after similar problems are considered in QLOGIC and GLOGIC*.

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* They are fully explicable (See MR-62, appendix).

# MATLOG Program in FORTRAN

## Part II : MATLOG Program for quantifier logic and their tests

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## Part II: MATLOG Programs for quantifier logic and their tests

### 1. Quantifier logic in BVMF

#### (a) General

The theory of quantifiers in BVMF has been developed first in MR-25A which has been published in Current Science (Parts I and II , Vol 52, pages 292 and 335, 1983). Since then, it has undergone several revisions, e.g. in MR-33, 45 46A, 48, 49, 50 and 51. However, a definitive treatment of the theory of quantifiers via well-formed formulae in SNSLOG was given in MR-52, and this laid the basis for the distinction between QL-1 and QL-2 as two standard types of relations in the BVMF description of quantified statements. The essential differences between these two are indicated by the following equations.

QL-1 : $(q_z x)(\underset{\sim}{a}x\ \underline{Z}(k,\ell)\ \underline{b}x)$ , e.g. $(\exists x)(\underline{a}x \Longrightarrow \underline{b}x)$
QL-2 : $\underset{\sim}{a}x\ \underline{Z}(k,\ell)\ \underline{b}y$ , e.g. $(\forall x)(\underline{a}x) \Longrightarrow (\exists y)(\underline{b}y)$

The treatment of these two types in BVMF algebra are clearly different. The latter, which is first considered, can be implemented in exactly the same manner as SNS, but using 3-vectors and 3x3 Boolean matrices. For doing this, a set of three basic states / (1 0 0),(0 1 0), (0 0 1), denoted by "For all", "For some", "For none", are used, and these lead to the eight possible quantifier states Q1 to Q8

which are described in MR-52. The implementation in OL-2

is particularly fully treated in MR-53, and the implementation

in OL-1 in MR-54. These are reviewed and summarized in MR-56,

Lectures 3 and 4, Series 2.

In our treatment below, we shall designate CL-2 as

QLOGIC and QL-1 as PLOGIC. In the case of QL-2, the algebraic

subroutines are identical with those in SNS (i) to (xvi) and

they are briefly noted here, but the logical subroutines are

described in fuller detail.

(b) Boolean algebraic formulae for quantifiers in general

These formulae are equally valid for QL-1 and QL-2,

and are the basis for the implementation of quantifier states

in EVMF. As already mentioned, they utilize Boolean 3-vectors

and Boolean 3x3 matrices. The basic 16 subroutines in SNSALG

can be taken over practically unchanged for the QLOGIC also.

The names of these are listed below, but only brief details are

given, since they could be obtained by replacing the names of

vectors, matrices, and functions, of the form SXXXX, by QXXXXX,

and remembering that M = N = 2 for the former get changed to

M = N = 3 for the latter. The symbols adopted for inquiring

LOGIC are given in Table 1.

Table 1. Description of quantifier states in the EVMF program MATLOG

| Description | Name printed | Symbol in program | Description in BA-3 | Symbol in QBA | Description in standard logic |
|---|---|---|---|---|---|
| For all | ALL | Q1 | (1 0 0) | $\forall$ | $(\forall x)(\underline{\ge}x)$ |
| Not for all | NAL | Q2 | (0 1 1) | $\wedge$ | $\neg(\forall x)(\underline{\ge}x)$ |
| For some | SOM | Q3 | (0 1 0) | $\zeta$ | $(\exists x)(\underline{\ge}x)\wedge(\exists x)(\neg\underline{\ge}x$ |
| All or none | AON | Q4 | (1 0 1) | $\ominus$ | $(\forall x)(\underline{\ge}x)\vee(\forall x)(\neg\underline{\ge}x$ |
| Not exists | NEX | Q5 | (0 0 1) | $\phi$ | $\neg(\exists x)(\underline{\ge}x)$ |
| There exists | EXS | Q6 | (1 1 0) | $\exists$ | $(\exists x)(\underline{\ge}x)$ |
| Indefinite | IND | Q7 | (1 1 1) | $\triangle$ | $(\exists x)(\underline{\ge}x)\vee\neg(\exists x)(\underline{\ge}x)$ |
| Impossible or Contradictory | XXX | Q8 | (0 0 0) | $\emptyset$ | $(\exists x)(\underline{\ge}x)\wedge\neg(\exists x)(\underline{\ge}x)$ |

The matrix connectives QMZ(I,J) are 3x3 matrices, with

I,J = 1, 2, 3 and the quantified terms are 3-vectors QVA(I),

I = 1, 2, 3 which can be any one of the eight BA-3 states

Q1 to Q8. The states Q1, Q3, Q5 form the generators of the

BA-3 algebra and are called "basic states". The logical

connectives $\underline{A}(k, \ell)$, $\underline{O}(k, \ell)$, $\underline{E}(k, \ell)$, $\underline{I}(k, \ell)$ are defined

in a manner closely similar to the SNS $\underline{Z}(k, \ell)$ — e.g.

QVA,QK,QL = Z INIT(QK, QL) (see (vi) below)

The subroutines for the BA-3 operators (i) to (xvi)

are as follows.

(i) Boolean sum of two vectors    $a \oplus b = c$    ($\oplus = U$ = union)

OUNION (QVA,QVA) = QVC

BSUM  (QVA(1),QVB(1)) = QVC(1)

BSUM  (QVA(2),QVB(2)) = QVC(2)

BSUM  (QVA(3),QVB(3)) = QVC(3)

(ii) Boolean product of two vectors    $a \otimes b = c$    ($\otimes = V$ = vidya)

QVIDYA(QVA,QVB) = QVC

BPDT  (QVA(1),QVB(1)) = QVC(1)

BPDT  (QVA(2),QVB(2)) = QVC(2)

BPDT  (QVA(3),QVB(3)) = QVC(3)

(iii) Boolean complement of Q-vector    $a^c = b$

QCOMP(QVA) = QVB

BCMP (QVA(1)) = QVB(1)

BCMP (QVA(2)) = QVB(2)

BCMP (QVA(3)) = QVB(3)

(iv) Boolean sum of two matrices    $P \oplus Q = R$

QMXSUM(QMP,QMQ) = QMR

BSUM(QMP(I,J), QMQ(I,J)) = QMR(I,J), I, J = 1, 2, 3

(v) Boolean product of two matrices    $P \otimes Q = R$

QMXPDT(QMP,QMQ) = QMR

Replace BSUM in (iv) by BPDT .

(vi) Boolean complement of a matrix operator $\underset{\sim}{P}^c = \underset{\sim}{Q}$

$$QMCMP(QMP) = QMQ$$

$$BCMP(QMP(I,J)) = QMQ(I,J), \quad I = 1, 2, 3$$

(vii) Scalar product of two vectors

$$\langle \underset{\sim}{a} \mid \underset{\sim}{b} \rangle = \underset{\sim}{c} = a_1 \otimes b_1 \oplus a_2 \otimes b_2 \oplus a_3 \otimes b_3$$

$$QSCPDT(QVA,QVB) = BC \qquad \text{(Note Boolean scalar on the r.h.s)}$$

(viii) Unary product of a vector with a matrix

$$\langle \underset{\sim}{a} \mid \underset{\sim}{Z} \mid = \langle \underset{\sim}{b} \mid, \quad \bigoplus_i a_i \otimes Z_{ij} = b_j, \quad i,j = 1,2,3$$

$$QUNPDT(QVA,QMZ) = QVB$$

(ix) Binary product for the matrix with two vectors

$$\langle \underset{\sim}{a} \mid \underset{\sim}{Z} \mid \underset{\sim}{b} \rangle = c ; \quad \bigoplus_i \bigoplus_j a_i \otimes Z_{ij} \otimes b_j = c, \quad i,j = 1,2,3$$

$$QBINPT (QVA,QMZ,QVB) = BC \qquad \text{(Note Boolean scalar on the r.h.s)}$$

(x) Matrix product of two matrices $\quad \bigoplus_j P_{ij} \otimes Q_{jk} = R_{ik}, \quad i,j,k = 1, 2, 3$

$$QMATPT(QMP,QMQ) = QMR$$

(xi) Direct product of two vectors $\quad \underset{\sim}{a} \times \underset{\sim}{b} = \underset{\sim}{Z}$

$$QDIRPT(QVA,QVB) = QMZ$$

$$QMZ(I,J) = QPDT(QVA(I),QVB(J))$$

(xii) Direct sum of two vectors $\quad \underset{\sim}{a} + \underset{\sim}{b} = \underset{\sim}{Z}$

$$QDIRSM(QVA,QVB) = QMZ$$

$$QMZ(I,J) = BSM(QVA(I),QVB(J))$$

(xiii) Direct equivalence of two vectors $(\underset{\sim}{a} \equiv \underset{\sim}{b}) = Z$

$$QDIREQ(QVA,QVB) = QMZ$$

$$QMZ(I,J) = EQU(QVA(I),QVB(J))$$

(xiv) Scalar-vector direct product $a \times \underset{\sim}{b} = \underset{\sim}{c}$

$$QSVDPT(PA,QVB) = QVC \quad (\text{Note PA on the l.h.s})$$

$$QVC(I) = EQU(PA,QVB(I))$$

(xv) Scalar vector direct sum $a + \underset{\sim}{b} = \underset{\sim}{c}$

$$QSVDSM(PA,QVB) = QVC \quad (\text{Note PA on the l.h.s})$$

$$QVC(I) = EQU(PA,QVB(I))$$

(xvi) Transpose of a matrix $\underset{\approx}{P}^t = \underset{\approx}{Q}$

$$QTRANS(QMP) = QMQ$$

$$QMQ(I,J) = QMP(J,I), \quad I,J = 1, 2$$

## (c) Logical operators in QL-2

Similar to the SNS logical functions described in Section 2(b) of Part I, the corresponding QL-2 functions, using Q1 to Q8 for the constants, can be defined. These are as follows.

## (xvii) Matrix elements of logical connectives

$$\underset{\approx}{A}(k, \ell), \underset{\approx}{O}(k, \ell), \underset{\approx}{E}(k, \ell), \underset{\approx}{I}(k, \ell), \quad k,\ell = 1 \text{ to } 8$$

$$QMATEL(QMZ,QK,QL) = QMM(I,J)$$

Replace S by Q, and use QK,QL with K,L = 1 to 8 in (xvii) of SNSLOG.

For QMA,QK,QL, QMM = QDIRPT(QK,QL)

For QMO,QK,QL, QMM = QDIRSM(QK,QL)

For QME,QK,QL, QMM = QDIREQ(QK,QL)

For QMI,QK,QL, QMM = QDIREQ(QCMP(QK),QL)

(xviii) Relative truth value of one quantified term
for another : $\underline{t}(\underline{a} \mid \underline{b}) = (\langle a \mid b \rangle, \langle a \mid b^c \rangle) = (c_1, c_2) = \underline{c}$

QRELTV(QVA,QVB) = SVC     (Note SNS in r.h.s)

SVC(1) = QSCPDT (QVA,QVB)

SVC(2) = QSCPDT (QVA,QCOMP(QVB))

Input is in QL-2, output is in SNS


(xix) SNS truth value of a binary relation for given
inputs $\underline{a}, \underline{b}$ : $\underline{t}(\underline{a} \, Z \, \underline{b}) = \underline{c} = (c_1, c_2)$

$c_1 = \langle \underline{a} \mid Z \mid \underline{b} \rangle$ , $c_2 = \langle \underline{a} \mid Z^c \mid \underline{b} \rangle$

QSTV (QVA,QMZ,QVB) = SVC          (Note SNS in r.h.s)

SVC(1) = QBINPT(QVA,QMZ,QVB)

SVC(2) = QBINPT(QVA, QMXCMP(QMZ),QVB)


(xx) SNS truth value of a relation involving the QL-2
logical connective $Z(k, \ell)$ for inputs $\underline{a}$ , $\underline{b}$

$\underline{t}(\underline{a} \, \underline{Z} \, \underline{b}) = \underline{c}$ for $\underline{Z} = Z(k, \ell)$ via relative truth values

QSTV2(QVA,QMZ,QVB) = SVC          (Note SNS in r.h.s)

QMZ = QMA,QK,QL  etc. (for $\underline{Z}$ = $\underline{A}$, $\underline{O}$, $\underline{I}$, $\underline{E}$)

If QMZ = QMA,QK,QL

SVC(1) = BPDT(QSCPDT(QVA,QK), QSCPDT(QVB,QL))

SVC(2) = BSUM(QSCPDT(QVA,QCOMP(QK)),
                        QSCPDT(QVB,QCOMP(QL)))

If QMZ = QMO,QK,QL

SVC(1) = BSUM(QSCPDT(QVA,QK),QSCPDT(QVB,QL))

SVC(2) = BPDT(QSCPDT(QVA,QCCMP(SK)),QSCPDT(QVB,QCCMP(QL)))

IF QMZ = QME,QK,QL

    QKP = QCOMP(QK), QLP = QCOMP(QL)

    SVCA = QSTV2(QVA,QMA,QK,QL, QVB)

    SVCB = QSTV2(QVA,QMA,QKP,QLP, QVB)

    SVC(1) = BSUM(QVCA(1),QVCB(1))

    SVC(2) = BPDT(QVCA(2),QVCB(2))


IF QMZ = QMI,QK,QL

    QKP = QCOMP(QK)

    QMZ = QMO,QKP,QL


As mentioned for SNS logical connectives, in the case

of QL-2 also, general 3x3 matrices can/occur, for relations
<div style="text-align:center">also</div>

in QL-2 that cannot be expressed in one of the four forms

described above.


(xxi)   Unary relation for $\underline{Z}(\underline{k}, \underline{\ell})$ <u>via</u> relative truth value

---

$\langle a \mid q(k) \rangle\, Z\, \langle q(\ell) \rangle \mid = b$ , $Z = \otimes, \oplus$ , etc.

    QUNPT2(QVA,QMZ) = QVB

IF QMZ = QMA,QK,QL

    BAP = QSCPDT(QVA,QK)

    QVB = QSVPDT(BAP,QL)

IF QMZ = QMO,QK,QL

    BAP = QSCPDT(QVA,QK)

    QVB = QSVDSM(BAP,QL)

IF   QMZ = QMI,QK,QL

$$BAP = QSCPDT(QVA,QCOMP(QK))$$

$$QVB = QSVDSM(BAP,QL)$$

IF   QMZ = QME,QK,QL

$$QKP = QCOMP(QK), \quad QLP = QCOMP(QL)$$

$$QVBA = QUNPT2(QVA, QMA,QK,QL)$$

$$QVBB = QUNPT2(QVA, QMA,QKP,QLP)$$

$$QVB = QUNION(QVBA,QVBB)$$

Here also, the formulae closely follow those in SNS

except that  S  is to be replaced by Q, leaving B unchanged,

in SUNPT2.

(xxii)  Binary matrix-Boolean operator  G "agree" for checking
the equivalence (agreement) of two quantifier terms
a and b : (a)G(b) = c

$$QMBG(QVA,QVB) = QVC \qquad \text{(Note SNS in r.h.s)}$$

$$BCA = BEQU(QVA(1),QVB(1))$$

$$BCB = BEQU(QVA(2),QVB(2))$$

$$BCC = BEQU(QVA(3),QVB(3))$$

$$BCBP = BPDT(BCA,BCB)$$

$$SVC(1) = BPDT(BCBP,BCC)$$

$$SVC(2) = BCMF(QVC(1))$$

(xxiii) Matrix-Boolean unary operators E, N, M, L

In this case also, replace S by Q, and 1, 2, by 1, 2, 3.
The equations are

$$QEQU(QVA) = QVB$$

$$QVB(1) = QVA(1), \quad QVB(2) = QVA(2), \quad QVB(3) = QVA(3)$$

$$QNOT(QBA) = QVB$$

$$QVB(1) = QVA(3), \quad QVB(2) = QVA(2), \quad QVB(3) = QVA(1)$$

$$QCOMP(QVA) = QVB$$

$$QVB(1) = BCMP(QVA(1)), \quad QVB(2) = BCMP(QVA(2)),$$
$$QVB(3) = BCMP(QVA(3)).$$

$$QELL(QVA) = QVB$$

$$QVB(1) = BCMP(QVA(3)), \quad QVB(2) = BCMP(QVA(2)),$$
$$QVB(3) = BCMP(QVA(1)).$$

The above subroutines appear to be sufficient for
dealing with all problems containing quantifiers which involve
QL-2 algebra. However, for QL-1, quite different subroutines
are needed and are treated in Section 4 under PLOGIC.

## 2. Problems in QL-2 algebra

### Problem 5

Make direct checks of the algorithms Q(i) to Q(xvi) for

agreement with ideas discussed in MR-53 and 56, Lecture 3.

### Problem 6

Check the agreement between QUNPDT and QUNPT2 for the cases

listed in      Problem 7,     and print them side by side for inputs

QVA = Q1 to Q8.  Verify that there is agreement for Q1 to Q7

as input and note examples where they disagree for Q8 as input.

### Problem 7

Check the 8x8 truth tables for (1) Q(xix)— QSTV and

(2) Q(xx) — QSTV2  for the following connectives.

$$A(1,1), \ A(1,6), \ A(2,5), \ A(2,6), \ A(5,6), \ A(2,7), A(1,3), \ A(2,4)$$
$$Q(1,1), \ Q(1,6), \ Q(2,5), \ Q(2,6), \ Q(5,6), \ Q(2,7), Q(1,3), \ Q(2,4)$$
$$I(1,1), \ I(1,6), \ I(2,5), \ I(2,6), \ I(5,6), \ I(2,7), I(1,3), \ I(2,4)$$
$$E(1,1), \ E(1,6), \ E(2,5), \ E(2,6), \ E(5,6), \ E(2,7), E(1,3), \ E(2,4)$$

Check that they agree for the 7x7 sub-table involving Q1 to Q7

and if so, print the output for QSTV2 and note whether they

disagree for Q8 as input  a or b .

2. Problems in QL-2 algebra

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$$E(1,1), \ E(1,6), \ E(2,5), \ E(2,6), \ E(5,6), \ E(2,7), E(1,3), \ E(2,4)$$

Check that they agree for the 7x7 sub-table involving Q1 to Q7

and if so, print the output for QSTV2 and note whether they

disagree for Q8 as input $a$ or $b$ .

Problem 8

This will closely follow the pattern in Problem 4.

Problem 8A — The analogous case will not be worked out in QL-2, as a more general case will be considered in QLOGIC, and it will be proved that, in general, the conjunction of the Boolean truth values of two relations is not equivalent to the Boolean truth value of the conjunction of the two relations (represented by the Boolean product of the two matrices), while the corresponding theorem is true for disjunctions.

Therefore, the test is made only for the particular result that $E(k, \ell)$ can be expressed both as $A(k, \ell) \oplus A(k^c, \ell^c)$ and $I(k, \ell) \otimes I(k^c, \ell^c)$, via QSTV (Problem 8B), but not by QSTV2 (Problem 8C). These are analogous to Problems 4B and 4C in SNS for SSTV and SSTV2 respectively.

Problem 8B :— Define QMZ,QK,QL, for $Z = A, O, E, I$, from subroutine Q(xviii) and check the agreement of the three truth values SVV1, SVV2, SVV3 for $c$, obtained by the following MATLOG statements for $(a \mid E(k, \ell) \mid b) = c$, employing only the QSTV subroutine.

---

* Although MATLOG statements for Problems 8B and 8C were formulated in June 1987, they could be put on the computer only by the end of September 87 (in Hyderabad) and they fully supported the theoretical expectations stated below. (See also Appendix of MATLOG-3, MR-62 for a brief discussion). (Added 12.10.87

SVV1 = QSIV(QVA,QMZ,QVB)

where

QMZ = QQIREQ(QK,QL)

SVV2 = SSTV(SVA1,SMO,S1,S1,SVA2)

where

SVA1 = QSIV(QVA,QMA,QK,QL,QVB)

SVA2 = QSIV(QVA,QMA,QCOMP(QK),QCOMP(QL), QVB)

SVV3 = SSTV(SVI1,SMA,S1,S1,SVI2)

where

SVI1 = QSTV(QVA,QMI,QK,QL,QVB)

SVI2 = QSTV(QVA,QMI,QCOMP(QK),QCOMP(QL),QVB)

Check SVV1, SVV2, SVV3 for agreement and if found true, print

only SVV1. (The test is to be done for all $k, \ell$ = 1 to 8, for

$k \leq \ell$).

Problem 8C : — Check the agreement of QSTV2(QVA,QME,QK,QL,QVB)

= SVC with the two different, but equivalent, definitions of

$E(k, \ell)$, following Problem 4C in SNS. The relevant equations

are as follows:

SVC1 = SSTV2(SVA1,SMO,S1,S1,SVA2)

where

SVA1 = QSTV2(QVA,QMA,QK,QL,QVB)

SVA2 = QSTV2(QVA,QMA,QCOMP(QK),QCOMP(QL),QVB)

$$SVC2 = SSTV2(SVI1,SMA,S1,S1,SVI2)$$

where

$$SVI1 = QSTV2(QVA,QMI,QK,QL,QVB)$$

$$SVI2 = QSTV2(QVA,QMI,QCOMP(QK),QCOMP(QL), QVB)$$

Print both SVC1 and SVC2, and check if they agree for all

a , b = Q1 to Q7, and also with SVV1 of Problem 8B, and note

where they disagree for a or b = Q8. (Test is to be repeated

for all (k, ℓ), with k, ℓ = 1 to 8, k ≤ ℓ).


3. Comments on the outputs from the problems in Section 2

The outputs for the functions given below are illustrative

of the nature of these functions. Note that Q8 is printed as

IMP and will be changed later to XXX.


Problem 5

It will be noticed from Table 1 that the four matrix-Boolean

operators QMBE, QMBN, QMBM, QMBL permute the four standard

quantifier states ALL, EXS, NAL, NEX among themselves, and

similarly interchange pairwise the remaining four states

(SOM, AON), (IND, IMP). Hence the collection of eight states

of BA-3 employed in quantifier logic are left invariant by the

operation of Boolean-matrix operators.

also
The operators QUNION and QVIDYA are/seen to produce

as outputs only one of the eight quantifier states, for all

combinations of the input states $(\underset{\sim}{a}, \underset{\sim}{b})$, and therefore, in

QLOGIC, the algebra is complete under "Boolean" operators.

The scalar product of two Q-vectors and the SNS relative

truth value of these are also listed in Table 1, for $\underset{\sim}{k}, \underset{\sim}{\ell} = 1$ to 8.

The former is a Boolean scalar 1 or 0 , and the latter a

Boolean SNS vector T, F, D or X. It can be verified that the

first component $c_\alpha$ of the SNS truth value $\underline{c} = (c_\alpha, c_\beta)$ for

the latter is always equal to the Boolean number corresponding

to QSCPDT, in all cases. This indicates the consistency of

algebra
our quantifier, employing BA-2 truth values in SNS, with the

classical quantifier calculus, using only BA-1 truth values.

The BA-2 calculus is analytical continuation, for logical

truth values T, F, D, X, of the BA-1 calculus employing only

T and F, and the latter always leads to the former if the

states D and X are not taken as independent logical truth values,

but expressible as $D = T \lor F$ and $X = T \land F$. (See MR-52 for

a treatment of this aspect.)

## Table 1. Outputs for some typical functions in QLOGIC

|     | QCOMP | QMBE | QMBN | QMBM | QMBL |
|-----|-------|------|------|------|------|
| ALL | NAL   | ALL  | NEX  | NAL  | EXS  |
| NAL | ALL   | NAL  | EXS  | ALL  | NEX  |
| SOM | AON   | SOM  | SOM  | AON  | ACN  |
| AON | SOM   | AON  | AON  | SOM  | SCM  |
| NEX | EXS   | NEX  | ALL  | EXS  | NAL  |
| EXS | NEX   | EXS  | NAL  | NEX  | ALL  |
| IND | IMP   | IND  | IND  | IMP  | IMP  |
| IMP | IND   | IMP  | IMP  | IND  | IND  |

### QUNION

| ALL | IND | EXS | AON | AON | EXS | IND | ALL |
|-----|-----|-----|-----|-----|-----|-----|-----|
| IND | NAL | NAL | IND | NAL | IND | IND | NAL |
| EXS | NAL | SOM | IND | NAL | EXS | IND | SOM |
| AON | IND | INC | AON | AON | IND | IND | AON |
| AON | NAL | NAL | AON | NEX | IND | IND | NEX |
| EXS | IND | EXS | IND | INC | EXS | IND | EXS |
| IND | IND | INC | IND | INC | IND | IND | INC |
| ALL | NAL | SOM | AON | NEX | EXS | IND | IMP |

### QVIDYA

| ALL | IMP | IMP | ALL | IMP | ALL | ALL | IMP |
|-----|-----|-----|-----|-----|-----|-----|-----|
| IMP | NAL | SOM | NEX | NEX | SCM | NAL | IMP |
| IMP | SCM | SOM | IMP | IMP | SCM | SCM | IMP |
| ALL | NEX | IMP | AON | NEX | ALL | ACN | IMP |
| IMP | NEX | IMP | NEX | NEX | IMP | NEX | IMP |
| ALL | SCM | SOM | ALL | IMP | EXS | EXS | IMP |
| ALL | NAL | SOM | AON | NEX | EXS | IND | IMP |
| IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

### QSCPD1

| 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 1 | 1 | 1 | 1 | 1 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 1 | 1 | 1 | 0 |
| 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### QRELTV

| T | F | F | T | F | T | T | F |
|---|---|---|---|---|---|---|---|
| F | T | D | D | D | D | T | F |
| F | T | T | F | F | T | T | F |
| D | D | F | T | D | D | T | F |
| F | T | F | T | T | F | T | F |
| D | D | D | D | F | T | T | F |
| D | D | D | D | D | D | T | F |
| X | X | X | X | X | X | X | X |

Problem 6

Four illustrations each for $\underset{\sim}{A}(k, \ell)$, $\underset{\sim}{Q}(k, \ell)$, $\underset{\sim}{E}(k, \ell)$, $\underset{\sim}{I}(k, \ell)$ are copied from the computer outputs obtained for this problem. It was verified that QUNPDT and QUNPT2 give identical results for $\underset{\sim}{a}$ = QVA = Q1 to Q7 for all examples. For Q8 as input QVA, however, $\underset{\sim}{b}$ = QVB = Q8 = IMP for $\underset{\sim}{A}(k, \ell)$ and $\underset{\sim}{E}(k, \ell)$, while QVB = QL for $\underset{\sim}{Q}(k, \ell)$ and $\underset{\sim}{I}(k, \ell)$, for QUNPT2. On the other hand, in the case of QUNPDT, QVB = Q8 whenever QVA = Q8. Hence only QUNPT2 results are reproduced, and, where they disagree for QUNPDT, they are marked by a ring.

Table 2. Typical examples of the output of QUNPT2 for QVA = 1 to 8

| Input QVA | Output QVB for | | | |
|---|---|---|---|---|
| | $\underset{\sim}{A}(1, 1)$ | $\underset{\sim}{A}(1, 6)$ | $\underset{\sim}{A}(2, 5)$ | $\underset{\sim}{A}(2, 7)$ |
| ALL | ALL | EXS | IMP | IMP |
| NAL | IMP | IMP | NEX | IND |
| SOM | IMP | IMP | NEX | IND |
| AON | ALL | EXS | NEX | IND |
| NEX | IMP | IMP | NEX | IND |
| EXS | ALL | EXS | NEX | IND |
| IND | ALL | EXS | NEX | IND |
| IMP | IMP | IMP | IMP | IMP |
| | $\underset{\sim}{Q}(1, 1)$ | $\underset{\sim}{Q}(1, 6)$ | $\underset{\sim}{Q}(2, 5)$ | $\underset{\sim}{Q}(2, 7)$ |
| ALL | IND | IND | NEX | IND |
| NAL | ALL | EXS | IND | IND |
| SOM | ALL | EXS | IND | IND |
| ADN | IND | IND | IND | IND |
| NEX | ALL | EXS | IND | IND |
| EXS | IND | IND | IND | IND |
| IND | IND | IND | IND | IND |
| IMP | (ALL) | (EXS) | (NEX) | (IND) |

Table 2. Contd.

| Input QVA | Output QVB for $\underline{E}(1, 1)$ | $\underline{E}(1, 6)$ | $\underline{E}(2, 5)$ | $\underline{E}(2, 7)$ |
|---|---|---|---|---|
| ALL | ALL | EXS | EXS | IMP |
| NAL | NAL | NEX | NEX | IND |
| SOM | NAL | NEX | NEX | IND |
| AON | IND | IND | IND | IND |
| NEX | NAL | NEX | NEX | IND |
| EXS | IND | IND | IND | IND |
| IND | IND | IND | IND | IND |
| IMP | IMP | IMP | IMP | IMP |
| | $\underline{I}(1, 1)$ | $\underline{I}(1, 6)$ | $\underline{I}(2, 5)$ | $\underline{I}(2, 7)$ |
| ALL | ALL | EXS | IND | IND |
| NAL | IND | IND | NEX | IND |
| SOM | IND | IND | NEX | IND |
| AON | IND | IND | IND | IND |
| NEX | IND | IND | NEX | IND |
| EXS | IND | IND | IND | IND |
| IND | IND | IND | IND | IND |
| IMP | (ALL) | (EXS) | (NEX) | IND |

Even these disagreements follow a rule as mentioned above.

Thus, they occur only for the O-type connectives QMO and QMI,

and irrespective of $(\underline{k}, \ell)$, the output for QUNPT2 is QVB = QL, for

QVA = Q8 = IMP. This is similar to the behaviour in SNS algebra,

as may be seen from Table 1B where for $\underline{O}(k, \ell)$ and $\underline{I}(k, \ell)$,

SVA = S4 = X always leads to SL as output for SVB. The

explanation of this will be given, from Boolean algebra, under

even more general conditions of GLOGIC, in Part III, and it

reduces to SNS and QLOGIC for M = N = 2, and 3 respectively.

Problem 7

In this case of binary relation also, the SNS truth

values calculated using QSTV and QSTV2 agree for the 7x7

sub-table employing Q1 to Q7 for $\underset{\sim}{a}$ and $\underset{\sim}{b}$. They, however,

disagree for some examples if one of either $\underset{\sim}{a}$ or $\underset{\sim}{b}$ is Q8.

Also, as in SNS, the disagreeing examples have the

truth value F, for QSTV2, for the connectives $\underset{\sim}{A}(\underset{\sim}{k}, \underset{\sim}{\ell})$ and

$\underset{\sim}{E}(\underset{\sim}{k}, \underset{\sim}{\ell})$, while they have the truth value T for the connectives

$\underset{\sim}{O}(\underset{\sim}{k}, \underset{\sim}{\ell})$ and $\underset{\sim}{I}(\underset{\sim}{k}, \underset{\sim}{\ell})$,/ as against X for the corresponding outputs for QSTV. This behaviour can also be proved

still more generally for similar logical connectives defined

in GLOGIC, and a description and proof are given in Part III

dealing with GLOGIC.   Print outs for QSTV2 in some typical

cases are given in Table 2, and the disagreeing outputs are

marked by rings.

Table 3. SNS output of QSTV2 for the 8x8 array of inputs QVA, QVB.

A(1.6)

| T | D | T | D | F | T | D | X |
|---|---|---|---|---|---|---|---|
| P | F | F | F | F | F | F | (F) |
| F | F | F | F | F | F | F | (F) |
| D | D | D | D | F | D | D | (F) |
| F | F | F | F | F | F | F | (F) |
| D | D | D | D | F | D | D | (F) |
| D | D | D | D | F | D | D | (F) |
| X | (F) | X | (F) | (F) | X | (F) | X |

A(5,6)

| F | F | F | F | F | F | F | (F) |
|---|---|---|---|---|---|---|---|
| D | D | D | D | F | D | D | (F) |
| F | F | F | F | F | F | F | (F) |
| D | D | D | D | F | D | D | (F) |
| T | D | T | D | F | T | D | X |
| F | F | F | F | F | F | F | (F) |
| D | D | D | D | F | D | D | (F) |
| X | (F) | X | (F) | (F) | X | (F) | X |

O(1,6)

```
T   T   T   T   T   T   T  (T)
T   D   I   D   F   T   D   X
T   D   I   D   F   T   D   X
T   D   T   D   D   T   D  (T)
T   D   T   D   F   T   D   X
T   D   I   D   D   T   D  (T)
T   D   T   D   D   T   D  (T)
(T) (T) (T) (T)  X  (T) (T)  X
```

O(2,5)

```
F   D   F   D   T   F   D   X
T   T   T   T   T   T   T  (T)
T   T   T   T   T   T   T  (T)
D   D   D   D   T   D   D  (T)
I   T   T   T   T   T   T  (T)
D   D   D   D   T   D   D  (T)
D   D   D   D   T   D   D  (T)
X  (T)  X  (T) (T)  X  (T)  X
```

E(1,1)

```
T   F   F   D   F   D   D   X
F   T   T   D   T   D   D   X
F   T   T   D   T   D   D   X
D   D   D   D   D   D   D  (F)
F   T   T   D   T   D   D   X
D   D   D   D   D   D   D  (F)
D   D   D   D   D   D   D  (F)
X   X   X  (F)  X  (F) (F)  X
```

E(5,6)

```
F   D   F   D   T   F   D   X
D   D   D   D   D   D   D  (F)
F   D   F   D   T   F   D   X
D   D   D   D   D   D   D  (F)
T   D   T   D   F   T   D   X
F   D   F   D   T   F   D   X
D   D   D   D   D   D   D  (F)
X  (F)  X  (F)  X   X  (F)  X
```

I(1,1)

```
F   F   F   D   F   D   D   X
T   T   T   T   T   T   T  (T)
T   T   T   T   T   T   T  (T)
T   D   D   D   D   D   D  (T)
T   T   T   T   T   T   T  (T)
T   D   D   D   D   D   D  (T)
T   D   D   D   D   D   D  (T)
(T)  X   X  (T)  X  (T) (T)  X
```

I(1,6)

```
T   D   T   D   F   T   D   X
T   T   T   T   T   T   T  (T)
T   T   T   T   T   T   T  (T)
T   D   T   D   D   T   D  (T)
T   T   T   T   T   T   T  (T)
T   D   T   D   D   T   D  (T)
T   D   T   D   D   T   D  (T)
(T) (T) (T) (T)  X  (T) (T)  X
```

It has been possible to deduce the following general
rule for the ringed entries that disagree with QSTV in Table 1.

$$
\left.
\begin{array}{l}
\underset{\sim}{a} = \emptyset, \ \underset{\sim}{b} \cap q(\ell) \neq \emptyset \\
\text{or} \\
\underset{\sim}{b} = \emptyset, \ \underset{\sim}{a} \cap q(k) \neq \emptyset
\end{array}
\right\} \longmapsto \quad (\underset{\sim}{a} \mid \underset{\sim}{Q}(\underset{\sim}{k}, \ell) \mid \underset{\sim}{b}) = T
\qquad (1)
$$

$$
\left.
\begin{array}{l}
\underset{\sim}{a} = \emptyset, \ \underset{\sim}{b} \cap q(\ell^c) \neq \emptyset \\
\text{or} \\
\underset{\sim}{b} = \emptyset, \ \underset{\sim}{a} \cap q(k^c) \neq \emptyset
\end{array}
\right\} \longmapsto \quad (\underset{\sim}{a} \mid A(\underset{\sim}{k}, \ell) \mid \underset{\sim}{b}) = F
\qquad (2)
$$

$$
\left.
\begin{array}{l}
\underset{\sim}{a} = \emptyset, \ \underset{\sim}{b} \cap q(\ell) \neq \emptyset \\
\text{or} \\
\underset{\sim}{b} = \emptyset, \ \underset{\sim}{a} \cap q(k^c) \neq \emptyset
\end{array}
\right\} \longmapsto \quad (\underset{\sim}{a} \mid \underset{\sim}{I}(\underset{\sim}{k}, \ell) \mid \underset{\sim}{b}) = T
\qquad (3)
$$

In fact, it can be verified that the data given in Table 1
of MR-60 obey the corresponding rules for SNSLOG. Thus the
condition $\underset{=}{a} \supseteq s(k)$ found there, has the property given
in (4) below and the right hand side is analogous to the
conditions employed in the above rules (1), (2), (3).

$$
(\underset{=}{a} \supseteq s(k)) \subseteq (\underset{=}{a} \cap s(k) \neq \emptyset)
\qquad (4)
$$

The conditions $\underset{\sim}{a} \cap q(k) \neq \emptyset$ and $\underset{\sim}{a} \cap q(k^c) \neq \emptyset$
are expressible in MATLOG as

$$
\text{QSCPDT}(\text{QVA},\text{QK}) = 1 \ , \quad \text{QSCPDT}(\text{QVA},\text{QCOMP}(\text{QK})) = 1
\qquad (5)
$$

We shall not prove the above rules (1), (2), (3) for

3x3 matrices in QLOGIC, but prove them for the even more

general case of m x n matrices of GLOGIC, employing clausal

relations, which are the analogues of the equations employed

in QLOGIC. This will be done by showing that the Boolean

truth value

$$QBTV(QVA,QMZ,QVB) \equiv QBINPT(QVA,QMZ,CVB) \qquad (6)$$

for QMZ = QMO, QK, QL, differs from

$$QBTV2(QVA,QMO,QK,QL,QVB)$$
$$\equiv BSUM(QSCPDT(QVA,QK),QSCPDT(QVB,QL)) \qquad (7)$$

for precisely the above conditions, namely

$$QVA = Q8, \quad QSCPDT(QVB,QL) = 1$$
or $\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad (8)$
$$QSCPDT(QVA,QK) = 1, \quad QVB = Q8$$

The analogues in QL-2 of Problems 2 and 3 of SNS in

Part I will not be given here, but will be briefly taken up

at the end of this report, after discussing the implementation

of elementary relations in PLOGIC(QL-1A) and PQLOG (QL-1B).

These are to be contrasted with a quantified relation expressed in QL-1. In/logical relation of the type QL-1, the relation really exists between the individual components $\underline{a}x$ and $\underline{b}x$ for each value of the variable $x$, and both $\underline{a}x$ and $\underline{b}x$ have the same domain of operation. This is particularly made clear by the form of the relation in QL-1 in general, namely

$$(\underline{q}_Z x)(\underline{a}x \; \underline{Z} \; \underline{b}x) \tag{2}$$

where there is only a single variable $x$ contained in the relation and the form of its unary and binary implementation is as in (3a), (3b) below:

Unary : $\underline{a}x$ , $(\underline{q}_Z x)(\underline{a}x \; \underline{Z} \; \underline{b}x) \longmapsto \underline{b}x$ $\qquad$ (3a)

Binary : $\underline{a}x$, $\underline{b}x$, $(\underline{q}_Z x)(\underline{a}x \; \underline{Z} \; \underline{b}x) \longmapsto c$ , SNS truth value of the relation for the given inputs. (3b)

A few simple examples of unary and binary relations are given in (4a-d) and (5a-c) below. The unary examples are

$$(\forall x)(\underline{a}x), \; (\exists x)(\underline{a}x \Rightarrow \underline{b}x) \longmapsto (\exists x)(\underline{b}x)$$

$$(\exists x)(\underline{a}x), \; (\exists x)(\underline{a}x \Rightarrow \underline{b}x) \longmapsto (\triangle x)(\underline{b}x) \qquad (4a\text{-}d)$$

$$(\exists x)(\underline{a}x), \; (\forall x)(\underline{a}x \wedge \underline{b}x) \longmapsto (\forall x)(\underline{b}x)$$

$$(\exists x)(\underline{a}x), \; (\exists x)(\underline{a}x \wedge \underline{b}x) \longmapsto (\exists x)(\underline{b}x)$$

Similarly, we have the binary examples

$$(\forall x)(\underline{a}x), \quad (\forall x)(\neg \underline{1}\underline{b}x), \quad (\exists x)(\underline{a}x \wedge \underline{b}x) \vdash \Rightarrow \underline{c} = F$$

$$(\exists x)(\underline{a}x), \quad (\exists x)(\neg \underline{1}\underline{b}x), \quad (\forall x)(\underline{a}x \wedge \neg \underline{1}\underline{b}x) \vdash \Rightarrow \underline{c} = D \qquad (5a\text{-}c)$$

$$(\exists x)(\underline{a}x), \quad (\forall x)(\underline{b}x), \quad (\exists x)(\underline{a}x \Rightarrow \underline{b}x) \vdash \Rightarrow \underline{c} = T$$

As will be seen from the form of the expressions in (3a) and

(3b), both vectors representing quantifiers, as well as SNS terms,

occur in these formulae, and it is not apparent as to how they can

be converted into BVMF containing essentially 3-vectors. However,

as has been briefly outlined in MR-56, Lecture-4, this class of

equations, which has been named QL-1A, can be formulated in a

compact and unified treatment, by first considering the quantifier

algebra QL-1, which is effectively equivalent to QL-1A for the

particular case when $(\underset{\sim Z}{q_Z}x) = (\forall x)$. Then Eqs. 3(a,b) respectively

go over into 6(b) and 6(a) below. We shall discuss these and

the nature of their relationship with 3(a,b) below.

$$\underline{a}x, \quad \underline{b}x, \quad \underline{a}x \underset{=}{\leq} \underline{b}x = \underline{c}x \longmapsto \underline{c}x \qquad \text{(Binary forward)} \qquad (6a)$$

$$\underline{c}x, \quad \underline{a}x, \quad \underline{a}x \underset{=}{\leq} \underline{b}x = \underline{c}x \longmapsto \underline{b}x \qquad \text{(Binary reverse)} \qquad (6b)$$

As has been described in MR-52, 53 and 54, the QL-1 type of

relation occurs from the existence of a logical relation between

the individual components $\underline{a}x$ and $\underline{b}x$ of the corresponding members

of the sets A and B having the quantifier states $\underline{a}x$ and $\underline{b}x$ ,

which lead to $\underline{c}x$ of the same individual x ,

and hence to the resultant set C described by the quantifier

state $cx$. Thus the implementation of the binary relation (7)

$$( \exists x)(ax), \ (\forall x)(bx), \ ax \wedge bx = cx \tag{7}$$

leads to $cx = (\exists x)(cx)$ for the following reasons. We are

given that "at least one member " $(\exists x)$ of the full set $\mathscr{S}$ has

the property of the set A, and that "all members" $(\forall x)$ of the

set $\mathscr{S}$ have the property of the set B, and we wish to find out

the quantified nature of the members of $\mathscr{S}$ that have the property

of both Set A and Set B. Denoting this set by C and the

corresponding vector by $cx$ , it is obvious that we can only say

that "there exists at least one member" $(\exists x)$ of $\mathscr{S}$ that have

the property of the set C. Thus a binary relation in QL-1 has

the individual output $cx$ in an SNS state, but this leads to

a quantifier state represented by a 3-vector $cx$ which is of

the same type as the quantifier states of the two inputs $ax$ and $bx$.

The principle of the procedure discussed above can be

used for working out the resultant quantifier state $cx$ of the

set C which is the logical sum, or product, of the quantifier states

$ax$ and $bx$, and more generally for a relation $z(k, \ell)$ connecting

gx and bx of sets A and B for their quantified states a and b being /any/ one of the eight standard states q(1) to q(8). This has been done ab initio and 3x3 tables for these operators have been obtained.

Table 4. 3x3 tables for the QL-1 connectives "and" and "or"

(a)   AND (A)                    (b)   OR (O)

| A | ∀ | ζ | ⊕ |
|---|---|---|---|
| ∀ | ∀ | ζ | ⊕ |
| ζ | ζ | ∧ | ⊕ |
| ⊕ | ⊕ | ⊕ | ⊕ |

| O | ∀ | ζ | ⊕ |
|---|---|---|---|
| ∀ | ∀ | ∀ | ∀ |
| ζ | ∀ | Ǝ | ζ |
| ⊕ | ∀ | ζ | ⊕ |

The detailed mathematical properties of these tables are discussed in MR-53,54, and a short summary given in MR-56, Lecture 4. Here, we shall only indicate those formulae which are relevant for working out the computer-implementable algorithms given below. As will be seen from Table 4(a, b), the outputs for the QL-1 relations gx A bx and gx O bx for the generator states q(1), q(3), q(5) have very simple representation in lattice algebra. Without going into this lattice theory as such, we may merely write the effective

algebraic formulae for $\underset{\sim}{c}x$ in terms of $\underset{\sim}{a}x$ and $\underset{\sim}{b}x$ as follows.

Denoting the states of $\underset{\sim}{a}$, $\underset{\sim}{b}$ and $\underset{\sim}{c}$ as $q(i)$, $q(j)$ and $q(k)$, we have

QL-1 :-: "and"

$$q(i) \; \underline{\underline{A}} \; q(j) = q(k)$$

where

$k = \text{Max}\,(i, j)$, and $k = 2$ if $i = j = 3$
and $\quad k = 8$ if $i$ or $j = 8$

QL-1 -:- "or"

$$q(i) \; \underline{\underline{O}} \; q(j) = q(k)$$

where

$k = \text{Min}\,(i, j)$, and $k = 6$ if $i = j = 3$
and $k = 8$ if $i$ or $j = 8$

The extension of these 3x3 tables into 8x8 tables is obtained

from noting that, a general quantifier state, $q(n)$, $n = 1$ to $8$,

is the Boolean sum of at most three basic states $q(1)$, $q(3)$, $q(5)$.

We therefore make use of the analogs of the formula (8) given

below, for obtaining the QL-1 conjunction or disjunction of

these mixed states. Thus, for $\underset{\sim}{a} = q(i_1) \oplus q(i_2)$ and

$\underset{\sim}{b} = q(j_1) \oplus q(j_2)$, $\underset{\sim}{a} \; \underline{\underline{A}} \; \underset{\sim}{b} = \underset{\sim}{c}$ is

$$\underset{\sim}{c} = (q(i_1) \oplus q(i_2)) \; \underline{\underline{A}} \; (q(j_1) \oplus q(j_2))$$

$$= (q(i_1) \; \underline{\underline{A}} \; q(j_1)) \oplus (q(i_2) \; \underline{\underline{A}} \; q(j_1)) \oplus (q(i_1) \; \underline{\underline{A}} \; q(j_2))$$

$$(q(i_2) \; \underline{\underline{A}} \; q(j_2)) \quad (8$$

The tables so obtained for $\underline{\underline{Z}} = \underline{\underline{A}}(1, 1)$ and $\underline{\underline{O}}(1, 1)$ are given in

page 12 of Lecture-4, MR-56. They will be reproduced as computer

outputs later in this report.

The natural question arises as to how we should proceed

to formulate the algorithms for the connectives $\underline{\underline{I}}$ and $\underline{\underline{E}}$ and

also for a general logical connective $\underline{Z}(k, \ell)$. Both these

are possible by noting that in QL-1, the following general

formula holds for $k, \ell = 1, 2$.

$$\underline{a}x \ \underline{\underline{Z}}(k, \ell) \ \underline{b}x \ = \ \underline{a}'x \ \underline{\underline{Z}}(1, 1) \ \underline{b}'x \tag{9}$$

where

$$\underline{a}'x = \underline{a}x, \ \underline{a}^n x \quad \text{according as } k = 1, 2$$
$$\underline{b}'x = \underline{b}x, \ \underline{b}^n x \quad \text{according as } \ell = 1, 2$$

The main point to be noted is that the appropriate operators

that are to be applied to $\underline{a}x$ and $\underline{b}x$, depending on the value

of $k, \ell = 1$ or $2$, are the matrix-Boolean operators QEQU and

QNOT respectively. Particularly, the occurrence of the

operator $\underline{N}$ is to be noted for QL-1, as contrasted with the

operator $\underline{M}$ which occurs in similar situations in QL-2.

The theoretical reason for this is the fact that it is the

predicate $\underline{a}x$ in $(qx)(\underline{a}x) = \underline{a}x$ that is negated and this

produces from $\neg \underline{a}x$ the quantifier state $\underline{a}x \ \underline{N}$ which is the

same as the MATLOG function QNOT(QVA).

With this introduction to the algebraic theory, we shall

formulate some of the essential functions required in PLOGIC

which are to be formulated in the QL-1 algebra described above.

(b) MATLOG functions for QL-1

The two operators QEQU and QNOT are particularly needed

in PLOGIC as they occur very frequently in it. Also, two

functions, namely QELEM and QVECEL, have to be explicitly

stated although they have been used in all the subroutines

in QLOGIC. These are given in (iii).

(iii) Decomposition of a vector QVA into its basic vectors QELEM

QELEM(QVA, QBAS,IB)

This program takes in as input a vector QVA having

components VA(1), VA(3), VA(5) which are given by VECEL(QVA),

and gives as output one, two, or three, basic vectors, out of

(1 0 0), (0 1 0), (0 0 1) — which are components of it — as

QBAS(I), I = 1, IB, where IB is the number of basic vectors

contained in QVA. It has the following structure.

$$QVECEL(QVA) = VA(3)$$

$$IB = 0$$

If   VA(1) = 1,   IB = IB + 1,   QBAS(IB) = Q1

If   VA(2) = 1,   IB = IB + 1,   QBAS(IB) = Q3

If   VA(3) = 1,   IB = IB + 1,   QBAS(IB) = Q5

Further, to take care of the lattice properties, two

subroutines QMIN and QMAX are also needed.

(iv)  Calculation of the minimum of two basic vectors QMIN

This function has the form  QMIN(QVA,QVB), where  QVA = QI
and QVB = QJ, are the inputs.  The output is  QMIN = QK,  where
K = min (I,J), the function  min  standing for the arithmetical
minimum of  I and J.

(v)   Determination of the maximum of two vectors  QMAX

As with QMIN, this has the form  QMAX(QVA,QVB).  With
the same inputs  QI,  QJ,  QMAX = QK, where  K = max(I,J).

It is to be noted that the minimum and maximum have to do
only with the indexes of the  Q-vectors  QVA,QVB  in the standard
form  Q1 to Q8  adopted by us, and the only values that occur
for  I, J, K  are 1, 3 and 5, since these functions are required /
only for basic vectors.

Unlike QLOGIC, there are not two different operators
U, Q  and  V, A  in PLOGIC, but a combination of both types
of functions occurs in the QL-1 operators "or" and "and" as
mentioned above.  However, these two are given the MATLOG
names PUNION and PVIDYA and the algebraic properties mentioned
above can be converted into the subroutines given below.

(vi)  QL-1 "or" — ax $\underline{\underline{O}}$ bx = cx : PUNION

PUNION(QVA,QVB) = QVC

QELEM(QVA) = QBAS1(IB1)

QELEM(QVB) = QBAS2(IB2)

IF  (QVA = Q8) . OR . (QVB = Q8)  QOUT = Q8

GO TO 1

ELSE  QOUT = Q8

DO  I = 1, IB1,  J = 1, IB2

10  QOUT = QUNION(QOUT,QMIN(QBAS1(I), QBAS2(J))

GO TO 1

1  QVC = QOUT

(vii)  QL-1 "and" — ax $\underline{\underline{A}}$ bx = cx : PVIDYA

PVIDYA(QVA,QVB) = QVC

Same as PUNION except that QMIN is replaced
by QMAX in 10 as

20  QOUT = QUNION(QOUT,QMAX(QBAS1(I), QBAS2(J))

Employing these two QL-1 connectives, the other two
connectives "imply" and "equivalent" can be simply defined
in QL-1 as in (10a), (10b). Hence no separate functions are
defined for these and they are calculable from the general
subroutine PBINPT given in (viii) below.

$$q(i) \, \underline{\underline{I}} \, q(j) \;=\; q(i^n) \, \underline{\underline{O}} \, q(j) \quad ; \tag{10a}$$

$$q(i) \, \underline{\underline{E}} \, q(j) \;=\; (q(i) \, \underline{\underline{A}} \, q(j)) \, \underline{\underline{O}} \, (q(i^n) \, \underline{\underline{A}} \, q(j^n)) \tag{10b}$$

(The corresponding Eq.(8) on page 11 of Lecture-4 is to be
replaced by (10a) and (10b) as being more precise.)

Combining (9), for dealing with a general $\underline{z}(k, \ell)$, with (10)

to extend the theory from $\underline{A}(1, 1)$ and $\underline{Q}(1, 1)$ to $\underline{E}(1, 1)$ and

$\underline{I}(1, 1)$, the algorithm for a general relation $\underline{a}x \; \underline{z}(k, \ell) \; \underline{b}x$

in the binary form of implementation can be made, as given in

the algorithm for the function PB1NPT below. As mentioned

earlier, the output of this binary product in QL-1 is a 3-vector

and not a Boolean scalar as in SKS and QL-1. Consequently, the

function has the form given in (viii) below.

(viii)   QL-1 Binary product —  $\underline{a}x \; \underline{z}(k, \ell) \; \underline{b}x = \underline{c}x$ : PBINPT

            PBINPT(QVA,SMX,SK,SL,QVB) = QVC

               SK = 1  ⟶  QOK = PEQU

               SK = 2  ⟶  QOK = PNOT

               SL = 1  ⟶  QOL = PEQU

               SL = 2  ⟶  QOL = PNOT

         QVC = PVIDYA(QOK(QVA), QOL(QVB))    for  SMX = SMA

         QVC = PUNION(QOK(QVA), QOL(QVB))    for  SMX = SMO

         QVC = PUNION(PNOT(QOK(QVA)), QOL(QVB))  for SMX = SMI

         QVC = PUNION(PVIDYA(QOK(QVA),QOL(QVB)) ,

            PVIDYA(PNOT(QOK(QVA)), PNOT(QOL(QVB))))  for  SMX = SME

The above are the subroutines which are required under

QL-1, which is not practically applied as such, but theoretically

very valuable for QL-1A calculus which is the standard form

employed in predicate calculus. It will be noticed that a

general SNS matrix SMZ is not defined in our QL-1 algebra,

but only the logical matrices $\underline{Z}(k, \ell)$ in SNS, namely SMX,SK,SL

with X = $\underline{A}$, $\underline{O}$, $\underline{E}$, $\underline{I}$ , are employed in the subroutines mentioned

above for PLOGIC. In fact, only these matrices are used in the

definition of logical relations in standard predicate calculus

which is our QL-1A algebra. (The extension to all 16 2x2

matrices $\underline{Z}$ can be done, but has not been included here,

since a very different formalism for implementing PLOGIC may

have to be employed for this purpose.) However, it should be

mentioned that the formulae given above are the analogs of

SSTV and QSTV, rather than SSTV2 and QSTV2, since as will be

seen from the outputs in Section 4(e) , the output of a binary

relation with $\underline{a}$ = $\emptyset$ , or $\underline{b}$ = $\emptyset$ in $\underline{a} \, \underline{Z} \, \underline{b} = \underline{c}$ is always $\underline{c}$ = $\emptyset$

for both $\underline{A}$-type and $\underline{O}$-type connectives $\underline{Z}$, (and naturally also

for $\underline{E}$-type and $\underline{I}$-type). As will be seen from the practical

examples to be worked out later, there is complete consistency

between logical relations, and we really deal with logical relations

in quantifier theory, and not with merely 3-member set theory.

This will become clearer after examining particular examples

of the application of the theory.

Before going to the algorithms in QL-1A, we shall give

a brief introduction to the algebraic theory of the relation

between OL-1 and QL-1A formulations and then give the algorithms

for obtaining the output of unary relations and the SNS truth

values of binary relations in QL-1A.

(c) Theory of the QL-1A binary relation

The equations that we will deal with are (3a) and (3b)

for unary and binary relations in QL-1A,which correspond

respectively to (6b) and (6a) of QL-1 for binary reverse and

binary forward relations, as mentioned in Section 4(a). We shall

now give the essential features of this relationship which

enables us to carry over the algorithms developed above for

QL-1 almost completely into QL-1A.

Considering first, the QL-1A binary relation (3b), we are

given a relation, whose connective is associated with a quantifier

$(q_zx)$ and an SNS connective $\underline{z}(k, \ell)$, and we wish to find out the

truth value of this relation for the inputs $\underline{a}x$ and $\underline{b}x$. In QL-1,

we have already worked out the subroutine PBINPT to work out

the quantifier state $\underline{c}x$ of the resultant term when $\underline{a}x$ and $\underline{b}x$

are related by $\underline{z}(k, \ell)$ as in (6a). To obtain (3b) from this,

it is only necessary to find out the relative truth value QRELTV

of $\underline{c}x$ for $q_zx$. We shall illustrate this by the second example

of QL-1A binary relation in (5b). In this, the inputs are

$\underset{\sim}{a}x = (\exists x)(\underline{a}x) = (1\ 1\ 0)$, $\underset{\sim}{b}x = (\exists x)(\neg \underline{a}x) = (0\ 1\ 1)$, which

are put in the relation $\underset{\sim}{a}x\ \underline{A}(1,\ 2)\ \underset{\sim}{b}x$ , and can be shown,

using FFINPT, to yield the output $\underset{\sim}{c}x = (1\ 1\ 1) = (\bigwedge x)(\underline{c}x)$.

We wish to obtain the truth value of this for the quantifier

state $(\bigvee x)$ in Eq.(5b). For this, we make use of QL-2 algebra,

and calculate $\underline{t}(\exists\ |\ \bigvee) = (\ \langle 1\ 1\ 1\ |\ 1\ 0\ 0\rangle\ ,\ \langle 1\ 1\ 1\ |\ 0\ 1\ 1\rangle\ )$

$= (1\ 1) = D$, which is the answer given in (5b). Looking at the

problem logically, if $(\exists x)(\underline{a}x)$ is true, and $(\exists x)(\neg \underline{b}x)$

is true, then we cannot say for certain that $\underline{c}x = \underline{a}x\ \bigwedge\ \neg\underline{b}x$

is true for at least one/x, since $x_1$ for which $\underline{a}x$ is true, may,

or may not, correspond with $x_2$ for which $\underline{b}x$ is true. On the

other hand, the input conditions also permit that both $\underline{a}x$ and

$\neg\underline{b}x$ could be true for all $x$ , so that the quantifier state

$(\bigvee x)(\underline{c}x)$ is also possible. Hence the relation $(\bigvee x)(\underline{a}x\ \bigwedge\ \neg\underline{b}x)$

could be either true or false, i.e. it is indefinite for the

given inputs, so that its truth value is the SNS state $(1\ 1) = D$.

If, on the other hand, in (5b) the second input is

$(\exists x)(\underline{b}x) = (1\ 1\ 0)$, then using the same subroutine FBINPT,

we obtain $\underset{\sim}{c}x = (0\ 1\ 1) = (\exists x)(\neg \underline{c}x)$. The relative truth value

(QRELTV) of this for $(\bigvee x)(\underline{c}x)$ is $\underline{t}(\bigwedge\ |\ \bigvee) = (\langle 0\ 1\ 1\ |\ 1\ 0\ 0\rangle\ ,$

$\langle 0\ 1\ 1\ |\ 0\ 1\ 1\rangle) = (0\ 1\ ) = F$ , which indicates that the inputs

definitely do not satisfy the relation. Logically, also, if

for some $x$, $\underline{b}x$ is true, then the expression $(\bigvee x)(\underline{a}x\ \bigwedge\ \neg\underline{b}x)$

which require $\underline{b}x$ to be false for all $x$, cannot be true, so that

the truth value of the relation for the given inputs $(\exists x)(\underline{a}x)$

and $(\exists x)(\underline{b}x)$ is $F = (0\ 1)$. It is interesting that all these

considerations are automatically taken care of in EVMF.

Thus, the algorithm for calculating the SNS truth

value of a QL-1A relation becomes straightforward and may be

given as follows in subroutine (ix) PSTV.

(ix) QL-1A binary truth value — $(q_z x)(ax \geq (k, \ell) \; bx) = cx$ : PSTV

> PSTV(QVZ,QVA,SMX,SK,SL,QVB) = SVC
>
> QVC = PHINPT(QVA,SMX,SK,SL,QVB)
>
> SVC = QRELTV(QVC,QVZ)

This is probably the only algorithm that is needed for binary

relations in QL-1A algebra. We shall now consider unary

relations of the type (3a).

(d) Theory of the QL-1A unary relation

Before considering the QL-1A unary relation (3a), we shall

consider Eq.(6b) giving the QL-1 binary reverse relation.

The principle of the binary reverse has been discussed in

several earlier reports, and as applied to QL-1, it is considered

in some detail in MR-53 and 54. Essentially, the idea is that

the relation $ax \geq bx = cx$ is given, and instead of calculating

$cx$ as output for given $ax$ and $bx$, we calculate $bx$ as output

for given $cx$ and $ax$. The word "reverse" is applied because

the output of the relation is available, and in reverse,

it restricts the relationship between $ax$ and $bx$, so that,

if one is given, the other can be calculated. If this idea

is kept in mind, then/Tables 4(a) and 4(b) can be reversed to
(3x3)

obtain $b$, given $c$ and $a$, as in Table 5(a) and 5(b) (copied from

MR-56). Note that the notation for the reverse relation $a\ (c,Z) = b$ corresponds to the forward relation $a\ Z\ b = c$.

Table 5. 3x3 truth table for QL-1 binary reverse relations

(a) $a(c, A) = b$          (b) $a(c, Q) = b$

| a \ c | ∀ | Σ | φ |
|---|---|---|---|
| ∀ | ∀ | Σ | φ |
| Σ | φ | ∃ | Λ |
| φ | φ | φ | Δ |

| a \ c | ∀ | Σ | φ |
|---|---|---|---|
| ∀ | Δ | φ | φ |
| Σ | ∃ | Λ | φ |
| φ | ∀ | Σ | φ |

These tables have been derived in MR-53 and 54. We /shall indicate here

how they /are algorithmised for application. As given in page 15,

Lecture-4 of MR-56, a very neat algorithm can be given for the

quantifier state of the output as given in Eqs.(11a,b) and (12a,b)/below.

Denoting the quantifier states of $a$ and $c$ by $q(i)$ and

$q(k)$, then, for $i,k = 1, 3,5$ for basic states, the quantifier state

$q(j)$ of the output $b$, is given by (11) and (12) respectively,

for the connectives $A(1, 1)$ and $Q(1, 1)$.

For $\underline{A}$

$$\text{If} \quad k < i , \quad q(j) = q(8) \tag{11a}$$

$$\text{If} \quad k \geq i, \quad q(j) = \bigoplus_{\ell=0}^{L} q(k - i + 1 + 2), \quad \text{where} \quad L = (i+1)/2 \tag{11b}$$

For $\underline{O}$

$$\text{If} \quad k > i , \quad q(j) = q(8) \tag{12a}$$

$$\text{If} \quad k \leq i, \quad q(j) = \bigoplus_{\ell=0}^{L} q(k + 2), \quad \text{where} \quad L = (7 - i)/2 \tag{12b}$$

To obtain $\underline{a}(\underline{c}, \underline{I}) = \underline{b}$, we make use of the equivalence

$(\underline{a}\, \underline{I}\, \underline{b} = \underline{c}) \Longleftrightarrow (\underline{a}^n\, \underline{O}\, \underline{b} = \underline{c})$, so that, in reverse, we can express

the operation of $\underline{I}$ by $\underline{O}$, as in (13).

$$(\underline{a}(\underline{c}, \underline{I}) = \underline{b}) \Longleftrightarrow (\underline{a}^n(\underline{c}, \underline{O}) = \underline{b}) \tag{13}$$

Also, more generally, for $\underline{A}(k, \ell)$ and $\underline{O}(k, \ell)$, we use (9)

for the forward relation, which yields in reverse, the equations

(14a,b,c).

$$(\underline{a}(\underline{c}, \underline{Z}(k, \ell)) = \underline{b}) \Longleftrightarrow (\underline{a}'(\underline{c}, \underline{Z}(1,1)) = \underline{b}') \tag{14a}$$

where

$$\underline{a}' = \underline{a}, \ \underline{a}^n \quad \text{according as} \quad k = 1, 2 \ ; \tag{14b}$$

$$\underline{b}' = \underline{b}, \ \underline{b}^n \quad \text{according as} \quad \ell = 1, 2 \tag{14c}$$

No simple expression can be given for $\underline{E}$, occurring in a QL-1

reverse relation, in terms of the corresponding relations for

$\underline{A}$ and $\underline{O}$, in this formalism. However, by directly inverting

the 3x3 truth table of the forward relation $a \underline{E} b = c$, it

can be shown that the reverse relation $a(c, \underline{E}) = b$ has exactly

the same 3x3 table in the format of Tables 5(a) and 5(b), as the

corresponding table for the forward relation in the format of

Tables 4(a) and 4(b). (See MR-53 and MR-56 for these tables —

the corresponding 8x8 tables are printed out in Section 4(e)).


We shall now show that the binary reverse relation in

QL-1 is equivalent to the unary relation in QL-1A, by taking

an example, as was done for the binary forward relation. The

inter-relationship between (3a) and (6b) $\overset{is}{/}$ that, if $cx$ of (6b)

is put equal to $(q_z x)$ of (3a), then the two equations are

logically equivalent to one another. Considering the example

(4a), we are given that there exists at least one $x''$ for which

$\underline{a}x$ implies $\underline{b}x$, and that $\underline{a}x$ is true for all $x''$. Then the

conclusion that "there exists an $x''$ for which $\underline{b}x$ is true is

obvious. In the QL-1 formalism, it takes the following form.

In the binary forward form, given that $ax$ is true for all $x$,

and $bx$ is true for at least one $x$, the relation $\underline{a}x \Longrightarrow \underline{b}x = \underline{c}x$

is true for at least one $x$. Therefore, given $\underline{a}x$ is true for all $x$,

and that the relation is true for at least one  x,  what is

the quantifier state of  $\underline{\underline{b}}x$  that will satisfy this condition?

Obviously, it is  $(\exists x)(\underline{\underline{b}}x)$.

Hence, Eqs.(11a,b) and (12a,b) can be taken over bodily

for Eq.(3a) with the equivalences  (15), for all  $\underline{\underline{Z}} = \underline{\underline{A}}, \underline{\underline{O}}, \underline{\underline{E}}, \underline{\underline{I}}$ :

$$(\underline{a} = q(i), \underline{b} = q(j), \underline{c} = q(k)) \text{ of } QL\text{-}1$$

$$\Longleftrightarrow (\underline{a} = q(i), \underline{b} = q(j), \underline{q}_{\underline{Z}} = q(k)) \text{ of } QL\text{-}1A \qquad (15)$$

for $\underline{A}$ and $\underline{\underline{O}}$
The    two subroutines/are named "unary and", "unary or" of

PLOGIC, and refer to QL-1A algebra with the inputs put in the

sequence  $\underline{q}_{\underline{Z}}$, $\underline{a}$ , and the output being labelled  $\underline{b}$ .   The

following MATLOG subroutines will become obvious from the above

considerations.


(x) <u>Unary and   for QL-1A</u>  —   $\underline{a}x$,  $(\underline{q}_{\underline{Z}}x)(\underline{a}x \underline{\underline{A}} \underline{\underline{b}}x) \longmapsto \underline{\underline{b}}x$ : PUNAND

        PUNAND(QVZ,QVA) = QVB

        QELEM(QVA) = QBAS1(IB1)

        QELEM(QVZ) = QBAS2(IB2)

    IF   QVZ . OR . QVA = Q8 ,  QOUT = Q8,    GO TO  1

    ELSE        QOUT = Q8

        DO  I1 = 1,  IB2,  J1 =   1,  IB1

Denote  QBAS2(I1)  as  QK

Denote  QBAS1(J1)  as  QI

If  $K < I$,  QJ = Q8

else  QJ = Q(K - I + 1) $\oplus$ Q(K - I + 3) $\oplus$ Q(K - I + 5)

where number of terms is $(I + 1)/2$

10       QOUT = QUNION(QOUT,QJ)

1        QVB = QOUT


(xi) <u>Unary or for QL-1A</u> — $\underset{\sim}{a}x$, $(\underset{\sim Z}{q}x)(\underline{a}x \ \underline{0} \ \underline{b}x) \mapsto \underline{b}x$ : PUNOR

PUNOR(QVZ,QVA) = QVB

Same as for PUNAND except that the conditions are replaced by

If  $K > I$,  QJ = Q8

else  QJ = Q(K - I + 1) $\oplus$ Q(K - I + 3) $\oplus$ Q(K - I + 5)

where number of terms is $(7 - I)/2$


We can generalize these two unary relations to a general
unary relation of the type in (3a) with $\underline{Z} = \underline{Z}(k, \ell)$,
$Z = A$, $O$, $E$ or $I$. In an obvious notation, this takes the form
of the function PUNPDT below.


(xii) General unary product for QL-1A — $\underset{\sim}{a}x$, $(\underset{\sim Z}{q}x)(\underline{a}x \ \underline{Z}(k, \ell) \ \underline{b}x)$
$\mapsto \underset{\sim}{b}x$ : PUNPDT

PUNPDT(QVZ,QVA,SMX,SK,SL) = QVB

SK = S1 $\longrightarrow$ QOK = QEQU

SK = S2 $\longrightarrow$ QOK = QNOT

SL = S1 $\longrightarrow$ QOL = QEQU

SL = S2 $\longrightarrow$ QOL = QNOT

If SMX = SMA ,

     QVBP = PUNAND(QVZ,QOK(QVA))

     QVB = QOL(QVBP)

If SMX = SMO ,

     QVB = QOL(PUNOR(QVZ,QOK(QVA)))

If SMX = SMI ,

     QVB = QOL(PUNOR(QVZ, QOK(PNOT(QVA))))

If SMX = SME ,

     QVB = PUNION(PVIDYA(QOK(QVA),QOL(QVZ)),

        PVIDYA(PNOT(QOK(QVA)),PNOT(QOL(QVZ))))

This completes all the subroutines that are needed for PLOGIC in its application to QL-1 type binary forward and binary reverse relations as well as QL-1A type unary and binary relations. No separate program is written for QL-1 binary reverse relations as they are equivalent to QL-1A unary relations and these subprograms for PUNAND and PUNOR can be used for this purpose with QVZ being replaced by QVC for the QL-1 binary relation.

As will be seen from the outputs of these subroutines which are obtained as solutions of some of the problems given below, the MATLOG algorithms, which are based on a semi-empirical algebraic notation employed for PLOGIC, agree completely with the truth tables obtained from an intuitive/logical approach in the earlier reports MR- 53, 54 and 56.

(e) Illustrations and problems in PLOGIC

The examples given below are of three types.

(a) Outputs of the subroutines PUNION, PVIDYA, PUNAND

PUNOR, PRINT , PUNPDT and PSTV. Not all of them are printed

out, but some examples are shown to illustrate the nature of

the output of these functions.

(b) Some formal relations between different logical

connectives are tested out in 8x8 tables by the application

of MATLOG. These include the description of equivalence

as the union of suitable conjunctions and as the vidya of

suitable implications, the De Morgan relation between a

conjunction and $/$ corresponding disjunction in QL-1, and
(its)

the equivalence (13) in QL-1A :

$$(q_z^\ell x)(\underset{\approx}{a}x \; \underline{z}^c \; \underline{b}x) \iff \neg (q_z x)(\underline{a}x \; \underline{z} \; \underline{b}x)$$
                                                                (13)

(c) Illustrative examples of some simple problems in

predicate calculus solved by PLOGIC subroutine of MATLOG.

Problems 9A and B

A. Obtain the 8x8 truth table for binary "or" and "and"

in QL-1, employing the MATLOG subroutines PUNION and PVIDYA.

B.   Similarly obtain the 8x8 truth table for the

connectives "imply" and "equivalent" using PBINPT.


Comments on the output of Problems 9A and 9B

   The four tables corresponding to 9A and 9B are given
                           on pages 45 and 46
below in Tables 5(a-d)/.  It can be verified that they agree

completely with Tables 7(a-d), pages 120, 121 of MR-53 which

were obtained by intuitively working out the logic of each

QL-1 binary relation from ab initio considerations based on

the definition of the QL-1 relation in terms of the input
                   of
truth values/$\underline{a}$x, $\underline{b}$x and $\underline{c}$x.


Problems 10 A and 10B

   Both the functions PUNOR and PUNAND as well as the general

subroutine PUNPDT can be checked by obtaining a set of tables

similar to Problem 9 from the MATLOG programs.  This is done

via problems 10A and B.

   10A : Print out the 8x8 truth tables for QL-1 "or"(PUNOR)

and QL-1 "and"(PUNAND) for the logical relation $(q_Z x)(\underline{a}x \ \underline{Z} = \underline{b}x)$

for $\underline{Z} = \underline{A}, \underline{O}$ .

**Table 5. Outputs of the subroutines PUNOR, PUNAND and of PUNPDT for E(1, 1) and I(1, 1)**

(a)  OUTPUT OF PUNOR

| OUA | AII | EXS | NAI | NEX | SOM | INU | AON | IMP |
|---|---|---|---|---|---|---|---|---|
| AII | AII | AII | ALI | AII | AII | AII | AII | IMP |
| EXS | AII | EXS | EXS | EXS | EXS | EXS | EXS | IMP |
| NAI | AII | EXS | INU | NAI | EXS | INU | INU | LMP |
| NEX | AII | EXS | NAI | NEX | SOM | INU | AON | LMP |
| SOM | AII | EXS | EXS | SOM | EXS | EXS | EXS | IMP |
| INU | AII | EXS | INU | INU | EXS | INU | INU | IMP |
| AON | AII | EXS | INU | AON | EXS | INU | AON | IMP |
| IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

(b)  OUTPUT OF PUNAND

| OUA | AII | EXS | NAI | NE | SOM | INU | AON | IMP |
|---|---|---|---|---|---|---|---|---|
| AII | AII | EXS | NAI | NEX | SOM | INU | AON | IMP |
| EXS | EXS | INU | NAL | NEX | NAL | INU | INU | IMP |
| NAI | NAI | NAI | NAI | NEX | NAI | NAI | NAI | IMP |
| NE | NE | NI | NEX | NE X | NEX | NEX | NE X | IMP |
| SOM | SOM | NAI | NAI | NEX | NAI | NAI | NAI | IMP |
| INU | INU | INU | NAL | NEX | NAI | INU | INU | IMP |
| AON | AON | INU | NAL | NE X | NAI | INU | AON | IMP |
| IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

**(c)** OUTPUT FOR PRINT OF I(1-1)

| OVA | QVB | ALL | EXS | NAL | NIX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|---|
| ALL | | ALL | EXS | NAL | NIX | SOM | IND | AON | IMP |
| EXS | | EXS | IND | TND | NAL | IND | | IND | IMP |
| NAL | | NAL | IND | IND | EXS | IND | IND | IND | IMP |
| NIX | | NIX | NAL | EXS | ALL | SOM | IND | AON | IMP |
| SOM | | SOM | IND | IND | SOM | IND | IND | IND | IMP |
| IND | | IND | IND | IND | IND | IND | IND | IND | IMP |
| AON | | AON | IND | IND | AON | IND | IND | AON | IMP |
| IMP | | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

**(d)** OUTPUT FOR PRINT OF I(1,1)

| OVA | QVB | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|---|
| ALL | | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
| EXS | | ALL | EXS | IND | NAL | EXS | IND | IND | IMP |
| NAL | | ALL | EXS | EXS | EXS | EXS | EXS | EXS | IMP |
| NIX | | ALL | ALL | ALL | ALL | ALL | ALL | ALL | IMP |
| SOM | | ALL | EXS | EXS | SOM | EXS | EXS | EXS | IMP |
| IND | | ALL | EXS | IND | IND | EXS | IND | IND | IMP |
| AON | | ALL | EXS | IND | AON | EXS | IND | AON | IMP |
| IMP | | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

10B : Print the truth tables for the above predicate

calculus relation for the connectives $E(1, 1)$, $I(1, 1)$, $A(2, 2)$,

$O(2, 2)$ using the general subroutine PUNPUT for these connectives.

Comments on the output of Problems 10A and 10B

The outputs are shown in Tables 6(a-d) in pages 48 and 49.

It will be noticed that the tables are not symmetric in the two

inputs, namely QVZ and QVA, since the former is the quantifier $q_z$

associated with the connective $Z$, while the latter is the input

$ax$. It can be checked that the first two tables 6(a,b) for the

Problem 10A are identical with Tables 6(a,b) in page 113 of MR-53,

obtained for the same problems by intuitive ab initio calculations,

using the logic of the relation between $ax$ and $bx$ for individual

members of the set that are quantified. The data in Tables 6(c,d)

of Problem 10B can be seen to agree with Tables 6(c,d) in page

113a of MR-53, while those in (e, f) indicate how other connectives

can be implemented.

The fact that simple algorithms could be formulated for

working out all the possibilities indicate the cogency and

consistency of the algebraic formalism which has been adopted

for predicate logic via the BA-3 algebra of quantifiers. The

tests given below are designed to check the interconnections

between these formulae.

Table 6. 8x8 truth tables for unary relations in QL-1 with the connectives A(1, 1), Q(1, 1), E(1, 1) and I(1, 1)

(a)  OUTPUT OF RUNOR

| QVZ \ QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|
| ALL | IND | IND | EXS | ALL | EXS | IND | IND | IMP |
| EXS | IND | IND | IND | EXS | IND | IND | IND | IMP |
| NAL | IMP | NAL | NAL | NAL | NAL | NAL | NAL | IMP |
| NEX | IMP | IMP | NEX | NEX | IMP | NEX | NEX | IMP |
| SOM | IMP | NAL | NAL | SOM | NAL | NAL | SOM | IMP |
| IND | IND | IND | IND | IND | IND | IND | IND | IMP |
| AON | IND | IND | IND | AON | EXS | IND | IND | IMP |
| IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

(b)  OUTPUT OF RUNAND

| QVZ \ QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|
| ALL | ALL | ALL | IMP | IMP | IMP | ALL | ALL | IMP |
| EXS | EXS | EXS | EXS | IMP | EXS | EXS | EXS | IMP |
| NAL | NAL | IND | IND | IND | IND | IND | IND | IMP |
| NEX | NEX | NAL | IND | IND | NAL | IND | IND | IMP |
| SOM | SOM | EXS | EXS | IMP | EXS | EXS | SOM | IMP |
| IND | IND | IND | IND | IND | IND | IND | IND | IMP |
| AON | AON | IND | IND | IND | NAL | IND | IND | IMP |
| IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

(c)  OUTPUT FOR PUNPUT OF E(1,1)

| QVZ \ QVA | QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|---|
| ALL | | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
| EXS | | EXS | IND | IND | NAL | IND | IND | IND | IMP |
| NAL | | NAL | IND | IND | EXS | IND | IND | IND | IMP |
| NEX | | NEX | NAL | EXS | ALL | SOM | IND | AON | IMP |
| SOM | | SOM | IND | IND | SOM | IND | IND | IND | IMP |
| IND | | IND | IND | IND | IND | IND | IND | IND | IMP |
| AON | | AON | IND | IND | AON | IND | IND | AON | IMP |
| IMP | | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

(d)  OUTPUT FOR PUNPUT OF I(1,1)

| QVZ \ QVA | QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|---|
| ALL | | ALL | EXS | IND | IND | EXS | IND | IND | IMP |
| EXS | | EXS | IND | IND | IND | IND | IND | IND | IMP |
| NAL | | NAL | NAL | NAL | IMP | NAL | NAL | NAL | IMP |
| NEX | | NEX | NEX | IMP | IMP | IMP | NEX | NEX | IMP |
| SOM | | SOM | NAL | NAL | IMP | NAL | NAL | SOM | IMP |
| IND | | IND | IND | IND | IND | IND | IND | IND | IMP |
| AON | | AON | IND | IND | IND | EXS | IND | IND | IMP |
| IMP | | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

**(e)**    OUTPUT FOR FUNPDT OF A(2,2)

| QVZ \ QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|
| ALL | ALL | ALL | IMP | IMP | IMP | ALL | ALL | IMP |
| EXS | EXS | EXS | EXS | IMP | EXS | EXS | EXS | IMP |
| NAL | NAL | IND | IND | IND | IND | IND | IND | IMP |
| NEX | NEX | NAL | IND | IND | NAL | IND | IND | IMP |
| SOM | SOM | EXS | EXS | IMP | EXS | EXS | SOM | IMP |
| IND | IND | IND | IND | IND | IND | IND | IND | IMP |
| AON | AON | IND | IND | IND | NAL | IND | IND | IMP |
| IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

**(f)**    OUTPUT FOR FUNPDT OF O(2,2)

| QVZ \ QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|
| ALL | IND | IND | EXS | ALL | EXS | IND | IND | IMP |
| EXS | IND | IND | IND | EXS | IND | IND | IND | IMP |
| NAL | IMP | NAL | NAL | NAL | NAL | NAL | NAL | IMP |
| NEX | IMP | IMP | NEX | NEX | IMP | NEX | NEX | IMP |
| SOM | IMP | NAL | NAL | SOM | NAL | NAL | SOM | IMP |
| IND | IND | IND | IND | IND | IND | IND | IND | IMP |
| AON | IND | IND | IND | AON | EXS | IND | IND | IMP |
| IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

Problem 11A, B

A : Verify that the two ways of writing the equivalence

relation, as disjunction of $\underline{A}(1,1)$ and $\underline{A}(2,2)$, and the conjunction

of $\underline{I}(1,1)$ and $\underline{I}(2,2)$, which is valid in SNS, can be carried over

to QL-1. This can done by writing the equivalence relation

$\underset{\sim}{a}x\ \underline{E}(1,1)\ \underset{\sim}{b}x$ as

$$(\underset{\sim}{a}x\ \ \underline{A}(1,\ 1)\ \underset{\sim}{b}x)\ \ \underline{O}(1,\ 1)\ \ (\underset{\sim}{a}x\ \ \underline{A}(2,\ 2)\ \ \underset{\sim}{b}x) \qquad (16a)$$

*and*

$$(\underset{\sim}{a}x\ \ \underline{I}(1,\ 1)\ \underset{\sim}{b}x)\ \ \underline{A}(1,\ 1)\ \ (\underset{\sim}{a}x\ \ \underline{I}(2,\ 2)\ \ \underset{\sim}{b}x) \qquad (16b)$$

B : In a similar manner, check the following De Morgan

relation given in (17), analogous to the De Morgan relation in

SNS, between the connectives $\underline{O}(1,\ 2)$ and $\underline{A}(1,\ 2)$ :

$$(\underset{\sim}{a}x\ \ \underline{O}(1,\ 2)\ \ \underset{\sim}{b}x)\ =\ \neg(\ \neg\underset{\sim}{a}x\ \ \underline{A}(1,\ 2)\ \neg\ \underset{\sim}{b}x) \qquad (17)$$

Solution of the Problems 11A and B

A: The MATLOG statements to be verified are as follows:

        QVC1 = PUNION(PVIDYA(QVA,QVB),

               PVIDYA(QNOT(QVA),QNOT(QVB)))

        QVC2 = PVIDYA(PIMPLY(QVA,QVB), PIMPLY(QNOT(QVA),QNOT(QVB)))

where         PIMPLY(QVA,QVB) = PUNION(QNOT(QVA),QVB )

The 8x8 tables for QVC1 and QVC2 obtained as computer outputs

are given below in Table 7, page 51.

## Table 7. Output for Problem 11A

QVC1

| QVB QVA | ALL | EXS | NAI | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|
| ALL | ALL | EXS | NAI | NEX | SOM | IND | AON | IMP |
| EXS | EXS | IND | IND | NAL | IND | IND | IND | IMP |
| NAI | NAL | IND | IND | EXS | IND | IND | IND | IMP |
| NEX | NEX | NAL | EXS | ALL | SOM | IND | AON | IMP |
| SOM | SOM | IND | IND | SOM | IND | IND | IND | IMP |
| IND | IND | IND | IND | IND | IND | IND | IND | IMP |
| AON | AON | IND | IND | AON | IND | IND | AON | IMP |
| IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

QVC2

| QVB QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|
| ALL | ALL | EXS | NAI | NEX | SOM | IND | AON | IMP |
| EXS | EXS | IND | IND | NAL | IND | IND | IND | IMP |
| NAL | NAL | IND | IND | EXS | IND | IND | IND | IMP |
| NEX | NEX | NAL | EXS | ALL | SOM | IND | AON | IMP |
| SOM | SOM | IND | IND | SOM | IND | IND | IND | IMP |
| IND | IND | IND | IND | IND | IND | IND | IND | IMP |
| AON | AON | IND | IND | AON | IND | IND | AON | IMP |
| IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

It will be seen that the two tables completely agree

and also that they agree with Table 7(d) of MR-53, where the

8x8 table, obtained from apriori logical considerations, is given.


11B :  The MATLOG statements for verifying the equivalence

of the r.h.s and l.h.s of (17) are


PBINPT(QVA,SMO,S1,S2,QVB)  =  QVC1

QNOT(PBINPT(QNOT(QVA),SMA,S1,S2, QNOT(QVB)))  =  QVC2

The check that QVC1 = QVC2 for all possible quantifier states,

Q1 to Q8, for the inputs QVA and QVB, may be seen from Table 8

below. This indicates the essential facility of MATLOG

algorithms to take care any general logical statement in

QL-1. It should be noted that "negation" is QNOT, corresponding

to the BA-3 operator $\underset{\sim}{N}$, for QL-1A logic, and this is uniformly

applicable to all formulae. The reason for this is because

the negation is for the predicate and not for the quantifier

as in QL-2 logic.

## Table 8.   Output for Problem 11B

QVC1

|  | QVB<br>QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|---|
| ALL |  | ALI | AI L | AEL | AI L | ALL | ALI | AI I | IMP |
| EXS |  | AI L | EXS | EXS | EXS | EXS | EXS | EXS | IMP |
| NAI |  | ALI | EXS | IND | NAL | EXS | IND | IND | IMP |
| NEX |  | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
| SOM |  | ALL | EXS | EXS | SOM | EXS | EXS | EXS | IMP |
| IND |  | ALL | EXS | IND | IND | EXS | IND | IND | IMP |
| AON |  | ALL | EXS | IND | AON | EXS | IND | AON | IMP |
| IMP |  | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

QVC2

|  | QVB<br>QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|---|---|---|---|---|---|---|---|---|---|
| ALL |  | AI I | ALI | ALL | ALL | ALL | ALL | ALL | IMP |
| EXS |  | AI I | EXS | EXS | EXS | EXS | EXS | EXS | IMP |
| NAI |  | ALI | EXS | IND | NAI | EXS | IND | IND | IMP |
| NEX |  | ALI | EXS | NAI | NEX | SOM | IND | AON | IMP |
| SOM |  | AI L | EXS | EXS | SOM | EXS | EXS | EXS | IMP |
| IND |  | ALI | EXS | IND | IND | EXS | IND | IND | IMP |
| AON |  | ALI | EXS | IND | AON | EXS | IND | AON | IMP |
| IMP |  | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |

## Problem 12A, B

These two problems require a new function PSMXCP, which is described in Section 4(f) in pages 56 and 57 which should be read before going further. These two problems will be based on Eq.(13) on page 43 and will check the validity of the equivalence given therein for chosen inputs. Problem 12A will check PSINPT and PSTV, while 12B will check PUMPDT.

12A : Check the equality of the resultant truth values $c_1$ and $c_2$ of the two equivalent statements (18a,b) for the $8 \times 8$ possible pairs $ax$, $bx$.

$$\neg(\exists x)(ax \equiv bx) \iff (\forall x)(ax \equiv^c bx) \qquad (18a,b)$$

## Solution to Problem 12A

The MATLOG statements for this can be given as follows:

PSTV(QCOMP(Q6),QVA,SMI,S1,S1,QVB)  =  SVC1

PSTV(Q1,QVA,PSMXCP(SMI,S1,S1),QVB) = SVC2

Input  QVA,QVB = Q1 to Q8 .

The equality of SVC1 and SVC2 may be seen from Tables 9(a, b) giving the outputs for SVC1 and SVC2 .

## Table 9 :  8x8 Tables of SVC1 and SVC2 for PSTV

SVC1

| QVA \ QVB | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|-----------|-----|-----|-----|-----|-----|-----|-----|-----|
| ALL | F | F | D | T | F | D | D | X |
| EXS | F | F | D | D | F | D | D | X |
| NAL | F | F | I | F | F | F | F | X |
| NEX | F | F | F | F | F | F | F | X |
| SOM | F | F | I | F | F | F | F | X |
| IND | F | F | U | D | F | D | D | X |
| AON | F | F | D | D | F | D | D | X |
| IMP | X | X | X | X | X | X | X | X |

| QVA \ QVB | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|-----------|-----|-----|-----|-----|-----|-----|-----|-----|
| ALL | F | I | D | T | F | D | D | X |
| EXS | F | F | D | D | F | U | D | X |
| NAL | F | F | F | F | F | F | F | X |
| NEX | F | F | F | F | F | I | F | X |
| SOM | F | F | F | F | F | F | F | X |
| IND | F | F | U | D | F | U | D | X |
| AON | F | F | D | D | F | D | D | X |
| IMP | X | X | X | X | X | X | X | X |

Problem 12B

Check the equality of the quantifier states of the two outputs $\underline{b}_1 x$ and $\underline{b}_2 x$ given by (19a,b) for all the eight possible quantifier states of $(q_z x)$ and $\underline{g}x$ .

$$\underline{g}x, \quad \overline{\phantom{-}}1(\underline{q}_z x)(\underline{g}x \,\underline{I}\, \underline{b}x) \mid \rightarrow \underline{b}_1 x \qquad (19a)$$

$$\underline{g}x, \quad (\underline{q}_z x)(\underline{g}x \,\underline{I}^c\, \underline{b}x) \mid \rightarrow \underline{b}_2 x \qquad (19b)$$

Solution to Problem 12B

The Fortran statements for the outputs QVB1 and QVB2, in terms of the inputs QVZ and QVA, are as follows:

FUNPDT(QCOMP(QVZ),QVA,SMI,S1,S1) = QVB1

FUNPDT(QELL(QVZ),QVA,PSMXCP(SMT,S1,S1)) = QVB2

For checking the equality of QVB1 and QVB2, the 8x8 tables are printed side by side in Table 10(a, b). It will be seen that the two tables agree, showing that the algorithms including the new function PSMXCP are well formulated.

.55a.

## Table 10 : 8x8 tables for QVB1 and QVB2 of Problem 12B

QVB1

| QVZ \ QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|-----------|-----|-----|-----|-----|-----|-----|-----|-----|
| ALL | NAL | NAL | NAL | IMP | NAL | NAL | NAL | IMP |
| EXS | NEX | NEX | IMP | IMP | IMP | NEX | NEX | IMP |
| NAL | ALL | EXS | IND | IND | EXS | IND | IND | IMP |
| NEX | EXS | IND | IND | IND | IND | IND | IND | IMP |
| SOM | AON | IND | IND | IND | EXS | IND | IND | IMP |
| IND | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |
| AON | SOM | NAL | NAL | IMP | NAL | NAL | SOM | IMP |
| IMP | IND | IND | IND | IND | IND | IND | IND | IMP |

QVB2

| QVZ \ QVA | ALL | EXS | NAL | NEX | SOM | IND | AON | IMP |
|-----------|-----|-----|-----|-----|-----|-----|-----|-----|
| ALL | NAL | NAL | NAL | IMP | NAL | NAL | NAL | IMP |
| EXS | NEX | NEX | IMP | IMP | IMP | NEX | NEX | IMP |
| NAL | ALL | EXS | IND | IND | EXS | IND | IND | IMP |
| NEX | EXS | IND | IND | IND | IND | IND | IND | IMP |
| SOM | AON | IND | IND | IND | EXS | IND | IND | IMP |
| IND | IMP | IMP | IMP | IMP | IMP | IMP | IMP | IMP |
| AON | SOM | NAL | NAL | IMP | NAL | NAL | SOM | IMP |
| IMP | IND | IND | IND | IND | IND | IND | IND | IMP |

(f) Complementation of the name of an SNS logical connective

The standard definition of the SNS truth value of a
relation $\underline{a} \; \underline{Z}(k, \ell) \; \underline{b} = \underline{c}$ is based on the two equations

$\langle \underline{a} \mid \underline{Z}(k,\ell) \mid \underline{b} \rangle = c_1$, $\langle \underline{a} \mid \underline{Z}^c(k,\ell) \mid \underline{b} \rangle = c_2$ , leading to

$\underline{c} = (c_1, c_2)$. For doing this, the procedure adopted so far

is to convert the 2x2 matrix $\underline{Z}(k, \ell)$ into its complement by

taking the Boolean complement of each of the components $Z_{\lambda\mu}$

of the matrix. In other words, the definition of the complement

of a matrix connective is the one associated with the Boolean

complement of its matrix. This, and the associated formulae
                           in QLOGIC
for 3x3 matrices/ have been the main basis of all the formulae

so far. However, in PLOGIC and also in PQLOGIC, we find
                                the
the need to associate/name $\underline{Z}(k, \ell)$ with the designation of

$\underline{Z}$ as one of $\underline{A}, \underline{O}, \underline{E}, \underline{I}$ and the indices $k$ or $\ell$ standing for

S1 or S2, and these are to be transformed into the name and

indices of the complement of the connective. Thus,

$\underline{A}^c(1, 1) = \underline{O}(2, 2)$; $\underline{I}^c(1, 1) = \underline{O}^c(2, 1) = \underline{A}(1, 2)$. Denoting

the original connective by $\underline{P}$, and its complement $\underline{P}^c$ by $\underline{Q}$, we wish

to associate the name P, and the indices $k_P$ and $\ell_P$ , with

the name Q and the indices $k_Q$ and $^* \ell_Q$ by a suitable function.

This is readily possible, in terms of our MATLOG notation,

by defining a new function called "name-complement" of an

SNS matrix in PLOGIC, namely PSMXCP, having the following

structure.

(xiii) Name-complement of an SNS connective in PLOGIC — $P^c = Q$

Denoting $P$ by SMXP,SKP,SLP and $Q$ by SMXQ,SKQ,SLQ, we have

PSMXCP(SMXP,SKP,SLP) = SMXQ,SKQ,SLQ

For  SMXP = SMA, SMXQ = SMO, SKQ = SCOMP(SKP), SLQ = SCOMP(SLP)

SMXP = SMO, SMXQ = SMA, SKQ = SCOMP(SKP), SLQ = SCOMP(SLP)

SMXP = SMI, SMXQ = SMA, SKQ = SKP, SLQ = SCOMP(SLP)

SMXP = SME, SMXQ = SME, SKQ = SKP, SLQ = SCOMP(SLP)

As mentioned in connection with Problems 12A and 12B, this / subroutine
has interesting possibilities for solving problems in PLOGIC

and PQLOG, using the algebra of QL-1A and QL-1B. It can be side

stepped and incorporated in the program for various functions as

given below. However, the function PSMXCP is very valuable for

theoretical analysis of the inter-relations between QL-2, QL-1

QL-1A and QL-1B, and the program is likely to    find application

in checking such formulae. In fact, the discussion given in the

next Section 4(g) arose as a result of testing such ideas in

connection with the algorithmization of QL-1B.

(g) Construction of the matrix of a "logical" connective in QL-2

As mentioned above, when the above ideas were fed in for

PQLOGIC (see Section 5), we observed that there is a need for a

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function for obtaining the matrix $|\underset{\sim}{Z}(k, \ell)|$ in GL-2 for a

relation of the type $(\forall x)(\underline{a}x)\ \underline{Q}(2, 1)\ (\exists y)(\underline{b}y)$, from the

three given data, namely $(\forall x), \underline{Q}(2, 1),(\exists y)$, which has not been

formulated as an algorithm in MATLOG, although it is implied

in (17) of Section 1(c) dealing with logical operators in GL-2.

In the above example, the SMS connective $\underline{Q}(2, 1)$, standing

for $\neg \underline{a} \lor \underline{b}$, is converted, by addition of the quantifiers

$(\forall x)$ and $(\exists y)$ to $\underline{a}$ and $\underline{b}$, into the form $\neg(\forall x)(\underline{a}x) \lor (\exists y)(\underline{b}y)$

corresponding to the GL-2 connective $\underset{\sim}{Q}(\underset{\sim}{2}, \underset{\sim}{6})$. The feature that

requires attention, from the point of view of the algorithm,

is that the operation of negation in SMS (SCOMP) in $\neg \underline{a} \lor \underline{b}$

is converted into the operation of "complementation" in GL-2

(GEM) in $\neg (\forall x)(\underline{a}x)$, and the latter is attached to the

quantifier $\forall$ which is associated with $\underline{a}$, and not to the

predicate $\underline{a}$ . This may look obvious, but it has to be properly

introduced into the set of subroutines, for it finds application

in various problems. In fact, in an analogous situation in GL-1B,

$(\forall x)(\exists y)(\neg \underline{a}x \lor \underline{b}y)$ has the equivalent form $(\forall x)(\neg \underline{a}x) \lor$

$(\exists y)(\underline{b}y)$ in GL-2, which corresponds to the connective $\underline{Q}(\underset{\sim}{5}, \underset{\sim}{6})$

and the SMS "negation" SCOMP in $\neg \underline{a}x$ is converted into the

"negation" of the quantified term in QL-2,corresponding to the

operation CHBN. The full algebraic theory of these is

derivable from the material presented in MR-53 and MR-54, but

it is found that they take a much simpler algorithmic form

which has been discovered in the attempt at computerizing them.

A brief summary of the theory of QL-1B, and the algorithm for

the implementation of elementary statements in QL-1B $\underline{via}$ their

QL-2 equivalents, are presented in Section 5 below. The

treatment here closely follows the notation adopted in MR-56,

Lecture-4, for QL-1A and which we have adopted for PLOGIC. Since

the form of the relation in QL-1B is similar to QL-1, while

it is implemented by converting it into QL-2, these types of

statements are put under the title "PQLOGIC" and discussed

in the next Sections 5(a,b). The procedure for obtaining

the equivalent matrix connective is closely similar to that

in PLOGIC and the corresponding algorithm is given

in Section 5(b).

In view of these, we give below the amended form of

Q(xvii) as Q(xviia) and Q(xvii b) , the first to obtain the

name of the matrix QMZ,QK,QL in terms of QVK,SMX,SK,SL,QVL,

and the second for obtaining the matrix elements of the 3x3

matrix corresponding to $QMZ,QK,QL$. We give, in Section 5(c),

the subroutine $PQ(i)$, corresponding to $Q(xviia)$, as applied to

FQLOGIC.

(xviia) <u>Matrix connective in QLOGIC defined via an SMS logical</u>
<u>connective:</u> $\underline{q}(k_Z)\ \underline{z}(k,\ell)\ \underline{q}(\ell_Z) = \underline{z}(k,\ell)$

$$QMCON(QVK,SMX,SK,SL,QVL) = QMZ,QK,QL$$

The notation may be illustrated by the example given above,

for which we have

$$QMCON(Q1,\ SMI,S2,S1,Q6) = QMI,Q2,Q6$$

The statements required for this subroutines are

$QMZ = QMA,QMO,QMI,QME$  according as  $SMX = SMA,SMO,SMI,SME$

$QK = QEQU(QVK),QNOT(QVK)$ according as  $SK = S1,S2$

$QL = QEQU(QVL),QNOT(QVL)$ according as  $SL = S1,S2$

(xviib) <u>Matrix elements of a QL-2 logical connective:</u>
$$\left|\underline{z}(k,\ell)\right| = z_{ij}\ ,\quad i = 1,2,3,\quad j = 1,2,3$$

$$QMATEL(QMZ,QK,QL) = QM(I,\ J)$$

If $QMZ = QMA$,  $QM = QLIRPT(QK,QL)$

If $QMZ = QMO$,  $QM = QDIRSM(QK,QL)$

If $QMZ = QMI$,  $QM = QDIRSM(QCOMP(QK),QL)$

If $QMZ = QME$,  $QM = QDIREQ(QK,QL)$

## 5. Multiple quantifiers and their implementation

### (a) Elementary statement in QL-1B

Just as QL-1A employs a quantifier associated with the

logical connective, for a statement composed of relations

between quantifier states, in the form $(q_x x)(\underline{a}x\ \underline{Z}(k,\ell)\ \underline{b}x)$,

the QL-1B type has the standard form,

$$(q_x x)(q_z y)\ (\underline{a}x\ \underline{Z}(k,\ell)\ \underline{b}y) \tag{1}$$

Typical examples are

$$(\forall x)(\exists y)\ (\underline{a}x \dashrightarrow \underline{b}y) \tag{2a}$$

$$(\exists x)(\exists y)\ (\underline{a}x \wedge \underline{b}y) \tag{2b}$$

In our notation and formulation of predicate logic, statements

of the type (2a) and (2b) have a very specific interpretation

which is somewhat different from the standard one ; but all

standard relations can be shown to be expressible, in one form

or the other, in terms of QL-2, QL-1A and QL-1B forms.

Thus, in (2a) and (2b), the variables x and y associated

with $\underline{a}$ and $\underline{b}$ are different in QL-1B, and the quantifier

associated with each is given by the variable that is adopted

for the corresponding quantifier. Otherwise, there is no

significance attached to the sequence of the quantifiers.

Thus for example

$$?a) \quad \equiv \quad (\exists y)(\forall x)(\underline{a}x \rightarrow \underline{b}y) \tag{3}$$

It becomes particularly obvious when this is written in terms

of the disjunction operator as

$$(2a) \quad \equiv \quad (\forall x)(\exists y)(\neg \underline{a}x \lor \underline{b}y) \equiv (\forall x)(\exists y)(\underline{b}y \lor \neg \underline{a}x) \tag{4a}$$
$$\equiv \quad (3) \quad \equiv (\exists y)(\forall x)(\neg \underline{a}x \lor \underline{b}y) \tag{4b}$$

In Eqs. ( 4a) and ( 4b), there is no first term and second term

within the bracket since they are connected by the commutative

operator "or", and no difference can be associated with the

sequence of the quantifiers, which are different in ( 4a) and ( 4b).

This feature will become particularly clear when we give below

the algorithms based on the essential principle of the standard

prenex normal form of quantifier calculus. This principle is,

however, applied in the reverse sense in EVMF, and the QL-1B

form is converted into an equivalent QL-2 statement, which is

then implemented by the techniques of Section 2. We shall

indicate this below, and only the essential formulae, required

for writing the computer programs, are given. A fuller treatment

will be given in a separate report.

## (b) Implementation of QL-1B statements in predicate logic via PQLOG.

The name PQLOG is applied to this sub-program because

it implements F-type statements (QL-1B), but _via_ the process

of converting them into QLOGIC statements (QL-2). The

subroutines for QLOGIC and FLOGIC, given in Sections 3 and 4 above,

are assumed to be available, although they may not all be

required. It is found that, just as in FLOGIC, we consider / only a

notation and formalism for describing a PQLOG relation and its

unary and binary forms of implementation, for the standard SKS

logical connectives of the type $\underline{A}$, $\underline{O}$, $\underline{I}$, $\underline{E}$, and not for a general

2x2 matrix. The relevant algebraic formulae, which are computerized

in Section (c), are given below.

### (i) QL-1B binary relation

A binary relation, in general, has the form

$$\underline{a}x, \underline{b}y, (q_Zx)(q_Zy)(\underline{a}x \ \underline{Z}(k, \ell) \ \underline{b}y) \longmapsto \underline{c} \qquad (5a)$$

which is equivalent to the QL-2 form

$$\underline{a}x, \underline{b}y, \underline{a}x \ \underline{Z}'(k, \ell) \ \underline{b}y \longmapsto \underline{c} \qquad (5b)$$

where
$$Z' = Z , \text{ for } A, O \qquad (5c)$$

and in ( 5b),

$$q(\underline{k}) = q_Zx \text{ or } q^n_Zx , \text{ according as } k = 1, 2 \text{ in } (5a) \qquad (6a)$$

$$q(\underline{\ell}) = q_Zy \text{ or } q^n_Zy , \text{ according as } \ell = 1,2 \text{ in } (5a) \qquad (6b)$$

Extensions of these for $Z = I$, $E$ are available, by using

the relation between implication and disjunction, and the

description of equivalence as the disjunction of $\underline{A}(1, 1)$ and

$\underline{A}(2, 2)$, as in QL-1A.

Thus, $I(1, 1) = O(2, 1)$, so that (5c) continues to be valid,

$Z' = Z = I$, but (6a) gets changed to

$$q(\underline{k}) = q_{\underline{Z}}^{\ell}x \text{ or } q_{\underline{Z}}^{m}x , \text{ according as } k = 1, 2 \text{ in (5a)} \qquad (6c)$$

leaving (6b) unaltered.

Similarly, for $\underline{Z}(k, \ell) = \underline{E}(k, \ell)$, $k, \ell = 1, 2$ , the

relation (5a) becomes

$$(q_{\underline{Z}}x)(q_{Z}y) \ ((\underline{a}x \ \underline{A}(k, \ell) \ \underline{b}y) \ \bigvee (\underline{a}x \ \underline{A}(k^c, \ell^c) \ \underline{b}y)) \qquad (7a)$$

$$= (\underline{a}x \ \underline{A}(\underline{k}, \underline{\ell}) \ \underline{b}y) \ \bigvee (\underline{a}x \ \underline{A}(\underline{k}' , \underline{\ell}') \ \underline{b}y) \qquad (7b)$$

where $q(\underline{k})$, $q(\underline{k}') = q_{\underline{Z}}x$ or $q_{\underline{Z}}^{n}x$, according as $k, k^c = 1, 2$ (7c)

$\quad\quad q(\underline{\ell})$, $q(\underline{\ell}') = q_{Z}y$ or $q_{Z}^{n}y$, according as $\ell, \ell^c = 1, 2$ (7d)

Eqs.(7b,c,d) are valid only for the interpretation (7a) of (5a)

for $\underline{E}(k, \ell)$, and not for $\underline{E}$ expressed as the conjunction of two

implications.

(ii) QL-1B unary relation

This has the form

$$\underline{a}x, (q_zx)(q_zy)(\underline{a}x\ \underline{Z}(k, \ell)\ \underline{b}y) \mapsto by \qquad (8a)$$

which takes the equivalent QL-2 form

$$\underline{a}x, \underline{a}x\ \underline{Z}'(\underline{k}, \underline{\ell}) \mapsto by \qquad (8b)$$

and the relation between Z and Z' follows what was given

above for the binary relation, with the change that, for $\underline{E}(k, \ell)$

$$\underline{a}x, \underline{a}x\ \underline{E}(\underline{k}, \underline{\ell}) = (\underline{a}x\ \underline{A}(k, \underline{\ell}) \oplus \underline{a}x\ \underline{A}(\underline{k}', \underline{\ell}')) \mapsto by \qquad (8c)$$

Essentially, the only new FORTRAN formulae that are needed

for PQLOG are those for the unary, and binary products, and the

SNS truth value of a binary relation, in PQLOG. These three

functions PQBNPT, PQSTV, PQUNPT are described below as PQ(ii)—(iv).

(These are preceded by PQMCON, in the algorithm PQ(i), analogous to

QMCON in Q(xviia)).

(c) Subroutines in PQLOG

PQ(i) : Matrix connective in QLOGIC corresponding to a
        logical relation in PQLOG defined via an SNS
        logical connective : $q(\underline{k}_Z)\ q(\underline{\ell}_Z)\ (\underline{a}\ \underline{Z}(k, \ell)\ \underline{b}) = \underline{Z}(k, \ell)$

PQMCON(QVK,SMX,SK,SL,QVL) = QMZ,QK,QL

For SMX = SMA,SMO,    QMZ = QMA,QMO

QK = QEQU(QVK),QNOT(QVK) according as  SK = S1,S2

QL = QEQU(QVL),QNOT(QVL) according as  SL = S1,S2

For SMX = SMI ,  QMZ = QMI

QK = QEJL(QVK),QCOMP(QVK) according as  SK = S1,S2

QL = QEQU(QVL), QNOT(QVL) according as  SL = S1,S2

No equivalent single logical QL-2 connective is available

for the case of SMX = SME. However, the corresponding unary

and binary relations can be implemented via Eqs 7(a-d) and (8c)

given above in Section 5(b).

Thus, for SMX = SME

$\quad$ PQMCON(QVK,SMX,SK,SL,QVL)

= QMXSUM(PQMCON(QVK,SMA,SK,SL,QVL) ,

$\quad$ PQMCON(QVK,SMA,SCOMP(SK),SCOMP(SL),QVL))

and is expressible as the Boolean sum of two 3x3 matrix connectives.

PQ(ii): PQBNPT for the formulae (5) to (7)

$\quad$ PQBNPT(QVA,QVK,SMX,SK,SL,QVL,QVB) = BC

where

$\quad$ QVA = $\underset{\sim}{a}$, QVB = $\underset{\sim}{b}$, QVK = $(q_z x)$, QVL = $(q_z y)$

For SMX = SMA,SMO,SMI,SME

$\quad$ BC = QBINPT(QVA,QMXP,QVB)

where

$\quad$ QMXP = PQMCON(QVK,SMX,SK,SL,QVL)

The calculation of the SNS truth value PQSTV, of a binary

relation in PQLOG, is also readily obtained in terms of the

above subroutine PQBNPT, as follows.

PQ(iii): PQSTV: - SNS truth value of a binary relation in QL-1B

The relevant subroutine is

$$PQSTV(QVA,QVK,SMX,SK,SL,QVL,QVB) = SVC$$

$$SVC(1) = QBINPT(QVA,QMXP,QVB)$$

$$SVC(2) = QBINPT(QVA,QCOMP(QMXP),QVB)$$

where
$$QMXP = PQMCON(QVK,SMX,SK,SL,QVL)$$

PQ(iv): PQUNPT : Unary product in QL-1B

The unary relation in PQLOG leads to a quantifier state

as output as described in (8a,b,c), and here also, the function

is evaluated by using the equivalent formula QUNPDT in QLOGIC

as follows:

$$PQUNPT(QVA,QVK,SMX,SK,SL,QVL) = QVB$$

This is evaluated as

$$QVB = QUNPDT(QVA,QMXP)$$

For SMX = SMA,SMO,SMI,SME

$$QMXP = PQMCON(QVK,SMX,SK,SL,QVL)$$

Note that for SME, we use the same definition of the equivalent
QL-2 matrix as adopted in PQ(i), and described in the algebra
mentioned in Section 5(b).

In this manner, all elementary statements in QL-1B are
convertible into the QL-2 form, either as a unary product, or
a binary product, and the whole structure of QL-2 can be used
for working them out. It is to be noted that for obtaining
the SNS truth value PQSTV in QL-1B, the complementary relation
is calculated in the equivalent QL-2 form. The need and
consequences of this procedure will be explained in the report
to be prepared, dealing with the improved algebraic theory of
QL-1B.

MATLOG Program in FORTRAN

Part III- Program for GLOGIC and the theory of relations in general

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Part III:  MATLOG Program for GLOGIC and the theory of relations
           in general

## 1.  Introduction

The theory of relations in Boolean algebra employing

m x n Boolean matrices was developed first in MR-42 and has

thereafter been applied in various ways.  Examples of its

application to data processing and graph theory were developed

in ALOG-31 to 34 and the application of general Boolean matrices

of this type to answer various questions in multivalued logic

considered in MR-46A.

In this report, we shall consider essentially the FORTRAN

programing of the relevant equations.  They follow essentially

the same pattern as for SNSLOG and QLOGIC, and deal with m x n

matrices, instead of 2x2 and 3x3 matrices for SNS and QL-2.

Since the numbers  m and n  have to be left open, to be chosen

suitably for each problem, the structure of the programs is

slightly different.  However, the subroutines that are listed

follow closely the sequence that was adopted for SNS and QL-2.

We shall call this as General Logic (GLOGIC) and indicate all

vectors and matrices and other operators in this logic by the

first letter  G  for their symbols.  The basic subroutines

are described in Sections 2 and 3 and thereafter various

applications of these are considered.

## 2. GLOGIC

In the general case, vectors and matrices will have the
format GVA (MI)  and  GMZ (MI,MJ). The former will have
components GVA(I), I = 1 , MI, and the latter GMZ (I,J)
I = 1, MI, J = 1 , MJ.

### (a) Additional subroutines for basic Boolean scalars

We define three functions "Boolean multiple sum" (BMSUM)
"Boolean multiple product" (BMPDT) and "Boolean multiple
Equivalence" (BMEQU) as follows

$$BMSUM(BA(I), MI) = BC$$

stands for    BMSUM = 0,  DO I = 1, MI

BMSUM = BSUM(BMSUM, BA(I))

The functions BMPDT and BMEQU are given by similar equations
as follows

BMPDT = 0, DO  I = 1, MI

BMPDT = BPDT(BMPDT,BA(I))

BMEQU = 1,  DO  I = 1, MI

BMEQU = BEQU(BMEQU, BEQU(BA(I),BA(1)))

### (b) GLOGIC functions and subroutines

i) Boolean sum of two vectors :  $\underline{a} \oplus \underline{b} = \underline{c}$

GUNION (GVA,GVB,MI) = GVC

Check that both GVA and GVB have MI components

BSUM (GVA(I), GVB(I)) = GVC(I)  I = 1, MI

(ii) <u>Boolean product of two vectors</u> : $\underline{a} \otimes \underline{b} = \underline{c}$

   GVIDYA (GVA,GVB,MI) = GVC

   Check that both GVA and GVB have   MI   components

   BPDT (GVA(I),GVB(I)) = GVC(I)   I = 1, MI

(iii) <u>Boolean complement of a vector</u> : $\overline{\underline{a}}^c = \underline{b}$

   GCOMP (GVA,MI) = GVB

   BCMP (GVA(I)) = GVB(I)   I = 1, MI

(iv) <u>Boolean sum of two matrices</u> : $\underline{P} \oplus \underline{Q} = \underline{R}$

   GMXSUM (GMP,GMQ,MI,MJ) = GMR

   Check that both GMP , GMQ  have MI,MJ components

   BSUM (GMP(I,J), GMQ(I,J)) = GMR(I,J)
                        I = 1 , MI
                        J = 1 , MJ

(v)  <u>Boolean product of two matrices</u> : $\underline{P} \otimes \underline{Q} = \underline{R}$

   GMXPDT (GMP,GMQ,MI,MJ) = GMR

   BPDT (GMP(I,J), GMQ(I,J)) = GMR(I,J)
                        I = 1 , MI
                        J = 1 , MJ

(vi) <u>Boolean complement of a matrix operator</u> : $\underline{P}^c = \underline{Q}$

   GMXCMP (GMP,MI,MJ) = GMQ

   BCMP (GMP(I,J)) = GMQ(I,J)   I = 1 , MI
                              J = 1 , MJ

(vii) <u>Scalar product of two vectors</u> $a_1 \otimes b_1 \oplus a_2 \otimes b_2 \oplus \ldots a_m \otimes b_m$

   GSCPDT (GVA,GVB,MI) = BC                          $= C$

   BMSUM(BPDT (GVA(I),GVB(I)),MI) = BC

**viii)** <u>Unary product of a vector with a matrix</u>

$$\bigoplus_i (a_i \otimes z_{ij}) = b_j$$

GUNPDT (GVA,GMZ,MI,MJ) = GVB

Input   GVA(MI), GMZ(MI,MJ)

Output  GVB(MJ)

BMSUM (BPDT (GVA(I),GMA(I,J), MI)),MI) = GVB(J)

$$J = 1, \ MJ$$

**ix )** <u>Binary product of a matrix with two vectors</u>

(Boolean truth value)    $\langle \underline{a} \mid \underline{\underline{Z}} \mid \underline{b} \rangle = c$

GBTV (GVA,GMZ,GVB,MI,MJ) = BC

GVAP = GUNPDT (GVA,GMZ,MI,MJ)

BC = GSCPDT (GVAP,GVB,MJ)

**x )** <u>Matrix product of two matrices</u> : $\mid \underline{\underline{P}} \mid \underline{\underline{Q}} \mid = \mid \underline{\underline{R}} \mid$

GMATPT (GMP,GMQ,MI,MJ,MK) = GMR

BMSUM (BPDT(GMP(I,J),GMQ(J,K)),MJ) = GMR(I,K)

$$I = 1 \ to \ MI , \quad K = 1 \ to \ MK$$

**xi )** <u>Direct product of two vectors</u> : $\underline{a} \ X \ \underline{b} = \underline{\underline{Z}}$

GDIRPT(GVA,MI,GVB,MJ) = GMZ

GMZ(I,J) = BPDT(GVA(I),GVB(J))

$$I = 1, \ MI$$
$$J = 1, \ MJ$$

**xii)** <u>Direct sum of two vectors</u> : $\underline{a} + \underline{b} = \underline{\underline{Z}}$

GDIRSM(GVA,MI,GVB,MJ) = GMZ

GMZ(I,J) = BSUM(GVA(I),GVB(J))

$$I = 1, \ MI$$
$$J = 1, \ MJ$$

xiii) <u>Direct equivalence of two vectors</u>:

$$\underset{\approx}{a} \equiv \underset{\approx}{b} = \underset{\approx}{Z}$$

$$\text{GDIREQ(GVA,MI,GVB,MJ)} = \text{GMZ}$$

$$\text{GMZ(I,J)} = \text{BEQU(GVA(I),GVB(J))}$$

$$I = 1, \text{ MI}$$
$$J = 1, \text{ MJ}$$

xiv ) <u>Scalar vector direct product</u> : $a \times \underset{\approx}{b} = \underset{\approx}{c}$

$$\text{GSVLPT(BA,GVB,MI)} = \text{GVC}$$

$$\text{GVC(I)} = \text{BPDT(BA,GVB(I))}$$

xv ) <u>Scalar vector direct sum</u> : $a + \underset{\approx}{b} = \underset{\approx}{c}$

$$\text{GSVDSM(BA,GVB,MI)} = \text{GVC}$$

$$\text{GVC(I)} = \text{BSUM(BA,GVB(I))} \qquad I = 1, \text{ MI}$$

xvi ) <u>Transpose of a matrix</u> : $\underset{\approx}{P}^t = \underset{\approx}{Q}$

$$\text{GTRANS(GMP,MI,MJ)} = \text{GMQ}$$

$$\text{GMQ(I,J)} = \text{GMP(J,I)}$$

$$I = 1, \text{ MI}$$
$$J = 1, \text{ MJ}$$

xvii) <u>Matrix transformation operators</u> $|Z|^o = |\underset{\sim}{z}|' = {}^e, {}^c, {}^t, {}^{ct}$

Since these operators occur frequently in general classical logic involving multiple element sets, they are given short names in MATLOG, namely GE, GC, GT, GTC. They are defined as follows:

a) $\text{GE(GMZ,MI,MJ)} = \text{GMZ1}$
   $\text{GMZ1(I,J)} = \text{GMZ(I,J)}$

b) $\text{GC(GMZ,MI,MJ)} = \text{GMZ2}$
   $\text{GMZ2,MI,MJ} = \text{GMXCMP(GMZ,MI,MJ)}$

c) $\text{GT(GMZ,MI,MJ)} = \text{GMZ3}$
   $\text{GMZ3,MJ,MI} = \text{GTRANS(GMZ,MI,MJ)}$

d) $\text{GTC(GMZ,MI,MJ)} = \text{GMZ4}$
   $\text{GMZ4,MJ,MI} = \text{GMXCMP(GTRANS(GMZ,MI,MJ)}$

## 3. Subroutines for GCLOG        .6.

The programs **for** General Classical Logic, (GCLOG)

are somewhat different from SNSLOG and are indicated by

general logical relational matrices $\underset{\sim}{A}(\underset{\sim}{k}, \underset{\sim}{\ell})$, $\underset{\sim}{O}(\underset{\sim}{k}, \underset{\sim}{\ell})$, etc.

The  subroutines suggested are as follows.

i) <u>Definition of relational matrices</u> :   $\underset{\sim}{Z}(\underset{\sim}{k}, \underset{\sim}{\ell})$,   Z = A, O, E, I

GMX, GVK, GVL,   X = A, O, E, I

These are analogous to (xvii) SMZ, SK,SL of SNSLOG

GMA,GVK,GVL = GLIRPT(GVK,GVL,MK,ML)

GMO,GVK,GVL = GDIRSM(GVK,GVL,MK,ML)

GME,GVK,GVL = GDIREQ(GVK,GVL,MK,ML) ;

GMI,GVK,GVL = GDIRSM(GCOMP(GVK),GVL,MK,ML)

ii) Relative (SNS) truth value of one term $\underset{\sim}{a}$ for another $\underset{\approx}{u}$ :

---

GRELTV(GVA,GVB,MI) = SVC        $t(\underset{\sim}{a} \mid \underset{\sim}{b}) = c$

This is the general case of (xviii) SRELTV

SVC(1) = GSCPDT(GVA,GVB)

SVC(2) = GSCPDT(GVA,GCOMP(GVB))

iii) SNS truth value of the relation $\underset{\sim}{a}\ \underset{\sim}{Z}\ \underset{\sim}{b}$

---

$\underline{t}(\underset{\sim}{a}\ \underset{\sim}{Z}\ \underset{\sim}{b}) = \underset{\sim}{c} = (c_1, c_2)$

where $c_1 = \langle \underset{\sim}{a} \mid \underset{\sim}{Z} \mid \underset{\sim}{b} \rangle$ ,   $c_2 = \langle \underset{\sim}{a} \mid \underset{\sim}{Z}^c \mid \underset{\sim}{b} \rangle$

GSTV(GVA,GML,GVB,MI,MJ) = SVC

This is the general case of (xix)  SSTV

$$SVC(1) = GBTV(GVA,GMZ,GVB,MI,MJ)$$

$$SVC(2) = GBTV(GVA,GMXCMP(GMZ),GVB,MI,MJ)$$

iv) SNS truth value for $\underset{\approx}{Z} = Z(\underset{\approx}{k}, \underset{=}{\ell})$, $Z = A, O, E, I$ via

relative truth value

$$GSTV2(GVA,GMX,GVK,GVL,GVB,MI,MJ) = SVC$$

If GMX = GMA

$$SVC(1) = BPDT(GSCPDT(GVA,GVK,MI), GSCPDT(GVB,GVL,MJ))$$

$$SVC(2) = BSUM(GSCPDT(GVA,GCOMP(GVK),MI),GSCPDT(GVB,GCOMP(GVL),MJ))$$

If GMX = GMO

$$SVC(1) = BSUM(GSCPDT(GVA,GVK,MI),GSCPLT(GVB,GVL,MJ))$$

$$SVC(2) = BPDT(GSCPDT(GVA,GCOMP(GVK),MI),GSCPDT(GVB,GCOMP(GVL),MJ))$$

If GMX = GME

$$GVKP = GCOMP(GVK)$$

$$GVLP = GCOMP(GVL)$$

$$GVCA = GSTV2(GVA,GMA,GVK,GVL,GVB,MI,MJ)$$

$$GVCB = GSTV2(GVA,GMA,GVKP,GVLP,GVB,MI,MJ)$$

$$SVC(1) = BSUM(GVCA(1), GVCB(1))$$

$$SVC(2) = BPDT(GVCA(2), GVCB(2))$$

If GMX = GMI

$$GVKP = GCOMP(GVK)$$

$$GMX = GMO,GVLP,GVL$$

v) Output of unary relation for $\underset{\approx}{Z}$ ($\underset{\approx}{k}$ , $\underset{\approx}{\ell}$) using the

relative truth value of the input


GUNPT2(GVA,GMX,GVK,GVL,MI,MJ) = GVB

If GMX = GMA

   BAP = GSCPDT(GVA,GVK,MI)

   GVB = GSVDPT(BAP,GVL,MJ)

If GMX = GMO

   BAP = GSCPDT(GVA,GVK,MI)

   GVB = GSVDSM(BAP,GVL,MJ)

If GMX = GME

   GVKP = GCOMP(GVK), GVLP = GCOMP(GVL)

   GVBA = GUNPT2(GVA,GMA,GVK,GVL,MI,MJ)

   GVBB = GUNPT2(GVA,GMA,GVKP,GVLP,MI,MJ)

   GVB  = GUNION(GVBA,GVBB,MJ)

If GMX = GMI

   BAP = GSCPDT(GVA,GCOMP(GVK),MI)

   GVB = GSVDSM(BAP,GVL,MJ)


vi) Binary operator "agree" ($\underset{\approx}{G}$) for the general case

   GMBG(GVA,GVB,MI) = SVC

   BCP(I) = BEQU(GVA(I),GVB(I))   I = 1, MI

   SVC(1) = BMPDT(BCP,MI)

   SVC(2) = BCMP(SVC(1))

vii )  <u>Matrix-Boolean unary operators</u>  $\underset{\sim}{E}$, $\underset{\sim}{N}$, $\underset{\sim}{M}$, $\underset{\sim}{L}$

GEQU(GVA,MI) = GVB

GVB(I) = GVA(I)    I = 1, MI

GNOT(GVA,MI) = GVB

GVB(I) = GVA(MI − I + 1)    I = 1, MI

GCOMP(GVA,MI) = GVB

GVB(I) = BCMP(GVA(I))    I = 1, MI

GELL(GVA,MI) = GVB

GVB(I) = BCMP(GVA(MI - I + 1)    I = 1, MI

## 4. Applications in examples

We shall first give illustrations of the unary and

binary relations involving a general matrix connective

$R_{ij}$ (i = 1 to m, j = 1 to n), taking for this purpose

illustrations from Tables 2 and 4 of IR-42. All of them

are based on the "student-professor" relation and its reverse

relation "professor-student" described in Table 1 of that

reference. For ready reference, the four matrices $\underset{\sim}{P}$, $\underset{\sim}{P}^{c}$

$\underset{\sim}{P}^{t}$, $\underset{\sim}{P}^{tc}$ which are derivable from the relation $\underset{\sim}{P}$, standing

for "professor of", are explicitly described in Table 1 below.

They make use of the four matrix-Boolean operators GE, GT,

GC, GTC as applied to the matrix $\underset{\sim}{P}$ .

Table 1. The four matrices* $\underset{\sim}{P}$, $\underset{\sim}{P}^c$, $\underset{\sim}{P}^t$, $\underset{\sim}{P}^{tc}$ employed for the relations† in the student-professor problem

$$\underset{\sim}{s} \xrightarrow{\underset{\approx}{P}} \underset{\approx}{p} \qquad\qquad \underset{\approx}{p} \xrightarrow{\underset{\approx}{S}} \underset{\approx}{s}$$

$|\underset{\sim}{P}|$ = GM1(5,4) =
= GE(GM1(5,4))

$\begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$

; $|\underset{\sim}{S}|$ = $|\underset{\sim}{P}^t|$ = GM3(4,5) =
= GT(GM1(5,4))

$\begin{pmatrix} 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 \end{pmatrix}$

$$\underset{\approx}{s} \xrightarrow{\underset{\approx}{P}^c} \underset{\approx}{p} \qquad\qquad \underset{\approx}{p} \xrightarrow{\underset{\approx}{S}^c} \underset{\approx}{s}$$

$\underset{\sim}{P}^c$ = GM2(5,4) =
= GC(GM1(5,4))

$\begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$

; $\underset{\sim}{S}^c$ = $\underset{\sim}{P}^{tc}$ = GM4(4,5) =
= GTC(GM1(5,4))

$\begin{pmatrix} 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 & 1 \end{pmatrix}$

*These are typical examples, and the same pattern holds for any matrix relation $|R|$ , and four related matrices $|R|$, $|R^c|$, $|R^t|$, $|R^{tc}|$ are definable.

†The single arrow ( $\longrightarrow$ ) is used to indicate a relation in the general theory of relations, which is to be distinguished from the double arrow ( $\Longrightarrow$ ) for implication, and the arrow with a vertical line ( $\longmapsto$ ) for "leads to ".

## (a) Illustration of unary relations

A number of chosen examples of these are given in

Table 2. The querries are self-explanatory and the relevant

vector-matrix equation that gives the solution to the problem

in each case is given on the r.h.s. The FORTRAN statement

in MATLOG is given below the query, and the input and output

are given below that. In these simple problems, they can be

readily verified to be true by inspection. However, the

FORTRAN statements indicate that a whole variety of related

questions that can be asked, can be expressed in terms of the

basic functions which we have defined in FORTRAN for GLOGIC.

## (b) Illustrations of binary relations in GCLOG

These are summarized in Table 3 which is given below,

and since they are taken from MR-42, no details are given and

only the corresponding Boolean algebraic vector-matrix formula

is given in each case. It can be readily verified that our

library of FORTRAN statements can readily translate them all

into a format suitable for application with our MATLOG subroutines

It is also evident that a whole variety of analogous problems

can be solved by combining these formulae of MATLOG in a

suitable manner.

Table 2. Some examples of unary relations in GCLOG*

*1. Who are all the professors (p') who teach the students in s' ?

$$\langle s \mid P \mid = \langle p \mid$$

GUNPDT(GVS,GM1,MI,MJ) = GVP

GVS = (1 1 0 0 0), GVP = (1 1 1 0)

4. Who are all the non-students(s") of some professors in p' ?

$$\langle p \mid P^{tc} \mid = \langle s \mid$$

GUNPDT(GVP,GM4,MJ,MI) = GVS

GVP = (1 1 0 0) , GVS = (1 1 1 1 0)

5. Who are all the students (s") attending classes taken by professors that teach s'?

$$\langle s1 \mid P \mid p^t \mid = \langle s2 \mid$$

GMATPT(GM1,GM3,MI,MJ,MI) = GM6

GUNPDT(GVS1,GM6,MI,MJ) = GVS2

GVS1 = (1 1 0 0 0) , GVS2 = (1 1 1 0 1)

7. Which professors (p") in p' teach some student in s'?

$$\langle s \mid P \mid \otimes \langle p1 \mid = \langle p2 \mid$$

GUNPDT(GVS1,GM1,MI,MJ) = GVP

GVIDYA(GVP,GVP1) = GVP2

GVS1 = (1 1 0 0 0), GVP1 = (0 1 0 1), GVP2 = (0 1 0 0)

10. Which are the students (s') that take the lectures of p', and do not attend the classes of p"?

$$\langle p1 \mid P^t \mid \otimes \langle p2 \mid P^{tc} \mid = \langle s \mid$$

GUNPDT(GVP1,GM3,MJ,MI) = GVS1
GUNPDT(GVP2,GM4,MJ,MI) = GVS2
GVIDYA(GVS1,GVS2,MI) = GVS
GVP1 = (1 1 0 0), GVP2(0 1 1 0 ), GVS(1 0 1 0 0)

* From Table 2, MR-42.

12. Which students ($\underline{s}'$) are taught by professors that do not belong to the subset $\underline{p}'$?   $\langle p^c | P^t | = \langle s |$

---

GUNPDT(GCOMP(GVP),GM3) = GVS

GVP = (0 1 0 1),   GVS = (1 1 1 0 1)

---

Table 3. Some examples of binary checking formulae in the BVM formulation of general classical relations*

*1 . Check if at least one of $\underline{s}'$ is related by $\underline{P}$ to one of $\underline{p}'$ .   $\langle s | P | p \rangle = 1$

---

GBTV(GVS,GM1,GVP,MI,MJ) = BTV ;   BEQU(BTV,B1) = BC

GVS = (1 1 0 0 0), GVP = (0 1 1 0), BTV = 1, BC = 1 (Yes)

3 . Check if at least one of $p'$ is non-related by $\underline{S}$ to one of $\underline{s}'$.   $\langle p | S^c | s \rangle = 1$

---

GBTV(GVP,GM4,GVS,MJ,MI) = BTV,   BEQ(BTV,B1) = BC

GVP = (0 1 1 0), GVS = (0 0 1 1 0), BTV = 1,   BC = 1 (Yes).

4 . Check if all of $p'$ are non-related by $\underline{S}$ to any of the students in $\underline{s}'$ .   $\langle p | S | s \rangle = 0$

---

GBTV(GVP,GM3,GVS,MJ,MI) = BTV,   BEQ(BTV,B2) = BC

GVP = (0 1 1 0), GVS = (0 0 1 1 0), BTV = 0,   BC = 1 (Yes)

All are non-related.

---

* From Table 4 of MR-42

Table 3 . Contd.

6 .  Check if the professors p'
teaching $\underline{s}'$ have any members
in common with those p"
teaching $\underline{s}''$ .

$\langle s1 \mid P \mid P^t \mid s2 \rangle = 1$

_____

GMATPT(GM1,GM3,MI,MJ,MI) = GM5

GBTV(GVS1,GM5,GVS2,MI,MI) = BTV,  BEQU(BTV,B1) = BC

GVS1 = (1 0 1 0 0), GVS2 = (0 1 0 0 1), BTV = 1 , BC = 1 (Yes

7 .  If the answer to Problem 6 is
"yes", list the set of
professors ($\underline{p}_o$) common to both.

$\langle s1 \mid P \mid = \langle p1 \mid ,$
$\langle s2 \mid P \mid = \langle p2 \mid ,$
$\langle p1 \mid \otimes \langle p2 \mid = \langle p_o \mid$

_____

GUNPDT(GVS1,GM1,MI,MJ) = GVP1

GUNPDT(GVS2,GM1,MI,MJ) = GVP2

GVIDYA(GVP1,GVP2,MJ) = GVPO

GVP1 = (1 0 1 1), GVP2 = (1 1 1 0), GVPO = (1 0 1 0)

8 .  Check if $\underline{s}'$ contains all the
students studying under the
professors ($p'$) who teach any
of them.

$\mid P \mid s \mid = \mid K \mid ,$
$\langle s \mid K \mid s^c \rangle = 0$

_____

GMATPT(GM1,GM3,MI,MJ,MI) = GMK

GBTV(GVS1,GMK,GCOMP(GVS1),MI,MJ) = BTV,  BEQU(BTV,B2) =BC

GVS1 = (1 0 0 0 1) , BTV = 1, BC = 0  (No).

9.  If "No", list the set of
professors (p') and also the
full list of students ($\underline{s}_o$)
studying under them.

$\langle s1 \mid P \mid = \langle p1 \mid$
$\langle p1 \mid S \mid = \langle s_o \mid$

_____

GUNPDT(GVS1,GM1,MI,MJ) = GVP1

GUNPDT(GVP1,GM3,MJ,MI) = GVSO

GVP1 = (1 1 1 0), GVSO = (1 1 1 0 1)

Before considering the particular case of GCLOG

which is relevant for clausal relations, and which employ

the subroutines involving GMX,GVK,GVL described above, we

shall make some comments regarding a general problem, namely

that of the application of Boolean connective operators

$\bigoplus$, $\bigotimes$ , $^c$ to matrix connectives, in addition to vectors

representing terms. This is given in the next Section 5

and some examples are given, after which the special case of

clausal relations, and their generalization into four types

as explained in MR-57, will be taken up.


## 5. Operations for the combination of terms and relations
### (a) Boolean operations on matrices

Quite apart from the vector-matrix products which are

relevant for the "matrix" type connective $\left| \underset{\sim}{P} \right|$ and the related

matrices $\left| \underset{\sim}{P}^c \right|$, $\left| \underset{\sim}{S} \right|$ and $\left| \underset{\sim}{S}^c \right|$, which were considered above,

we have also made use of "Boolean" connective operators denoted

by the symbols $\bigoplus$ , $\bigotimes$ and $^c$ in some of the problems

considered in Tables 2 and 3. However, in these applications,

these Boolean connectives follow the same definitions as in

MR3 and ML-2, namely that they correspond to the combination

of information regarding the same entity, but coming from

different sources. When applied to "terms" represented by

vectors in EVMF, the two operators, $\oplus$ , standing for

"union" ($\underset{\sim}{U}$) , and $\otimes$ , standing for "vidya" ( $\underset{\sim}{V}$ ) , correspond

to the union and intersection of the sets A1 and A2, whose

members are represented by the 1's in the corresponding

vectors $\underset{\sim}{a1}$ and $\underset{\sim}{a2}$ respectively. Thus, we obtain the resultant

set A, represented by $\underset{\sim}{a}$ , given by the vector equations

$$\underset{\sim}{a1} \oplus \underset{\sim}{a2} = \underset{\sim}{a} \quad , \quad \underset{\sim}{a1} \otimes \underset{\sim}{a2} = \underset{\sim}{a} \qquad (1)$$

in these two cases.

Similarly, it is also possible to apply the operations

of "union" and "vidya" to "relations" represented by matrices.

We have already employed this intuitively in SNS and QLOG

(e.g. in the definition of the equivalence operator), but

this requires more specific definition and description for the

general case. Thus, the operations $\oplus$ and $\otimes$ are required

in situations involving the simultaneous existence of two

relations $\underset{\sim}{R1}$ and $\underset{\sim}{R2}$ as applied to the same state vector $\underset{\sim}{a}$ ,

leading to two outputs $\underset{\sim}{b1}$ and $\underset{\sim}{b2}$ respectively for $\underset{\sim}{b}$ . The

question then arises : "What is the effective resultant output

vector $\underset{\sim}{b}$ if the two relations are both applicable to a problem,

either conjunctively, or disjunctively, as the case may be?"

As we shall see below, this can be done by applying the

Boolean sum, or product, operation to the corresponding elements

of the connective matrices $\underset{\sim}{R1}$ and $\underset{\sim}{R2}$ respectively. Thus,

for two relations in conjunction

$$\underset{\sim}{a}\,\underset{\sim}{R1} = \underset{\sim}{b1}\ ,\ \ \underset{\sim}{a}\,\underset{\sim}{R2} = \underset{\sim}{b2}\ \longmapsto\ \underset{\sim}{a}\,(\underset{\sim}{R1} \otimes \underset{\sim}{R2}) = \underset{\sim}{b}\ ,\ \ \ \ (2a)$$

and for two relations in disjunction, they lead to

$$\underset{\sim}{a}\,(\underset{\sim}{R1} \oplus \underset{\sim}{R2}) = \underset{\sim}{b}\ \ \ \ \ \ \ \ \ \ \ \ \ (2b)$$

We have already defined the Boolean operations of sum,

and product, as applied to matrices, in the subroutines

GMXSUM and GMXPDT, and they are definable as in (3a) and (3b)

below :

Boolean product: $\ \ R1_{i,j} \otimes R2_{i,j} = R_{i,j}$, $\ \ \ \ \ \ \ \ \ \ (3a)$

Boolean sum : $\ \ R1_{i,j} \oplus R2_{i,j} = R_{i,j}$, $\ \ \ \ \ \ \ \ \ \ (3b)$

for each $(i,j)$, $i = 1$ to $m$ , $j = 1$ to $n$ $\ \ \ \ \ \ \ \ \ \ (3c)$

The logical reason as to why the Boolean operation

has to be applied between the components of the matrix

elements representing the connectives, rather than vector

components of the two outputs $\underset{\sim}{b1}$ and $\underset{\sim}{b2}$, in order to obtain $\underset{\sim}{b}$

can be given as follows. A general matrix relation $\underline{a}\,\underline{R} = \underline{b}$

in GLOGIC really stands for the set of mxn relations

$a_i\,R_{ij} = b_j$ , $i = 1$ to m, $j = 1$ to n, in which each individual

element $a_i$ of the input is either related, or non-related, to

each individual element $b_j$ of the output, according as

$R_{ij} = 1$, or 0. Therefore, when two relations $R1_{i,j}$ and $R2_{i,j}$

are coexistent, then each equation $a_i\,R_{ij} = b_j$ leading to

$b_j$ from $a_i$ has to be treated separately. Effectively, we thus

have mxn equations, for all pairs (i,j) with $i = 1$ to m,

$j = 1$ to n, as in (4a) and (4b), corresponding to (3a) and (3b)

respectively.

$$b_j = a_i\,R1_{i,j} \oplus a_i\,R2_{i,j} = a_i(R1_{i,j} \oplus R2_{i,j}) \qquad (4a)$$

$$b_j = a_i\,R1_{i,j} \otimes a_i\,R2_{i,j} = a_i(R1_{i,j} \otimes R2_{i,j}) \qquad (4b)$$

These evidently lead to the matrix equations (2a) and (2b),

respectively, in BVMF, for the general case.

To complete the discussion, we shall indicate that the

Boolean complementation operator, applied to a <u>relation</u>, leads

from the equation $\underline{a}\,\underline{R} = \underline{b}$ to $\underline{a}\,\underline{R}^c = \underline{b}'$ , where the elements

of $R_{ij}$ are negated, as in

$$\underline{R}^c_{i,j} = (\underline{R}_{i,j})^c \qquad (5)$$

It should be noted that $\underset{\sim}{b}' \neq \underset{\sim}{b}^c$ in general, and the two operations, of complementing a relation, and of complementing the output, represented by the MATLOG functions GMXCMP and GCOMP respectively, are entirely different from one another.

If these distinctions regarding the Boolean operations as applied to terms and relations are borne in mind and properly applied, then practically any problem involving relations in GLOGIC can be treated, requiring only one or more of the operations that have been described above in MATLOG. We shall give a few illustrative examples below.

Table 4.　Illustrations of Boolean operations on connectives in GCLOG

Inputs to the problem

Two sets represented by the vectors $\underset{\sim}{a}$ and $\underset{\sim}{b}$ are related by two different properties, namely

$\underset{\sim}{P}$ : $b_j$ is the professor of $a_i$

$\underset{\sim}{E}$ : $b_j$ is the examiner of $a_i$

$$|\underset{\sim}{P}| = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix} \quad , \quad |\underset{\sim}{E}| = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

A : Conjunction of two relations

Find those members of the set, / $\underline{b3}$ , who are both professors of, as well as the examiner of the set represented by

vector $\underline{a}$ = (0 1 1 0 0).

The relevant matrix is

$$|\underline{R1}| = |\underline{P}| \otimes |\underline{E}| = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

and $\langle \underline{b3}| = \langle \underline{a}|\underline{R1}| = (0\ 0\ 0\ 1)$

Note that the relations $\underline{P}$ and $\underline{E}$ individually lead to the sets $\underline{b1}$ and $\underline{b2}$ given by the vectors

$$\langle \underline{b1}| = (1\ \ 1\ \ 0\ 1), \quad \langle \underline{b2}| = (1\ \ 1\ \ 1\ \ 1)$$

and

$$\langle \underline{b3}| \neq \langle \underline{b1}| \otimes \langle \underline{b2}| = \langle \underline{b4}| = (1\ 1\ 0\ 1)$$

B : Disjunction of two relations

Find the representive vector $\underline{b5}$ corresponding to the members of the set $\underline{b}$ , who are either professors of at least one member of $\underline{a1}$ , or are the examiners of one member of $\underline{a1}$ .

Clearly, the relevant matrix is

$$|\underset{\sim}{R2}| = |\underset{\sim}{P}| \oplus |\underset{\sim}{E}| = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

and $\langle\underset{\sim}{b5}| = \langle\underset{\sim}{a}|\underset{\sim}{R2}| = (1 \quad 1 \quad 1 \quad 1)$

Note that, in this case, $\langle\underset{\sim}{b5}| = \langle\underset{\sim}{b1}|\oplus\langle\underset{\sim}{b2}|$ , which can be

proved quite generally for the Boolean _sum_ of two relations.

If we represent $\underset{\sim}{P}$ and $\underset{\sim}{E}$ by GMP and GMC, and

$\underset{\sim}{a}$ and $\underset{\sim}{b}$ by GVA and GVB, then the FORTRAN statement for

calculating $\overline{\underset{\sim}{b3}}$ in MATLOG is

GVB3 = GUNPDT(GVA,GMXPDT(GMP,GME),MI,MJ))

while the corresponding statement for calculating $\underset{\sim}{b4}$ is

GVB4 = GVIDYA(GUNPDT(GVA,GMP), GUNPDT(GVA,GME), MI,MJ)

The MATLOG formula for calculating $\underset{\sim}{b5}$ is, similarly,

GVB5 = GUNPDT(GVA,GMXSUM(GMP,GME),MI,MJ)

## C. Complementation of a relation

The Boolean complement has already been discussed in

Table 2, both for a matrix relation, as well as for a vector term.

In particular, for our problem, the complement of $\underset{\sim}{P}$ leads

to the BVMF equation

$$\langle\underset{\sim}{a}|\underset{\sim}{P}^c| = \langle\underset{\sim}{b6}|$$

giving the set of $\underset{\sim}{b}$'s who do not take classes for some of the $\underset{\sim}{a}$'s in the set under consideration.

To illustrate the distinction between $\langle \underset{\sim}{a} \mid \underset{\sim}{p}^c \mid = \langle \underset{\sim}{b6} \mid$ and $(\langle a \mid \underset{\sim}{p} \mid)^c = \langle \underset{\sim}{b7} \mid$ , we note that they are programable in MATLOG by the statements

    GVB6 = GUNPDT(GVA,GMXCMP(GMP),MI,MJ)
and
    GVB7 = GCOMP(GUNPDT(GVA,GMP), MI, MJ))

Taking the example given above, we obtain

$$\langle \underset{\sim}{b6} \mid = (0\ 1\ 1\ 0\ 0) \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} = (1\ 1\ 1\ 1)$$

$$\langle \underset{\sim}{b7} \mid = \langle \underset{\sim}{b1} \mid^c = (1\ 1\ 0\ 1)^c = (0\ 0\ 1\ 0)$$

The difference between $\underset{\sim}{b6}$ and $\underset{\sim}{b7}$ arises from the following distinction:

$\underset{\sim}{b6}$ : Set of all b(j)'s who do not teach at least one of the members of the set - a(i).

$\underset{\sim}{b7}$ : Set of all b(j)'s who do not belong to the set $\underset{\sim}{b1}$ that teach at least one member of the set (ai), i.e. who do not teach any member of a(i).

It is very satisfying to know that such hair-splitting differences can be straightaway programed in MATLOG, once the logic is written in BVMF notation.

(b)  Combination of two matrix relations

In the previous subsection 5(a), we had considered the

combination of two matrices via the Boolean operations of sum

and product as applied to relations.  However, these are

the
not/only ways in which two matrix relations can be combined.

                                          the
There are more possibilities which belong  to/matrix type

                            the
of operations rather than/Boolean type of operations.  We shall

describe these below with special reference to two problems

and then consider the nature of such combinations more generally.


(i) Transitive combination of two matrix relations

Consider three sets A, B, C  whose members are represented

by the vectors $\underset{\sim}{a}(i)$, $\underset{\sim}{b}(j)$, $\underset{\sim}{c}(k)$, i = 1 to I, j = 1 to J,

k = 1 to K, which are related by two successive relations

as in (6):

$$\underset{\sim}{a}(i) \ \underset{\sim\sim}{R1}\ (i,j) = \underset{\sim}{b}(j) \tag{6a}$$

$$\underset{\sim}{b}(j) \ \underset{\sim\sim}{R2}\ (j,k) = \underset{\sim}{c}(k) \tag{6b}$$

Then the matrix relation directly connecting $\underset{\sim}{a}(i)$ with $\underset{\sim}{c}(k)$ is

$$\underset{\sim}{a}(i) \ \underset{\sim\sim}{R3}(i,k) = \underset{\sim}{c}(k) \tag{7a}$$

where

$$\left| \underset{\sim\sim}{R3} \right| = \left| \underset{\sim\sim}{R1} \right| \ \underset{\sim\sim}{R2} \ \bigg| \tag{7b}$$

In fact, this type of combination of two matrices has already been defined in one of the MATLOG functions listed above, namely GMATPT whose description is as follows. The function giving $|\underset{\sim}{R3}|$ = GMR3 , in terms of $|\underset{\sim}{R1}|$ = GMR1 , $|\underset{\sim}{R2}|$ = GMR2 is

GMR3 = GMATPT(GMR1, GMR2, MI,MJ,MK)

Simple examples of GMATPT are:

$\underset{\sim}{R1}$ = "parent of", $\underset{\sim}{R2}$ ="parent of" $\longmapsto$ $\underset{\sim}{R3}$ = "grandparent of"  (8a)

$\underset{\sim}{R1}$ = "parent of", $\underset{\sim}{R2}$ ="brother of" $\longmapsto$ $\underset{\sim}{R3}$ = "uncle of"  (8b)

It is bovious that, having obtained the matrix $\underset{\sim}{R3}(i,k)$ , we can forget $\underset{\sim}{R1}$ and $\underset{\sim}{R2}$ from which this was derived, and consider it as a single matrix relation and manipulate it in all the processes described above. For example, $\underset{\sim}{R3}^{t}$ will be the relation "nephew or niece of". It can be extended to any number of such successive relations connected in a transitive fashion. Therefore, this type of combination finds much application in predicate calculus.

## (ii) Tensor combination of two matrix relations

This does not strictly come under the purview of this report, but is included here so as to make the list of possible

ways in which two relations can be combined fairly complete.

It is illustrated by the following examples. Consider three

sets  A, B1, B2  having the relations represented by the

$\mathcal{B}V\!M\!F$ equations

$$\underline{a}\ \underline{R1} = \underline{b1}\ ,\quad \underline{a}\ \underline{R2} = \underline{b2} \tag{9a}$$

where

$$\underline{a} = \underline{a}(i),\ \underline{b1} = \underline{b1}(j),\ \underline{b2} = \underline{b2}(k) \tag{9b}$$

A good example is,

$$c.(i) = child,\quad f(j) = father,\quad m(k) = mother \tag{10a}$$

Then we can talk of the relation between a child and the "parents"

(father and mother) of its parents, which we shall

represent by $\underline{p}(j,k)$. Then this relation takes the form

$$\underline{c}(i)\ \underline{R3}(i,j,k) = \underline{p}(j,k) \tag{10b}$$

It is obvious that the entity $\underline{R}(i,j,k)$ is not a rectangular

matrix, but a third order tensor in this particular case.

More generally, a general tensor relation can be built up of

elementary relations of the type given in (9), each of which

requires a second order tensor, or rectangular matrix, for its

definition. An introduction to such general tensor relations

in Boolean algebra has been given in MR-59, and the implementation

of such relations/for practical problems will be considered later.
generally

Two examples are given below for the relations "parent"

standing for "either father or mother", and "parents"

corresponding to "the pair consisting of father and mother"

of child c(i). These are described in the Sections (c) and (d)

below.


## (c) Union of two matrix relations

The relation "parent of" is not a third order tensor

relation, but the Boolean sum of two matrix relations. Taking

(10a), if


c(i) has i = 1 to I , f(j) has j = 1 to J, m(k) has k = 1 to K

$$\tag{11a}$$

we can define $p(\ell)$ for "parent" with the domain


$$p(\ell) \ , \quad \ell = 1 \text{ to } L = J + K \tag{11b}$$

where

$$p(\ell) = f(j) \quad \text{for } \ell = j \quad , \quad j = 1 \text{ to } J \tag{11c}$$

and

$$p(\ell) = m(k) \quad \text{for } \ell = J + k \ , \quad k = 1 \text{ to } K \tag{11d}$$


Then, in the extended domain $(i, \ell)$, the relations $\underset{\sim}{R1'}$

(father of) and $\underset{\sim}{R2'}$ (mother of), as in (12a,b)

$$\underset{\sim}{R1}'(i, \ell) = |c(i)\rangle \otimes \langle p(\ell)|, \quad 1 \leqslant \ell \leqslant J \tag{12a}$$

$$\underset{\sim}{R2}'(i, \ell) = |c(i)\rangle \otimes \langle p(\ell)|, \quad J + 1 \leqslant \ell \leqslant L \tag{12b}$$

will lead to $\underset{\sim}{R3}'$ for "parent of", as in (13)

$$\underset{\sim}{R3}'(i, \ell) = \underset{\sim}{R1}'(i, \ell) \oplus \underset{\sim}{R2}'(i, \ell) \tag{13}$$

This formalism can be suitably applied for all problems requiring the relation between the union of sets such as F and M into a single set $P = F \cup M$, and a third set such as C, in the above example.

### (d) Expanded product of two matrix relations

On the other hand, the relation "parents of" is a tensor of the third order, obtained by the expansion of the two matrices $\underset{\sim}{R1}(i, j)$ and $\underset{\sim}{R2}(i, k)$ to yield $\underset{\sim}{R3}(i,j,k)$ whose components R3(i,j,k) are given by (14).

$$R3(i,j,k) = R1(i, j) \times R2(i, k) \tag{14}$$

Then, Eq.(10b) holds for obtaining the pair $p(j, k) =$ "f(j) and m(k)" of the child c(i). We shall consider the MATLOG statement for this in due course in a later more comprehensive report of tensor relations.

(e) Tensor relation $\underset{\sim}{R}(i,j,k)$ applied in three ways

It is to be noted, however, that formally, the tensor

components $R(i,j,k)$ correspond to a mutual relation between

$c(i)$, $f(j)$, $m(k)$, and we can obtain equally well the

"mother-child" combinations who are related to a given father

$f(j)$ by the equation (15).

$$\underset{\sim}{f(j)} \quad \underset{\sim}{R}(i,j,k) \;=\; \underset{\sim}{q}(i, k) \quad (\equiv \text{the combinations} \quad c(i), \; m(k)) \quad (15)$$

So also, the "father-child" combinations related to a set of

given mothers $\underset{\smile}{m(k)}$ is given by Eq.(16),

$$\underset{\smile}{m(k)} \quad \underset{\sim}{R}(i,j,k) \;=\; \underset{\pmb{r}}{r}(i, j) \tag{16}$$

These also can be incorporated in programs, but will be deferred

for a more comprehensive treatment of tensor relations.

## 6 . Generalized Clausal Relations

In all that has been considered so far, we have used only

one type of relation, for an mxn matrix $R(i, j)$, namely the

standard form of a relation :

$$a_1 \wedge a_2 \wedge \ldots \wedge a_m \xrightarrow{R} b_1 \vee b_2 \vee \ldots \vee b_n \tag{17}$$

in which R(i, j) is a general mxn matrix, including in

particular, the case of the matrix corresponding to the

clausal form, namely

$$( \underset{\sim}{R} = \underset{\sim}{k} \times \underset{\sim}{\ell} ) \equiv (GMZ = GDIRPT(GVK,GVL,MI,MJ)) \tag{18}$$

As already mentioned, the MATLOG equation for (17) takes

the forms (19a) and (19b) for unary and binary implementation of (17)

Unary : GUNPDT (GVA, GMZ, MI, MJ) = GVBP                    (19a)

Binary : GBTV(GVA, GMZ, GVB, MI, MJ) = BC                    (19b)

However, as considered in MR-57, (17) has three more

analogues, leading to$^a$/set of four types of relations R1, R2,

R3, R4 as below.

$$R1 : a_1 \wedge a_2 \wedge \ldots \wedge a_m \xrightarrow{R} b_1 \vee b_2 \vee \ldots \vee b_n \tag{20a}$$

$$R2 : a_1 \vee a_2 \vee \ldots \vee a_m \xrightarrow{R} b_1 \vee b_2 \vee \ldots \vee b_n \tag{20b}$$

$$R3 : a_1 \wedge a_2 \wedge \ldots \wedge a_m \xrightarrow{R} b_1 \wedge b_2 \wedge \ldots \wedge b_n \tag{20c}$$

$$R4 : a_1 \vee a_2 \vee \ldots \vee a_m \xrightarrow{R} b_1 \wedge b_2 \wedge \ldots \wedge b_n \tag{20d}$$

For implementing these in MATLOG, three additional

subroutines are needed in GLOGIC, and one extra one, with four

parts, in GCLOG. These are as follows:

Add under GLOGIC

(vii a) <u>General conjunctive product (Type A) of two vectors</u>

$$(a_1^c \otimes b_1 \oplus a_2^c \otimes b_2 \oplus \ldots \oplus a_m^c \otimes b_m)^c = c$$

GCJPTA(GVA,GVB, MI) = BC

BCMP(GSCPDT(GCOMP(GVA),GVB, MI)) = BC


(vii b) <u>General conjunctive product (Type B) of two vectors</u>

$$(a_1 \otimes b_1^c \oplus \ldots \oplus a_m \otimes b_m^c)^c = c$$

GCJPTB(GVA,GVB,MI) = BC

BCMP(GSCPDT(GVA,GCOMP(GVB), MI)) = BC


(viiia) <u>Unary conjunctive product of a vector with a matrix</u>

$$( \bigoplus_i (a_i^c \otimes z_{ij}))^c = b_j$$

GUNCJP(GVA,GMZ,MI,MJ) = GVB

GCOMP(GUNPDT(GCOMP(GVA), GVZ,MI,MJ)) = GVB


Add under GCLOG

(viii) <u>Boolean truth value of a relation for given inputs</u>

$\underset{\sim}{a} \ \underset{\sim}{Z} \ \underset{\sim}{b} = c$ , for the four types of relations

R1, R2, R3, R4. It is better to have four different

functions, rather than using a variable having the

values 1 to 4. Thus, we have

(viii a)  Type  R1

GBTVR1 (GVA,GMZ,GVB,MI,MJ) = BC

GUNPDT(GVA,GMZ,MI) = GVBP

GSCPDT(GVBP,GVB,MJ) = BC


(viii b)  Type  R2

GBTVR2(GVA,GMZ,GVB,MI,MJ) = BC

GUNCJP(GVA,GMZ,MI) = GVBPP

GSCPDT(GVBPP,GVB,MJ) = BC


(viii c)  Type  R3

GBTVR3(GVA,GMZ,GVB,MI,MJ) = BC

GUNPDT(GVA,GMZ,MI) = GVBP

GCJPTB(GVBP,GVB,MJ) = BC


(viii d)  Type  R4

GBTVR4(GVA,GMZ,GVB,MI,MJ) = BC

GUNCJP(GVA,GMZ,MI) = GVBPP

GCJPTB(GVBPP,GVB,MJ) = BC


The theory of these is given completely in MR-57, and
therein
Table 3 on page 27/contains the formulae which have been
converted to MATLOG notation in this section.

7.  Aristotle's Syllogisms in GLOGIC

As mentioned in MR-57, the generalized clausal relations of

Types 1 and 2 can be expressed via quantifiers as follows:

Type R1 :   $(\forall i)(ai) \xrightarrow{R} (\exists j)(bj)$    (I-type)          (21a)

Type R2 :   $(\exists i)(ai) \xrightarrow{R} (\exists j)(bj)$    (A-type)          (21b)

where the notation A, I follows standard symbolism adopted for

the names of syllogisms such as Barbara, Darii etc. (See MR 34).

Thus

Type A :   All A are B $\Longleftrightarrow$ If there is an ai present, then (22a)
                                        there is a bj present

Type I :   Some A are B $\Longleftrightarrow$ If all ai are present, then    (22b)
                                        there is a bj present

To this may be added the two classical types, E and O, of

categorical statements, as in (22c,d).

Type E :   All A are not B $\Longleftrightarrow$ If there is an ai present, then
                                        there is a non-bj present (22c)

Type O :   Some A are not B $\Longleftrightarrow$ If all ai are present, then
                                        there is a non-bj present (22d)

These correspond, in terms of quantifiers as :

Type R1 :   $(\forall i)(ai) \xrightarrow{R} (\exists j)(\neg bj)$    (O-type)          (21c)

Type R2 :   $(\exists i)(ai) \xrightarrow{R} (\exists j)(\neg bj)$    (E-type)          (21d)

An analysis of the logical content of (22a-d) indicates

that the relation R in (21a-d) has the form ai - bj .

However, without writing out in detail the matrix R, the

statements (21a-d) can be taken over in QL-2 formulation, and

the 18 categorical syllogisms that are valid, out of the 64

that can be stated, can be worked out _via_ QL-2 algebra.

(This will be discussed in detail in a later theoretical report.)

We shall show how this can be deduced using MATLOG, as below.

From the form of the relations (21a-d), the four 3x3

matrices for categorical statements of the type A, E, I, O

can be calculated (using (xvii) of MATLOG-2, MR-61) :

$$\text{Type } \underline{A} \; : \; \text{QMZA} = \text{QMI,Q6,Q6} = \text{QMO,Q5,Q6} = \begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{pmatrix} \qquad (23a)$$

$$\text{Type } \underline{E} \; : \; \text{QMZE} = \text{QMI,Q6,Q2} = \text{QMO,Q5,Q2} = \begin{pmatrix} 0 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix} \qquad (23b)$$

$$\text{Type } \underline{I} \; : \; \text{QMZI} = \text{QMI,Q1,Q6} = \text{QMO,Q2,Q6} = \begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix} \qquad (23c)$$

$$\text{Type } \underline{O} \; : \; \text{QMZO} = \text{QMI,Q1,Q2} = \text{QMO,Q2,Q2} = \begin{pmatrix} 0 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix} \qquad (23d)$$

We shall first show that the two classical syllogisms

which are basic to the working out of all valid ones in

Aristotle's list, namely <u>Barbara</u> and <u>Darii</u>, follow from the

form of the matrices in (23a-d). In words, these two are

as follows:

All A is B,  All B is C  $\Longleftrightarrow$  All A is C    (<u>Barbara</u>)    (24a)

Some A is B,  All B is C  $\Longleftrightarrow$  Some A is C   (<u>Darii</u>)    (24b)

The corresponding MATLOG equations to be checked are :

QMATPT(QMZA, QMZA)  =  QMZA                                      (25a)

QMATPT(QMZI, QMZA)  =  QMZI                                      (25b)

These are seen to be valid from (23a,c).

Other syllogisms follow similarly. For example, <u>Celarent</u>

follows as in (25c), and <u>Ferio</u> from (25d).

QMATPT(QMZA, QMZE)  =  QMZE                                      (25c)

QMATPT(QMZI, QMZE)  =  QMZO                                      (25d)

So also, the non-existence of any syllogisms of the Aristotelean

type other than the 18 listed in MR-34 can be proved by showing

that the six others in Table 3, page 33 of MR-34 (other than

Sl. No.1, <u>Barbara</u> and Sl. No. 3, <u>Darii</u>) lead to no definite

conclusions. Thus, as will be seen from (26a-f) below, the

output matrix is QMDDD in all these cases, and it is a matrix

that will give an output Q8 = (1 1 1), corresponding to the

state $\triangle$ for universal doubt, whatever may be the input (except

the impossible state (0 0 0)).

| Sl. No. 2: | QMATPT(QMZE, QMZA) | = | QMDDD | (26a) |
|---|---|---|---|---|
| Sl. No. 4: | QMATPT(QMZO, QMZA) | = | QMDDD | (26b) |
| Sl. No. 5: | QMATPT(QMZA, QMZI) | = | QMDDD | (26c) |
| Sl. No. 6: | QMATPT(QMZE, QMZI) | = | QMDDD | (26d) |
| Sl. No. 7: | QMATPT(QMZI, QMZI) | = | QMDDD | (26e) |
| Sl. No. 8: | QMATPT(QMZO, QMZI) | = | QMDDD | (26f) |

where

$$QMDDD = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix} \qquad (26g)$$

These simple calculations using MATLOG give a proof of the

consistency and completeness of all the syllogisms of the

Aristotlean type in classical logic. The precise conditions in

BVMF, under which the additional 18 syllogisms of the "new" type

listed in MR-34 are valid, have not yet been worked out.

## 8. Mixed statements containing quantified and unquantified terms via GLOGIC

The formulae given below are based on the treatment

in MR-53. In quantifier algebra (QL-2) we can have statements

of the type QSN and SNQ with 3x2 matrices and 2x3 matrices

respectively. These are best treated using the GLOGIC formulae

developed in this report. We only give below the MATLOG

statements which will serve for QSN unary and binary products

and SNQ unary and binary products, with the definition of

QSN and SNQ matrices.via the logical relations represented

by them. They are not given any number since, in each case,

a MATLOG statement in GLOGIC using the formulae given in

Sections 2 and 3 will serve the purpose. They are listed below.

QSN — unary product :  $\underline{a} \ \underline{Z} = \underline{b}$

      GUNPDT(QVA,QSMZ, 3,2) = SVB

QSN — binary product :  $\underline{a} \ \underline{Z} \ \underline{b} = c$

      GBTV(QVA,QSMZ,SVB, 3,2) = BC

QSN — relational matrices :  $\underline{k} \times \underline{\ell}, \ \underline{k} + \underline{\ell}, \ \underline{k} \equiv \underline{\ell}, \ \underline{k} \longrightarrow \underline{\ell}$

      QSMZ = QSMX,  QVK,SVL = GMX,GVK,GVL

        where  GVK = QVK,  GVL = SVL

<u>SNQ — unary product</u> :  $a \, \underline{Z} = b$

       GUNPDT(SVA,SQMZ, 2, 3)  =  QVB

<u>SNQ — binary product</u> :  $a \, \underline{Z} \, b = c$

       GBTV(SVA,SQMZ,QVB, 2,3)  =  BC

<u>SNQ — relational matrices</u> :  $k \times \ell$, $k + \ell$, $k \equiv \ell$, $k \longrightarrow \ell$

       SQMZ = SQMX,SVK,QVL = GMX,GVK,GVL,

       where GVK = SVK, GVL = QVL

Other logical functions such as relative truth value, SNS

truth value of a relation etc., can also be formulated

using suitable subroutines from GLOGIC. Details are not given.

As discussed in MR-53, it can happen that the input for

a quantified term in a relation is a non-quantified term, and

vice-versa, and similarly for the output. However, the necessary

formulae for this purpose which were developed in MR-53 are

based on <u>logic</u> and are expressible in Boolean <u>algebra</u> according

to the logical conditions that are imposed. Since this is a

problem in logic, it is deferred for consideration in the

detailed report on the application of BVMF for practical problems.

## 9. Comment on the combination of two relations in general

In Part II of MATLOG series (MR-61), it was mentioned in Section 2, page 12, that, in general, the conjunction of the Boolean truth values of two relations is not necessarily equivalent to the Boolean truth value of the conjunction of the two relations, while the corresponding result is always true for disjunctions, and that this will be considered in Part III dealing with GLOGIC. This has in fact been done in Section 5, page 20 of this report, where, in a practical example, the non-equivalence of the two types of conjunction, as also the equivalence of the two types of disjunction, has been illustrated. A general proof of this is reserved for a theoretical report dealing with theory of relations in general. However, the proof that the two ways of combining two relations will lead to the same result for the output when one of the conditions (a), (b), (c) of Section 4, page 35, of MATLOG-1 (MR-60) is satisfied, can be given, as mentioned therein, for GLOGIC, and hence for QLOGIC and SNSLOG. This is briefly indicated below.

## (a) Unary relations and Boolean truth values of binary relations

Following the notation employed in this part, we shall

consider, in particular, the following unary relations,

$$\underset{\sim}{a} \; \underset{\sim}{R1} = \underset{\sim}{b1} \; , \quad \underset{\sim}{a} \; \underset{\sim}{R2} = \underset{\sim}{b2} \qquad \qquad (27a,b)$$

which lead,by conjunction and disjunction of the relations

R1 and R2, to (28a,b):

$$\underset{\sim}{a} \; (\underset{\sim}{R1} \otimes \underset{\sim}{R2}) = \underset{\sim}{c} \; , \quad \underset{\sim}{a} \; (\underset{\sim}{R1} \oplus \underset{\sim}{R2}) = \underset{\sim}{d} \qquad (28a,b)$$

So also, by conjunction and disjunction of the output vectors

$\underset{\sim}{b1}$ and $\underset{\sim}{b2}$ of (27a,b) we obtain the outputs of (29a,b):

$$\underset{\sim}{b1} \otimes \underset{\sim}{b2} = \underset{\sim}{c}' \; , \quad \underset{\sim}{b1} \oplus \underset{\sim}{b2} = \underset{\sim}{d}' \qquad (29a,b)$$

We shall prove below that

$$\underset{\sim}{c} \neq \underset{\sim}{c}' \text{ in general, while } \underset{\sim}{d} = \underset{\sim}{d}' \text{ always} \qquad (30a,b)$$

The proof is from pure Boolean algebra and does not

involve any logical considerations. It follows from the

definition of the unary matrix product contained in Eqs.(27a,b).

Thus, with the notation adopted in this report, for a given

input a(i), we have*

$$a(i) \otimes R1(i,j) = b1(j), \quad a(i) \otimes R2(i,j) = b2(j), \quad j = 1 \text{ to } n$$
$$(31a,b)$$

---

\* In the succeeding equations of this section, the symbol $\otimes$ is
  omitted wherever it is obvious.

Then, it necessarily follows that

$$a(i) \quad \because \quad (R1(i,j) \otimes R2(i,j)) = b1(j) \otimes b2(j) \qquad (32)$$

and similarly for the Boolean sum. However, if the vector $\underset{\sim}{a}$

is not a basic vector, with $a(i) = 1$ for a single $i$ and

equal to 0 for all other $i = 1$ to $m$, but is a general vector

with more than one component, $a(i)_1 = 1$, then Eq.(32) does not

follow, and it takes the form of non-equality of (33) :

$$\bigoplus_{i=1}^{m} a(i) \, (R1(i,j) \otimes R2(i,j)) \quad \Longleftrightarrow\!\!\!\!/ \quad b1(j) \otimes b2(j), \quad j = 1 \text{ to } n$$
$$(33)$$

On the other hand, because of the associative property of the

Boolean sum, we obtain (34) in support of Eq.(30b), which can be

written in the form (35):

$$\bigoplus_{i=1}^{m} a(i) \, (R1(i,j) \oplus R2(i,j)) \quad \Longleftrightarrow \quad b1(j) \oplus b2(j), \quad j = 1 \text{ to } n$$
$$(34)$$

$$\underset{\sim}{a} \, (\underset{\sim}{R1} \oplus \underset{\sim}{R2}) = \underset{\sim}{a}\,\underset{\sim}{R1} \oplus \underset{\sim}{a}\,\underset{\sim}{R2} \qquad (35)$$

Thus, we have obtained a general proof of the result,

associated with the particular examples given in Section 5,

showing the non-equivalence for conjunction, and equivalence

for disjunction, of the two ways of performing this logical

operation with two relations applied to the same input. For

the same reasons as given above, it can be shown that (36a,b)

are valid for the Boolean truth value of two binary relations,

applied to the same input and output in two different ways, as

indicated there.

$$(\underset{\sim}{a} \, \underset{\approx}{R1} \, \underset{\sim}{b}) \, \bigvee \, (\underset{\sim}{a} \, \underset{\approx}{R2} \, \underset{\sim}{b}) \iff \underset{\sim}{a} \, (\underset{\approx}{R1} \oplus \underset{\approx}{R2}) \, \underset{\sim}{b} \qquad (36a)$$

$$(\underset{\sim}{a} \, \underset{\approx}{R1} \, \underset{\sim}{b}) \, \bigwedge \, (\underset{\sim}{a} \, \underset{\approx}{R2} \, \underset{\sim}{b}) \not\iff \underline{a1} \, (\underset{\approx}{R1} \otimes \underset{\approx}{R2}) \, \underset{\sim}{b} \qquad (36b)$$

This follows for the reason that $\bigvee$ and $\bigwedge$ in these equations

are represented by the Boolean operations $\oplus$ and $\otimes$ , so that

(36a) follows from the associative property of the Boolean sum,

while no such result follows when Boolean sum and Boolean product

occur together in an expression as in (36b). As mentioned in

Section 5, the MATLOG equations for these are quite different

and are not expected to be equal, but the equivalence holds in

the particular case of disjunction for the reason mentioned

above. The extension to SNS truth values is discussed below.


## (b) Conditions for conjunction identity

Before considering this, we shall first consider statements

(a),(b),(c) of MR-60 mentioned above, under which the identity of the

the
two ways of performing/conjunction of two relations is obtained.
change
With slight/in nomenclature, we have to show that $\underset{\sim}{c} = \underset{\sim}{c}'$

under the following three conditions.

(a) $|\underset{\sim}{R1}| \supset |\underset{\sim}{R2}|$, or $|\underset{\sim}{R1}| \subseteq |\underset{\sim}{R2}|$

(b) $|\underset{\sim}{R1}| = |\underset{\sim}{R2}^c|$, which is equivalent to $|\underset{\sim}{R1}^c| = |\underset{\sim}{R2}|$    (3 ?

(c) $|\underset{\sim}{R1}| \otimes |\underset{\sim}{R2}| = |\underset{\sim}{E}(\underset{\sim}{k}, \emptyset)|$

We shall prove the first and the third of these, while the

second condition (b) is found to be inaccurate. (This arose

because of a mistake in writing that $\underset{=}{P}^c = \underset{\equiv}{Q}$ in page 32 of

MR-60)


Condition (a)

Taking the first condition, it follows that

$R1(i,j) \implies R2(i,j)$, which means that the set of 1's in $R1(i,j)$

contains all the elements having the value 1 in the set $R2(i,j)$.

Under these conditions, the l.h.s of (33) becomes


$$\bigoplus_{i=1}^{m} a(i) \, R2(i,j) = b2(j)$$


Under the same conditions, we also obtain the result that

$b1(j) \implies b2(j)$, so that the r.h.s of (33) is also $b2(j)$

as in (37).

Similarly, for the second condition, it follows that

$R2(i,j) \implies R1(i,j)$ and $b2(j) \implies b1(j)$, so that both the

l.h.s and the r.h.s of (33) become equal to $b1(j)$, and hence

equal to one another.

## Condition (b)

On applying this condition to Eq.(32), we obtain the

result that the r.n.s is a null vector and that tnis is not

modified in any way as a consequence of the multiple sum in

the l.h.s of (33). On the other hand, in the r.h.s of (33),

$b1(j)$ is not, in general, equal to $b2^c(j)$ when the multiple

sum is applied twice over, as in (38). Thus,

$$c'(j) = b1(j) \otimes b2(j) = (\ \bigoplus_{i=1}^{m} a(i)\ R1(i,j))\ \otimes\ (\ \bigoplus_{i=1}^{m} a(i)\ R2(i,j)) \tag{38}$$

is not necessarily equal to 0 for all j, as is the case for $c(j)$.

It can in fact be verified that, even for SNS,

$$|R1| = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \quad , \quad |R2| = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} \tag{39}$$

do not lead to an identity for (33). The condition (b) is

therefore insufficient as a general rule for (33) to be an

equality. A complete output table of Problem 4A of MATLOG-1 has

been prepared[*] which shows that condition (a) is valid in a

---

[*]See Report MR—65 for details.

vast majority of cases and only six examples are not derivable

by this condition. Of these, only two possibilities, as in

(39a,b)

$$|R1| = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \quad , \quad |R2| = \begin{pmatrix} 1 & 1 \\ 0 & \end{pmatrix} \quad , \quad |R| = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \qquad (39a)$$

$$|R1| = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \quad , \quad |R2| = \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \quad , \quad |R| = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \qquad (39b)$$

satisfy the condition (c). For the remaining four cases, both

$|R1|$ and $|R2|$ have two 1's in a row and two 1's in a column,

or **vice versa**. These have logical properties of the type

"$\underline{a}$ is always T(or F)" and "$\underline{b}$ is always T(or F)", so thatt

from pure logic, if both of them are simultaneously true, we

get the result "$\underline{a}$ = T and $\underline{b}$ = T", "$\underline{a}$ = T and $\underline{b}$ = F", etc.,

whose matrices are of the form $A(k, \ell)$, $k,\ell$ = 1, 2 . So we

replace the condition (b) in (37) by the following:

(b)   Both $|R1|$ and $|R2|$ are singular matrices, one with
       two 1's in a row and the other with two 1's in a column.

With these three conditions (a), (b), (c), of (37), all the

cases where the conjunction identity is satisfied are covered,

and they can also be generalized to QLOGIC and GLOGIC as

described in the next section 9(c).

## Condition (c)

This is best proved for SNS first, and generalized to GLOGIC.

We consider first the two cases $|E(1, 1)|$ and $|E(1, 2)|$ as given

in (40):

$$|E(1, 1)| = |E| = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \quad , \quad |E(1, 2)| = |N| = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \qquad (40)$$

For each of these, there is only one combination of $|R1|$ and $|R2|$

which will lead to their Boolean product being equal to $|E|$ or $|N|$

as the case may be. We shall prove each of these cases. The

formulae to be proved are:

$$(a_1 \quad a_2) \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \otimes (a_1 \quad a_2) \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} = (a_1 \quad a_2) \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \text{ for } \underline{E} \qquad (41a)$$

and

$$(a_1 \quad a_2) \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \otimes (a_1 \quad a_2) \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} = (a_1 \quad a_2) \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \text{ for } \underline{N} \qquad (41b)$$

It is readily verified that (41a) and (41b) are satisfied as

shown in (42a,b):

$$(a_1 \quad a_1 \oplus a_2) \otimes (a_1 \oplus a_2 \quad a_2) = (a_1 \quad a_2) \qquad (42a)$$

$$(a_2 \quad a_1 \oplus a_2) \otimes (a_1 \oplus a_2 \quad a_1) = (a_2 \quad a_1) \qquad (42b)$$

Similarly, we can work out in detail that the equality holds for

all combinations $(k, \ell)$ of SNS. A table of these is, in fact,

given in MR-60, page 39, dealing with Problem 4B.

## (c) Extension to GLOGIC

The condition (a) in (37), for the equality of the two ways of performing conjunction, is followed quite generally in GLOGIC, and hence also in SNS and QL-2, to which it reduces when we put $m = n = 2$ and $m = n = 3$, respectively. Thus $\underset{\sim}{c} = \underset{\sim}{c}'$ is valid under condition (a), for a general mxn Boolean matrix.

On the other hand, condition (c) has been proved in (42) only for SNS algebra employing 2x2 matrices. We shall now indicate how this can be generalized to GLOGIC when only logical connectives of the type $\underset{\sim}{a} \, \underset{\sim}{Z}(\underset{\sim}{k}, \, \underset{\sim}{\ell}) \, \underset{\sim}{b}$ are considered (where $Z = A, O, I$ or $E$). This is possible since, for logical relations of the type $\underset{\sim}{k} \wedge \underset{\sim}{\ell}, \, \underset{\sim}{k} \vee \underset{\sim}{\ell}, \, \underset{\sim}{k} \Rightarrow \underset{\sim}{\ell}, \, \underset{\sim}{k} \Leftrightarrow \underset{\sim}{\ell}$, the relevant mxn matrix of the relation takes the form given below in (43a to d), by a suitable rearrangement of the indices $i = 1$ to $m$, and $j = 1$ to $n$.

$$\wedge \longleftrightarrow \begin{pmatrix} 1's & 0's \\ \hline 0's & 0's \end{pmatrix} \quad , \quad \vee \longleftrightarrow \begin{pmatrix} 1's & 1's \\ \hline 1's & 0's \end{pmatrix} \qquad (43a,b)$$

$$\Rightarrow \longleftrightarrow \begin{pmatrix} 1's & 0's \\ \hline 1's & 1's \end{pmatrix} \quad , \quad \Leftrightarrow \longleftrightarrow \begin{pmatrix} 1's & 0's \\ \hline 0's & 1's \end{pmatrix} \qquad (43c,d)$$

These are seen to have a close correspondence with the structure

of the 2x2 matrices in SNS, for the four connectives A, O, I, E

respectively.

From Boolean matrix algebra, it is obvious that the four

relations represented by the matrices in (43a,b,c,d) will have

properties identical with the four SNS matrices for A, O, I, E,

with the only difference that they refer to $\underset{\sim}{a}\ \underset{\sim}{Z}\ \underset{\sim}{b}$ , as

distinct from $\underline{a}\ \underline{Z}\ \underline{b}$ for SNS. In the case of 2x2 matrices,

if $a_i$ and $b_j$ are related by $Z_{ij} = 1$, then the unary relation

$\underline{a}\ \underline{Z} = \underline{b}$ leads to $b_j$ if $a_i$ is input. This same sentence can be

carried over to the general relation in GLOGIC. However, in SNS,

there can only be one $a_i$ and one $b_j$ in the four quadrants

marked in (43a,b,c,d) while there can be a larger number of them

(such as $k\ell$ , $k(n - \ell)$ etc.) in GLOGIC. A little reflection

will show that we have in this manner completely transferred the

logical properties of the relations $\underset{\sim}{k} \wedge \underset{\sim}{\ell}$ ,..., $\underset{\sim}{k} \Longleftrightarrow \underset{\sim}{\ell}$ ,

connecting the two sets represented by $\underset{\sim}{k}$ and $\underset{\sim}{\ell}$ , into the matrix

formalism₁by using this procedure.

Therefore, if we consider the condition (c) mentioned in

(37) and worked out in (43) for SNS, the whole proof from (40)

to (42) can be carried over into GLOGIC, with $\underline{E}(k, \ell)$ being

replaced by $\underset{\sim}{E}(\underset{\sim}{k}, \underset{\sim}{\ell})$, subject to the generalizations mentioned

above. Thus, we can prove, for instance, that (44) is valid

for GLOGIC.

$$(\underset{\sim}{a} \underset{\sim}{I} \underset{\sim}{b}) \wedge (\underset{\sim}{a}^c \underset{\sim}{I} \underset{\sim}{b}^c) = \underset{\sim}{a} \underset{\sim}{E} \underset{\sim}{b} = (\underset{\sim}{a} \underset{\sim}{A} \underset{\sim}{b}) \vee (\underset{\sim}{a}^c \underset{\sim}{A} \underset{\sim}{b}^c) \qquad (44)$$

This is also applicable to logical connectives of the type

involving a clausal relation, and their generalizations

considered in earlier sections, for which most of the results

worked out in SNS algebra can be taken over. On the other hand,

as already mentioned, the results proved in (31) to (36) and

for condition (a) of (37) are valid, not only for logical matrix

operators, but also for any general m x n matrix occurring for

the relation in GLOGIC.

## (d) SNS truth values

The extension of the above considerations to SNS truth

values is fairly straightforward. For this purpose, we have

only to show the identity of $\underset{=}{c}$ of (45a) and $\underset{=}{c}'$ of (45b) for

the conjunction of two relations $\underset{\sim}{R1}$ and $\underset{\sim}{R2}$ .

$$\underset{=}{c} = \underset{=}{t}(\underset{\sim}{a} \underset{\sim}{R1} \underset{\sim}{b}) \wedge \underset{=}{t}(\underset{\sim}{a} \underset{\sim}{R2} \underset{\sim}{b}) \qquad (45a)$$

$$\underset{=}{c}' = \underset{=}{t}(\underset{\sim}{a} (\underset{\sim}{R1} \otimes \underset{\sim}{R2}) \underset{\sim}{b}) \qquad (45b)$$

(using GSTV)

It can be readily shown that this equality/is equivalent to

two equations (46a) and (46b) given below.

$$\langle \underset{\sim}{a} \mid \underset{\approx}{R1} \mid \underset{\sim}{b} \rangle \otimes \langle \underset{\sim}{a} \mid \underset{\approx}{R2} \mid \underset{\sim}{b} \rangle = \langle \underset{\sim}{a} \mid \underset{\approx}{R1} \otimes \underset{\approx}{R2} \mid \underset{\sim}{b} \rangle \tag{46a}$$

$$\langle \underset{\sim}{a} \mid \underset{\approx}{R1}^{c} \mid \underset{\sim}{b} \rangle \oplus \langle \underset{\sim}{a} \mid \underset{\approx}{R2}^{c} \mid \underset{\sim}{b} \rangle = \langle \underset{\sim}{a} \mid \underset{\approx}{R1}^{c} \oplus \underset{\approx}{R2}^{c} \mid \underset{\sim}{b} \rangle \tag{46b}$$

Of these, the second equation (46b) for disjunction is universally true, following from the result (35). Therefore, the condition for the equality of the SNS truth values of (45a) and (45b) is the same as that for the equality to hold between the l.h.s and r.h.s for the Boolean truth values in (46a). This has been
                                                       (GBTV)
discussed above, and we have listed the three conditions (a), (b), (c) in Section 9(b) above. Therefore, if one of these conditions is satisfied, not only is there a conjunction identity for Boolean truth values, but also for SNS truth values defined in (45a,b).

Since we have proved this for a general relation denoted by m x n matrices, it is equally true for QL-2 and SNS logic. Therefore, all the results obtained in the problems discussed in MR-60 and 61 become explainable. The most interesting consequence from the point of view of the logic, is that the conjunction of two relations is not a uniquely definable property, and that its logical consequences will depend upon whether the conjunction is applied to the two relations as such, or to the truth values of the consequences of each of the two relations applied individually. This will be pursued further in a theoretical paper. (See MR-65)

## 10. Comparison of GSTV and GSTV2

This section is what has been referred to as Appendix

to MR-62 in the previous two parts MR-60 and 61, with special

reference to the analysis of Problems 4B, 4C of the former

and Problems 8B, 8C of the latter. The discussion has also

relevance to Problems 1A and 4A of the former and Problem 7

of the latter. They form a natural extension of the previous

section 9 of this report with special reference to the differences

between the application of GSTV and GSTV2 to calculate truth

values. We shall first discuss the essential feature with

reference to SNS logic and then indicate how these can be

extended to QLOGIC and to GLOGIC in general.


### (a) Difference between SBTV and SBTV2

Although the two functions SBTV and SBTV2 have not been

specially defined in MR-60 dealing with SNS logic, they are

particular cases of GLOGIC where GSTV and GSTV2 have been

defined. The crucial differences between the two, for relations

expressible by logical connectives in propositional calculus,

can be traced to the following example of Boolean truth values,

calculated via the matrix formalism, and via the relative

truth values, for the simple case of a disjunction $\underline{a} \lor \underline{b}$ .

For the case when $\underline{b} = (0\ 0) = X$, the following two results

(47a,b) are obvious.

$$(1\ 0) \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \end{pmatrix} = 0 \quad (\text{SBTV}) \tag{47a}$$

$$(1\ 0) \begin{pmatrix} 1 \\ 0 \end{pmatrix} \ (1\ 0) \begin{pmatrix} 0 \\ 0 \end{pmatrix} = 1 \qquad 0 = 1 \ (\text{SBTV2}) \tag{47b}$$

Thus, for the relation $\underline{a} \lor \underline{b}$ , and the inputs $\underline{a} = T$ and $\underline{b} = X$, the truth value calculated using the MATLOG formulation of BVMF is 0, while that calculated using relative truth values as in the "FORTRAN Program NYAYA2" (MR-10) is 1. Even at the time that NYAYA was developed, we had noticed that the formulae led to a non-X state for the SNS truth value of $\underline{a} \lor \underline{b}$ for the inputs $\underline{a} = T$, $\underline{b} = X$, even though one of the inputs is the impossible state which would intuitively indicate that the relation should also have the contradictory state X. This was taken care of by introducing an additional condition in all NYAYA formulae, (except union), namely "X-priority check", which made the truth value X if either $\underline{a}$ or $\underline{b} = X$ (see MR-12, p 40 and MR-16, p9, in particular). With this condition, it can be readily verified that SBTV and SBTV2 will agree for all inputs and outputs and for all matrix relations in SNS logic.

In fact, when MATLOG was first developed and the first

program for MATLOG was written in MR-18, we noticed that the

X-priority check was no longer necessary for truth values, and

therefore, this has never been mentioned explicitly thereafter,

until in MR-62 PLOGIC formulae were discussed, where, once again,

this was found to be necessary. For instance, we have put the

condition

$$\text{IF} \quad \text{QVA . OR . QVB} = \text{Q8} \; , \quad \text{QVC} = \text{Q8} \tag{48}$$

Therefore it suggests itself that, for truth values requiring

the condition that impossible as input always leads to

impossible as output (except for union), BVMF with matrix

multiplication can always be employed, including the case of

union in SUNION, GUNION and QUNION.

### (b) Extension to SSTV2, QSTV2, GSTV2

The question arises as to why SBTV2, and the derived truth

value formulae SSTV2, QSTV2 and GSTV2, are necessary. This is

best made clear by taking the general example of GBTV2

corresponding to the disjunction $\underline{Q}(\underline{k}, \underline{\ell}) = \underline{k} \vee \underline{\ell}$, which

means, in multivalued logic corresponding to sets $\mathcal{A}$ and $\mathcal{B}$

having the subsets $A_k$ and $B_\ell$, that "$A_k$ or $B_\ell$ is true".

In this case, if $B_\ell$ is a null set , since $A_k$

is not a null set, if one element of this is present, then

the given relation is true. This feature can be obtained only

by using GBTV2, and not by GBTV, which will automatically

make the output false for all inputs if $B = \emptyset = (0 \ 0 \ ... \ 0 \ 0)$.

Therefore, the converse of the X-priority check is demanded

for set theory and multivalued logic. The matrix formulation

is universally applicable so long as neither A nor B is a null set.

If one of them is a null set and the relation is a logical

disjunction of $^\mu A_k$ or $B_\ell^{\prime\prime}$ , then GBTV2 and the analog GSTV2

must be employed. In fact, even for the sets A and B having

only two members, or three members, and the vector $\underline{a}$, or $\underline{a}$, is

not a truth value, but represents a general 2-vector or 3-vector $\left(\underline{\underline{a}}\right)$

for sets, the same considerations hold.

These are not defects of BVMF, but rather very much to its

credit, as it is possible to distinguish clearly the application

of m x n matrices for all m, n $\geqslant$ 2 , to the truth values in

SNS and QL, and to set theory in general, for all applications.

It might even happen, in set theory, that the occurrence of the

null set is an impossibility, in which case, this can be

incorporated in a suitable manner in the formulae. In general,

if the conditions of the problem are precisely stated, it is

always possible to construct a vector-matrix formula to take

care of the logical features of the problem, and then we can

make it computer implementable for solution.

If these features are understood, then the solutions

to all the problems mentioned above become understandable.

Some of these are discussed below, and in MR-65, and full tables

of data are contained there.

(c) Explanation of the differences between SSTV and SSTV2 for $\underline{E}(k, \ell)$

Analogous to (47a,b) above, the two BVMF formulae for

$\underline{a} \wedge \underline{b}$ are as follows.

$$(1 \quad 0) \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \end{pmatrix} = 0 \quad (SBTV) \tag{49a}$$

$$(1 \quad 0) \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes (1 \quad 0) \begin{pmatrix} 0 \\ 0 \end{pmatrix} = 1 \otimes 0 = 0 \quad (SBTV2) \tag{49b}$$

In this case, the two agree. In fact, as a general proposition,

it can be shown that all conjunctions lead to agreement between

SBTV and SBTV2, while for disjunctions if either $\underline{a}$ or $\underline{b}$ is X,

then it is possible that SBTV and SBTV2 disagree.

Combining these two results in (47) and (49), we can

make the following general statement of the likely differences

between SSTV and SSTV2.

> "For all conjunctions $\underline{a} \wedge \underline{b}$ , if either $\underline{a}$ or $\underline{b}$
>
> is X, then it is possible that SSTV2 gives F while
>
> SSTV gives X" . (50a)

> "For all disjunctions $\underline{a} \vee \underline{b}$, if either $\underline{a}$ or $\underline{b}$ is X,
>
> it is possible that SSTV2 leads to T, while SSTV
>
> always leads to X". (50b)

This covers all of Problems 1A and 4A of Part I and Problem 7

of Part II.

In fact, the same considerations, extended to GLOGIC

(as indicated in theory in MR-65) also explain the conditions

under which SSTV and SSTV2 formulae give different results,

for $\underline{a} \, E \, \underline{b}$ in SNS, and also correspondingly for $a \, E \, b$ in QL-2,

and $\underline{a} \, E \, \underline{b}$ in GLOGIC. We shall not describe this in detail,

but give some illustrative tables which, when proved, will make it

clear. Since tables illustrating Problems 1A and 4A of Part I,

and Problem 7 of Part II, are given in those parts, and as the

complete output for Problems 4C and 8C will be given in MR-65,

we shall only give selected tables from the latter two for

illustration below. Thus, Table 5(a) contains the computer

outputs for $\underset{=}{E}(1, 1)$ and $\underset{=}{E}(1, 2)$ for SSTV2 implementation

of the mathematical formulae (51a,b), for $\underset{=}{E}(1, 1)$, and the

corresponding formula generated for $\underset{=}{E}(1, 2)$ by replacing $\underset{=}{b}$ by $\neg\underset{=}{b}$.

$$\underset{=}{a} \; \underset{=}{E} \; \underset{=}{b} \;\; = \;\; (\underset{=}{a} \; A \; \underset{=}{b}) \;\; \vee \;\; (\neg \underset{=}{a} \; A \; \neg \underset{=}{b}) \tag{51a}$$

$$\underset{=}{a} \; \underset{=}{E} \; \underset{=}{b} \;\; = \;\; (\underset{=}{a} \; \underset{=}{I} \; \underset{=}{b}) \;\; \wedge \;\; (\neg \underset{=}{a} \; \underset{=}{I} \; \neg \underset{=}{b}) \tag{51b}$$

The first two columns correspond to the inputs $\underset{=}{a}$, $\underset{=}{b}$ . The

third column is a common output of (51a) and (51b) obtained

by SSIV, while the entries in the 4th and 5th column are

respectively those obtained with (51a) and (51b) using SSTV2

implementation, which differ from SSTV. It will be noticed

that such differences occur only when either $\underset{=}{a}$ or $\underset{=}{b}$ is X

and also that the differring entry is always F for the fourth

column corresponding to (51a) and T for the fifth column

corresponding to (51b). This feature is true for all $\underset{=}{E}(k, \ell)$,

as will be seen from the other examples given in Tables 5(b,c).

A complete output for this problem is included in MR-65.

## Table 5.  Check of Problem 4C for  E(1, 1) and E(1, 2)

### (a) Both k and $\ell$ not 3, or 4

PROB  4C    (K,L) = (1,1)

| | | | | |
|---|---|---|---|---|
| T | T | T | | |
| T | F | F | | |
| T | D | D | | |
| T | X | X | | |
| F | T | F | | |
| F | F | T | | |
| F | D | D | | |
| F | X | X | | |
| D | T | D | | |
| D | F | D | | |
| D | D | D | | |
| D | X | X | F | T |
| X | T | X | | |
| X | F | X | | |
| X | D | X | F | T |
| X | X | X | | |

PROB  4C    (K,L) = (1,2)

| | | | | |
|---|---|---|---|---|
| T | T | F | | |
| T | F | T | | |
| T | D | D | | |
| T | X | X | | |
| F | T | T | | |
| F | F | F | | |
| F | D | D | | |
| F | X | X | | |
| D | T | D | | |
| D | F | D | | |
| D | D | D | | |
| D | X | X | F | T |
| X | T | X | | |
| X | F | X | | |
| X | D | X | F | T |
| X | X | X | | |

Table 5. Contd.

(b)    One of  k, $\ell$ = 3 or 4

```
        PROB 4C   (K,L) = (1,3)
T   T     T
T   F     T
T   D     T
T   X     X
F   T     F
F   F     F
F   D     F
F   X     X
D   T     D
D   F     D
D   D     D
D   X     X    F   T
X   T     X
X   F     X-
X   D     X
X   X     X
```

```
        PROB 4C   (K,L) = (1,4)
T   T     F
T   F     F
T   D     F
T   X     X
F   T     T
F   F     T
F   D     T
F   X     X
D   T     D
D   F     D
D   D     D
D   X     X    F   T
X   T     X
X   F     X
X   D     X
X   X     X
```

(c)  Both k and $\ell$ = 3, or 4

```
        PROB 4C    (K,L) = (3,3)
T   T        T
T   F        T
T   D        T
T   X        X
F   T        T
F   F        T
F   D        T
F   X        X
D   T        T
D   F        T
D   D        F
D   X        X
X   T        X
X   F        X
X   D        X
X   X        X


        PROB 4C    (K,L) = (3,4)
T   T        F
T   F        F
T   D        F
T   X        X
F   T        F
F   F        F
F   D        F
F   X        X
D   T        F
D   F        F
D   D        F
D   X        X
X   T        X
X   F        X
X   D        X
X   X        X
```

Table 6 contains the computer outputs for $\underset{\sim}{E}(1, 1)$, $\underset{\sim}{E}(1, 6)$ $\underset{\sim}{E}(1, 7)$ and $\underset{\sim}{E}(7, 7)$ for the QSTV2 implementation of QL-2 formulae analogous to (51a) and (51b). In this case, the full 8x8 table is given in a square array for QSTV (first table) and also for both (51a) and (51b) employing QSTV2 (second and third tables). Of these, the first table corresponds to the agreeing results using both (51a) and (51b), and also matrix multiplication, for QSTV It will be seen that the second and third tables for each $(k, \ell)$ agree with the first table for the first seven columns and first seven rows, in all the examples given. However, for the last row and/or last column, QSTV gives always X as output, while QSTV2 gives some entries, which are F in the case corresponding to (51a), and T in the case corresponding to (51b). The behaviour is closely similar to that observed in SNS. It is only necessary to replace the contradictory state, X, for the SNS terms, by the contradictory state, $\emptyset$, for quantifier states.

It will be noticed that there are 2x3 = 6 entries when both $\underset{\sim}{k}$ and $\ell$ are not equal to 7 or 8, three entries if one of them equal to 7 or 8 while the other is not , and no differences at all between QSTV and QSTV2 if both are equal to 7 or 8.

Table 6. Check of equations similar to (51a,b) for $E(k,\ell)$

FROB 8C (k,L) = (1,1) — 6 differences

|      | ALL | NAL | SOM | AON | NEX | EXS | IND | IMP | ALL | NAL | SOM | AON | NEX | EXS | IND | IMP | ALL | NAL | SOM | AON | NEX | EXS | IND | IMP |
|------|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| ALL  | T   | F   | F   | D   | F   | D   | D   | X   | T   | F   | F   | D   | F   | D   | D   | X   | T   | F   | F   | D   | F   | D   | D   | X   |
| NAL  | F   | T   | T   | D   | T   | D   | D   | X   | F   | F   | F   | D   | F   | D   | D   | X   | F   | T   | T   | D   | T   | D   | D   | X   |
| SOM  | F   | T   | T   | D   | T   | D   | D   | X   | F   | F   | T   | D   | F   | D   | D   | X   | F   | T   | T   | D   | T   | D   | D   | X   |
| AON  | D   | D   | D   | D   | D   | D   | D   | ⓕ   | D   | D   | D   | D   | D   | D   | D   | ⓣ   | D   | D   | D   | D   | D   | D   | D   | ⓣ   |
| NEX  | F   | T   | T   | D   | T   | D   | D   | X   | F   | T   | T   | D   | T   | D   | D   | X   | F   | T   | T   | D   | T   | D   | D   | ⓣ   |
| EXS  | D   | D   | D   | D   | D   | D   | D   | ⓕ   | D   | D   | D   | D   | D   | D   | D   | ⓣ   | D   | D   | D   | D   | D   | D   | D   | ⓣ   |
| IND  | D   | D   | D   | D   | D   | D   | D   | ⓕ   | D   | D   | D   | D   | D   | D   | D   | ⓣ   | D   | D   | D   | D   | D   | D   | D   | ⓣ   |
| IMP  | X   | X   | X   | ⓕ   | X   | ⓕ   | ⓕ   | X   | X   | X   | X   | ⓣ   | X   | ⓣ   | ⓣ   | X   | X   | X   | X   | ⓣ   | X   | ⓣ   | ⓣ   | X   |

FROB 8C (k,L) = (1,4) — 6 differences

|      | ALL | NAL | SOM | AON | NEX | EXS | IND | IMP | ALL | NAL | SOM | AON | NEX | EXS | IND | IMP | ALL | NAL | SOM | AON | NEX | EXS | IND | IMP |
|------|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| ALL  | T   | D   | T   | D   | F   | D   | D   | X   | T   | F   | F   | D   | F   | D   | D   | X   | T   | F   | F   | D   | F   | D   | D   | X   |
| NAL  | F   | D   | F   | D   | F   | D   | D   | X   | F   | F   | F   | D   | F   | D   | D   | X   | F   | F   | F   | D   | F   | D   | D   | X   |
| SOM  | F   | D   | F   | D   | T   | D   | D   | X   | F   | F   | F   | D   | T   | D   | D   | X   | F   | F   | F   | D   | T   | D   | D   | X   |
| AON  | D   | D   | D   | D   | D   | D   | D   | ⓕ   | D   | D   | D   | D   | D   | D   | D   | ⓤ   | D   | D   | D   | D   | D   | D   | D   | ⓣ   |
| NEX  | F   | D   | F   | D   | T   | D   | D   | X   | F   | F   | F   | D   | T   | D   | D   | X   | F   | F   | F   | D   | T   | D   | D   | X   |
| EXS  | D   | D   | D   | D   | D   | D   | D   | X   | D   | D   | D   | D   | D   | D   | D   | X   | D   | D   | D   | D   | D   | D   | D   | ⓣ   |
| IND  | D   | D   | D   | D   | D   | D   | D   | X   | D   | D   | D   | D   | D   | D   | D   | ⓤ   | D   | D   | D   | D   | D   | D   | D   | ⓣ   |
| IMP  | X   | X   | X   | X   | X   | X   | X   | X   | X   | X   | X   | ⓤ   | X   | ⓤ   | ⓤ   | X   | X   | X   | X   | ⓣ   | X   | ⓣ   | ⓣ   | X   |

PROB 8C (A,L) = (1,7) — 3 differences

ALL NAL SOM AON NEX EXS IND IMP

ALL NAL SOM AON NEX EXS IND I

|     | ALL | NAL | SOM | AON | NEX | EXS | IND | IMP |
|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| ALL | T | T | F | F | D | F | D | D | X |
| NAL | F | F | F | D | F | D | D | X |
| SOM | F | F | F | D | F | D | D | X |
| AON | D | D | D | D | D | D | D | X |
| NEX | F | F | F | D | F | D | D | X |
| EXS | D | D | D | D | D | D | D | X |
| IND | D | D | D | D | D | D | D | X |
| IMP | X | X | X | X | X | X | X | X |

PROB 8C (A,L) = (7,7) — No differences

ALL NAL SOM AON NEX EXS IND IMP

ALL NAL SOM AON NEX EXS IND IMP

|     | ALL | NAL | SOM | AON | NEX | EXS | IND | IMP |
|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| L | T | T | T | T | T | T | T | X |
| L | T | T | T | T | T | T | T | X |
| M | T | T | T | T | T | T | T | X |
| M | T | T | T | T | T | T | T | X |
| X | T | T | T | T | T | T | T | X |
| S | T | T | T | T | T | T | T | X |
| D | T | T | T | T | T | T | T | X |
| P | X | X | X | X | X | X | X | X |

This is closely similar to the behaviours in SNS as mentioned

for Table 5. In fact, it is readily verified that the

differences between QSTV and QSTV2 are produced only if

$$\underline{t}(\underline{a} / \underline{k}) = D \quad \text{and} \quad \underline{b} = \emptyset',$$

or

$$\underline{t}(\underline{b} / \underline{\ell}) = D \quad \text{and} \quad \underline{a} = \emptyset \qquad (52)$$

The corresponding formulae for SNS can be put in the same

pattern as

$$\underline{t}(\underline{a} / s(k)) = D \quad \text{and} \quad \underline{b} = X$$

or

$$\underline{t}(\underline{b} / s(k)) = D \quad \text{and} \quad \underline{a} = X \qquad (53)$$

These results are readily derivable from BVMF algebra, and this

is left to the reader who may prove them even more generally for

GLOGIC, after reading MR-65.

## Appendix

The appendix to Part III, MR-62 has been referred to in MR-60 and 61. These references deal essentially with Problems 4, 6, 7 and 8. In some cases the reference is made to MR-65 for the same problems. Since the theoretical discussions have to be given in some detail, no appendix is added here, but both the materials intended for this appendix as well as for MR-65 are given in some detail in the report MR-65 forming Part IV of MATLOG series.

The proposed detailed treatment of QL-1B ,mentioned in page 62, MR-61, and of tensor relations and GLOGIC, mentioned in MR-62 at various places, have not been written up in MR-65 for lack of technical assistance.

# STEREOCHEMISTRY OF COLLAGEN

(Invited Review Article to appear in
Intl. J. Peptide and Protein Res.)

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# Stereochemistry of Collagen

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## ABSTRACT

This review article, based on a lecture delivered in Madras in 1985, is an account of the author's experience in the working out of the molecular structure and conformation of the collagen triple-helix over the years 1952 - 1978. It starts with the first proposal of the correct triple-helix in 1954, but with three residues per turn, which was later refined in 1955 into a coiled-coil structure with approximately 3.3 residues per turn. The structure readily fitted proline and hydroxyproline residues and demanded glycine as every third residue in each of the three chains. The controversy regarding the number of hydrogen bonds per

tripeptide could not be resolved by x-ray diffraction or

energy minimization, but physicochemical data, obtained in

other laboratories during 1961-1965, strongly pointed to two

hydrogen bonds, as suggested by the author. However, it was

felt that the structure with one straight NH ... O bond was

better. A reconciliation of the two was obtained in Chicago

in 1968, by showing that the second hydrogen bond is _via_ a

water molecule, which makes it weaker, as found in the

physicochemical studies mentioned above. This water molecule

was also shown, in 1973, to take part in further cross-linking

hydrogen bonds with the OH group of hydroxyproline, which

occurred always in the location previous to glycine, and is

at the right distance from the water. Thus, almost all

features of the primary structure, x-ray pattern, optical

and hydrodynamic data, and the role of hydroxyproline in

stabilising the triple helical structure, have all been

satisfactorily accounted for. These also lead to a confirmation

of Pauling's theory that vitamin C improves immunity to

deseases, as explained in the last section.

# Stereochemistry of Collagen

## 1. Introduction

It was very kind of Dr. Thyagarajan, Director of CLRI, to have invited me to inaugurate the renewed activities in the field of basic sciences in this Institute. I am particularly happy about this since I have been associated with CLRI right from its inception in the 1950's. I could not think of a better topic than the 'Stereochemistry of Collagen' for the inaugural address, since it has been one of the main fields of activity, in our laboratories of Molecular Biophysics, both in the University of Madras, and in the Indian Institute of Science, Bangalore. As I have not been doing original work in this field for the past few years, I propose to give a short account of the various developments in this subject made in our laboratories, with special reference to the molecular structure of collagen. Starting from x-ray diffraction studies in the early 1950's and the proposal of the correct triple helix for collagen in 1954-55, which were followed up by conformational analysis and energy minimization for its refinement, they have led to the delineation of the role of hydroxyproline in collagen in

the early 1970's and the linking up of this to the role

of vitamin C in immune response in 1978. These are broadly

reviewed in this lecture.

## 2. Protein helices

Stable structures of protein molecules depend upon two

aspects, namely the unbranched peptide chain of the protein

molecule and the occurrence of hydrogen bonds between

different parts of the protein chain that stabilizes its

secondary structure. The importance of hydrogen bonds for

helix formation became established with the theoretical

elucidation of the now-well-known $\alpha$-helix by Pauling in

1951. Once the parameters of the helix were found from

stereochemistry, it took no time at all to prove the

correctness of these by x-ray diffraction, as was done by

Perutz. Within a year or two of the demonstration of the

$\alpha$-helix in the KMEF group of fibrous proteins, came the

even more exciting double helix for DNA as proposed by

Watson and Crick in 1952. The helix era had begun.

I mention all these because I was encouraged to enter

this field by reading the beautiful series of papers

published by Pauling and coworkers in 1951, and, when I was

appointed Professor and Head of the newly started Physics

Department in Madras in 1952, I chose x-ray diffraction, and

x-ray crystallography in particular, as the main theme of

our laboratory, and their application to biomolecules as the

main aspect of this field that is to be pursued vigorously.

However, I did not know where to begin, and which molecules,

or biopolymers, were the most profitable ones to study.

This problem was resolved by the happy coincidence of a

visit by Prof. J.D. Bernal to Madras in the early 50's.

When I put this question to him, he told me that he was not

very happy with the various structures of collagen that

had been proposed in the literature at that time, and that

the problem was wide open. Even more than that, he indicated

that there were some specimens of shark fin ray collagen

(elastoidin) in the Department of Biochemistry of Madras

itself.

3. First contacts with collegen via CLRI

This helped us in starting our studies. But this

material was not the best one suited for x-ray diffraction

purposes. We wanted kangaroo tail tendon (KTT) or beef

Achilles tendon. Here again, it was very lucky that the

Central Leather Research Institute was in Madras; in fact

it was our neighbour in Adyar. I contacted Dr. Nayudamma

in CLRI and they obtained for us a big tubeful of KTT from

Australia. For beef tendon, the usual procedures of obtaining

the pure protein was done in CLRI itself. The availabality

of x-ray diffraction photographs in our own laboratory

for first hand studies was very useful in the solution of

the correct structure, as we shall see later.

Coming to the story of the progressive steps in obtaining

the detailed molecular architecture (secondary structure) of

the collagen helix, I shall only touch upon the highlights,

particularly with reference to the work done in our laboratories.

This report is not intended to be a review, or an authoritative

account, but only a reminiscence of the studies made in Madras ,

Chicago and Bangalore.*

*Reference may be made to the reviews (1-4) by the author,
and, therefore, this lecture will not cite full references
to the literature, but only the vital ones.

The most important feature of collagen, that has relevance to the molecular architecture of its protofibril, is the characteristic amino acid composition of the protein,which had been determined at that time for a wide range of biological sources (Fig. 1). It is widely different from that of other proteins. The most important feature is the existence of one-third the number of amino-acid residues as glycine, and the existence of almost another one-third of the residues as the imino-acid residues proline and hydroxyproline. The positively and negatively charged amino acids are few, but occur in approximately equal amounts. The other features that had been reported in the literature were the data on infra red dichroism of the fibre and optical rotation of the solution. The infra red data showed that the groups N — H and C = O, of the backbone of the protein chain, are approximately at right angles to the fibre axis, and the specific rotation indicated that the helix was left handed, having a sense apposite to that of the $\alpha$ -helix. The infra red data also indicated that all the peptide units were trans, and that no cis residues were present. The x-ray pattern did not have too much of data, and had been interpreted diffierently by different workers.
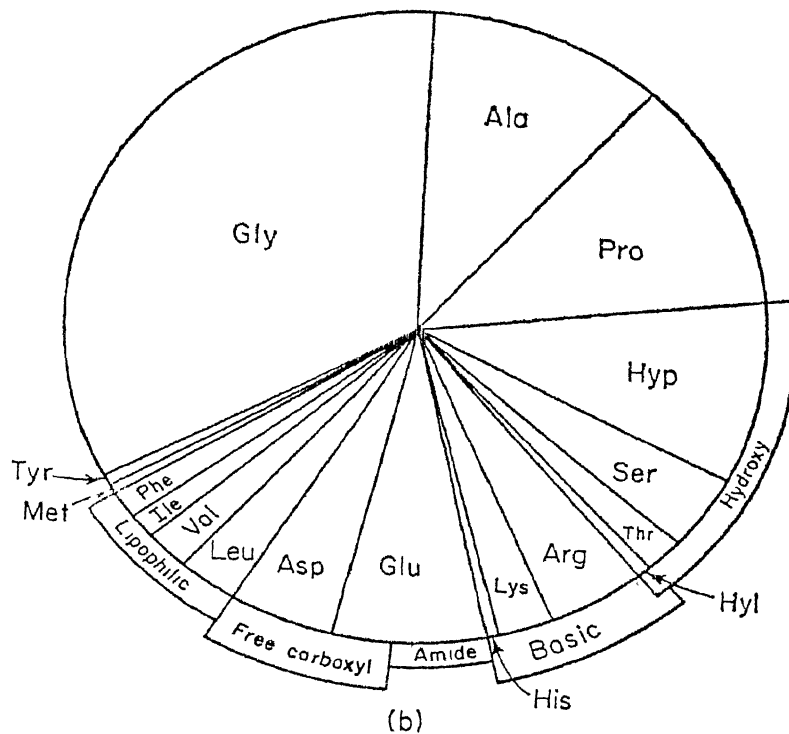
Fig.1. Amino acid composition of collagen. Note the
occurrence of 1/3 the number of residues as
glycine and practically another 1/3 as the
imino-acid residues proline and hydroxyproline,
having rigid side chains which impose a left-handed
twist for the helix (see below). (From the book (1)).

for the chain configuration in collager

There were several proposals in the literature, / some

of which are shown in Fig. 2(a-d). As will be seen from

these figures, the nature of the helix proposed by different

workers widely differred from one another, , and the one

shown in Fig.2(c) was not helical at all. Some had cis

peptide units in addition to the trans peptide units, a

feature that was contrary to the infra red information

that only trans residues were present. It must be

mentioned that there was no special feature in any of

them, which required that one-third of the residues must

be glycine.

## 4. Proposal of the first structure in 1954

I had the good fortune of having a trained crystallographer

and an excellent scientist in Dr. Gopinath Kartha, who joined

me as a postdoctoral worker, in our studies on collagen. This

gave a great spurt to our activities. Kartha and I made

various attempts at building a structure that will fit the

x-ray diffraction data, in addition to all the chemical and

physico-chemical data mentioned above. A fresh indexing of

the x-ray pattern (see Fig. 5), indicated that the protofibril

was about 12Å in diameter, and that it had a repeat along the
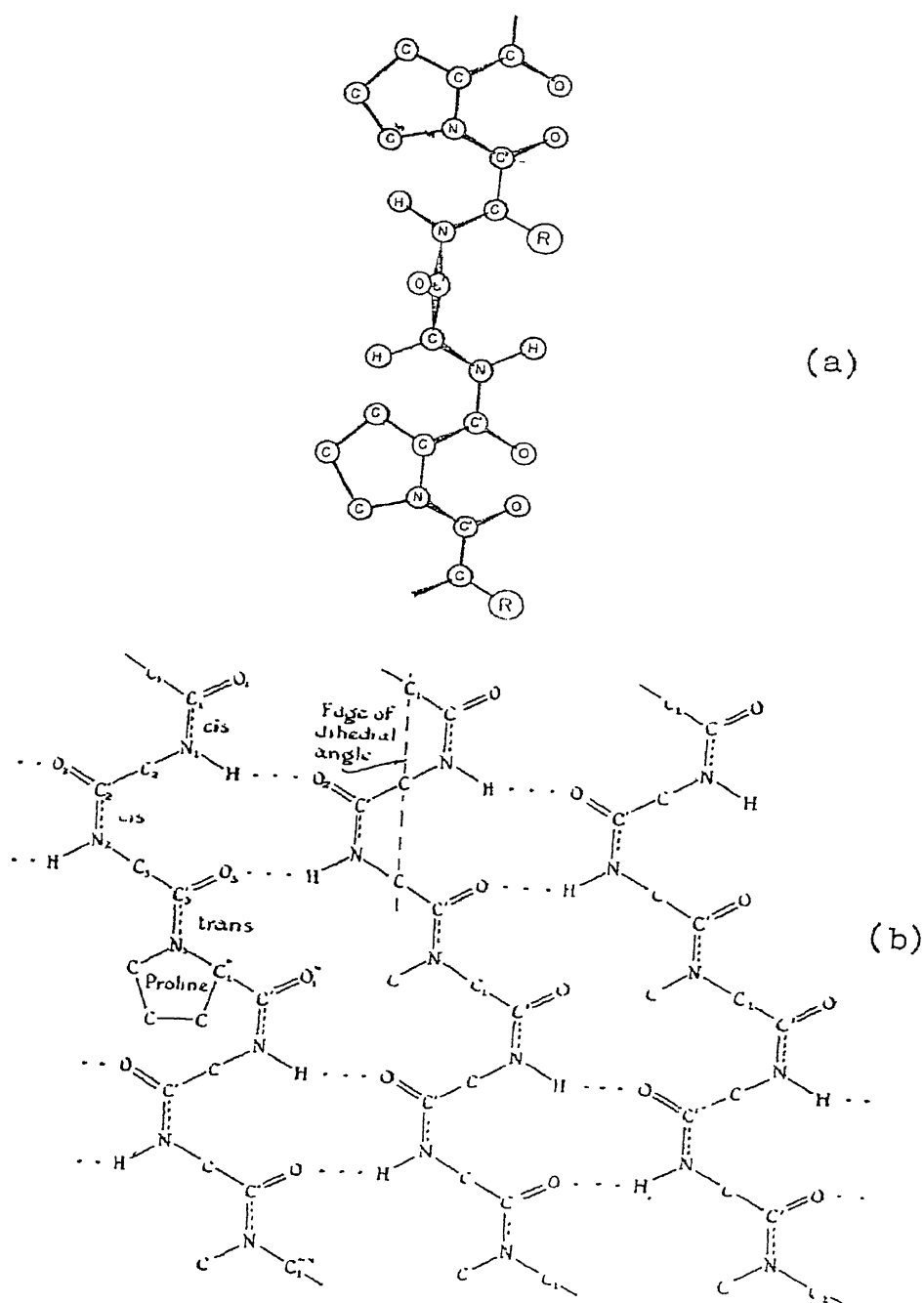
(a)

(b)

**Fig.2.** Diagrams illustrating the structures proposed for collagen by (a) Astbury and Reid (1940), (b) Pauling and Corey (1951), (c) Randal and coworkers (19 2), and (d) Crick (1954). Note that the structures in (a), (b) and (d) employ cis residues for proline, while (c), which is all trans, is not helical at all. None of them demand that glycine should form one-third the number of residues. (From (2)).
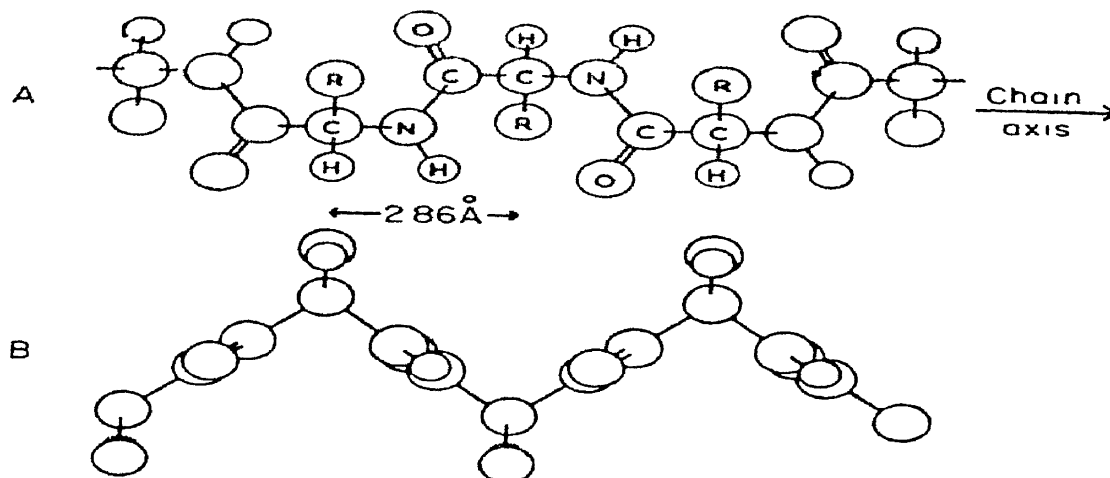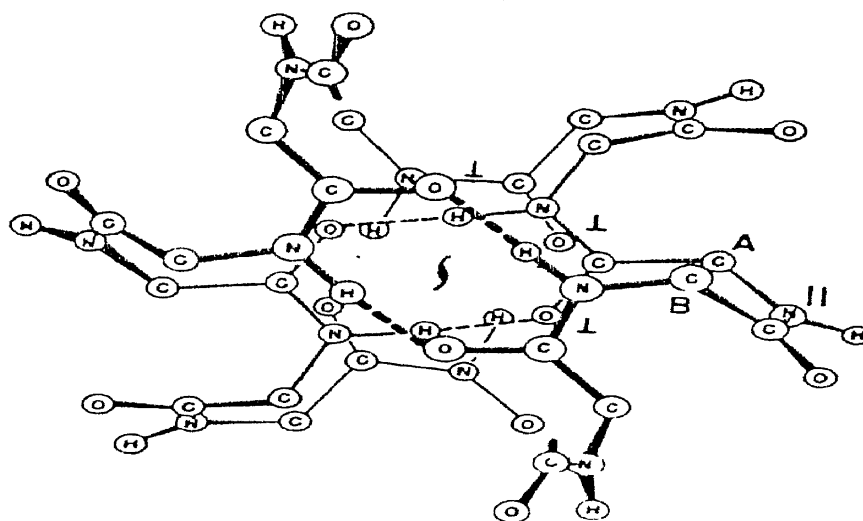
FIGURE 9(c)



FIGURE 9(d)

Fig.2. Continued

fibre axis of approximately 9-10Å with probably three residues

in the repeat, corresponding to the meridional spacing of 3Å,

and having three such chains in the unit cell as indicated

by the density. Kartha and I, being both crystallographers,

gave the greatest value to the diffraction pattern and started

building structures which were left-handed helices, as required

by optical rotation, and which had N—H and C = O bonds

at right angles to the helix axis, as required by infra red

data. We also put in the condition that no cis residues were

present, as had been found also from infra red data. All of

these were conditions that were happily chosen, and found

to be even in the final model.

But the most interesting feature was that the occurrence

of proline and hydroxyproline in the L-configuration in

collagen restricts the freedom of orientation of the chain in

relation to the fibre axis and permits only a left-handed helix.

We, therefore, assumed that each individual chain had a helical

configuration of three residues per turn, with the crystallographic

symmetry $3_2$. When three such chains, each having the

left-handed 3-fold screw axis symmetry, were put together,

and hydrogen bonds were fitted in, we found that each of them

was again related to the other two by a similar 3-fold screw

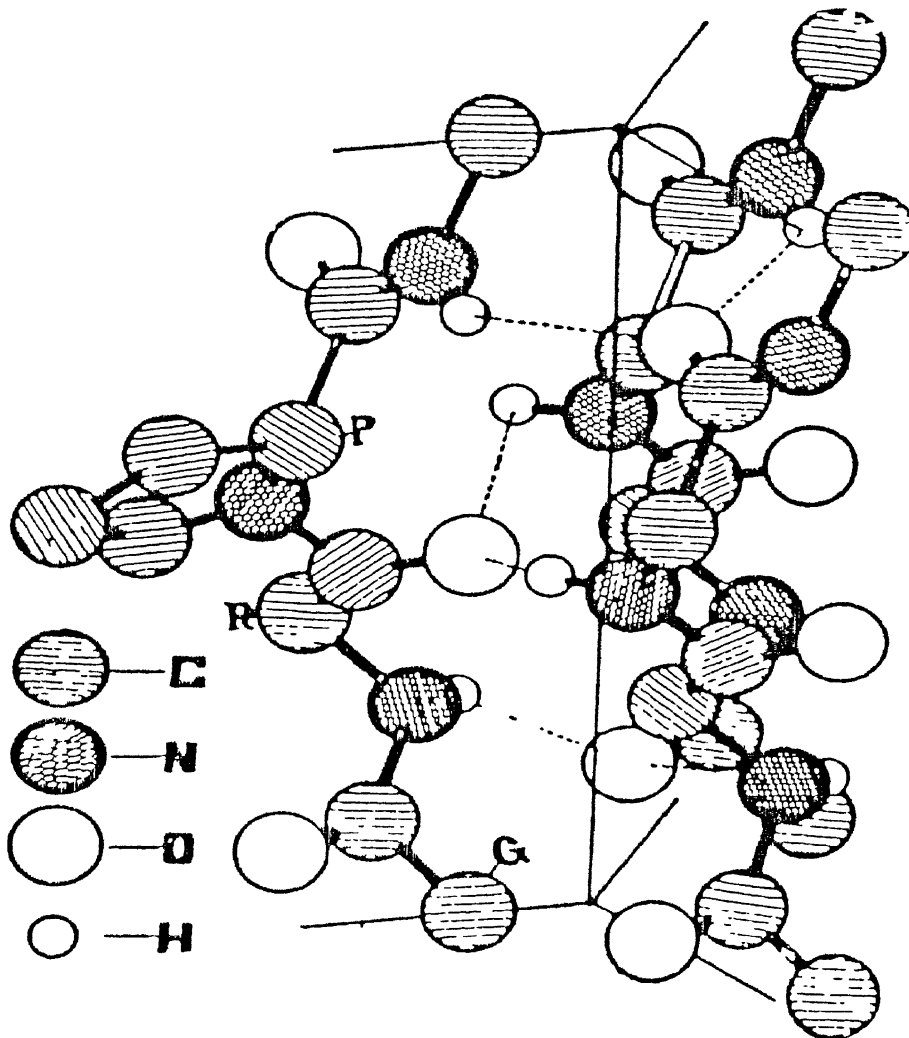axis. A perspective view of the structure so obtained,

Fig.3. Perspective view, from the side, of the contents
of one unit cell, for a height of 9— 10Å, of the
collagen triple helical structure (5). Note
that the hydrogen bonds are nearly at right
angles to the helical axis, all residues are
trans, and that L-proline side chains readily
fit in the structure, giving it a left-handed
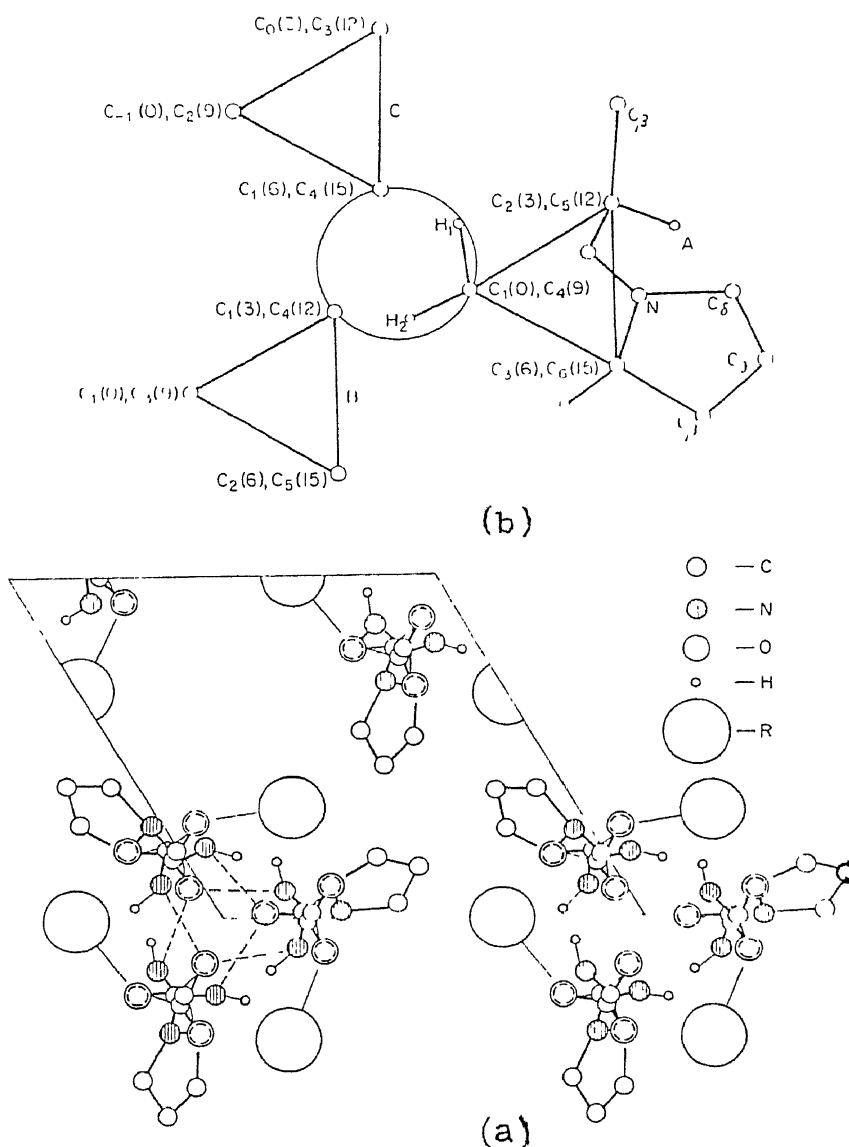twist for the helix. (From (5)).

C₀(2),C₃(12)(?)

C₋₁(0),C₂(9)

C₁(6),C₄(15)

C₁(3),C₄(12)

C₂(6),C₅(15)

C₂(3),C₅(12)

C₁(0),C₄(9)

C₃(6),C₆(15)

H₁

H₂

N

A

C₃

C₆

C₇

(b)

O — C
⊙ — N
◯ — O
∘ — H
◯ — R

(a)

Fig.4. The structure in Fig.3 projected down the helical axis — (a) triple helices packed in an hexagonal unit cell. Only the hydrogens in N — H groups are shown. Note the feature that one of the three $\alpha$-carbon atoms in the three residues occurring for one full turn cannot have any side-chain $\beta$-carbon atom, while the other two can have any side chain attached to them, including especially proline. This feature is particularly clear from (b) where the two hydrogen atoms of glycine attached to $C^{\alpha}$ , and the $C^{\beta}$-atoms of the side chain in the other two residues, are also marked in one of the chains. (From (6)).

(which we published in Nature (5) in 1954) is shown in

Fig. 3. Fig. 4(a) contains a projection down the helical

axis of the proposed crystal structure (reproduced from (6)).

As will be seen from Fig.3, the hydrogen bonds are

nearly at right angles to the helical axis and the peptide

units are all in the trans configuration. Also the proline

side chain comes on the outside of the triple-chain

protofibril. This is also seen very clearly in Fig. 4(a).

However the most significant feature of the molecular

structure that automatically came out of the analysis was

that every third residue in each of the three chains must

be glycine. The essence of this is shown in Fig. 4(b), where

it will be seen that every third residue cannot accommodate

the $\beta$ -carbon of any of the other nineteen amino acids

if inter-chain hydrogen bonds are to be made in the
                as in Fig.4(a).
protofibril,/ Therefore, the experimentally observed composition

of one-third glycine, which is a very characteristic feature
                                        feature
of collagen, comes out as a fundamental  / required for the

molecular fit of the three helical chains with one another in

the triple helix. It is also seen from Fig. 4(a), that

proline or hydroproline could redily occur in the other two

locations between two glycine residues, in the sequence

—Gly—X—Y—, of the molecular chain.

It is interesting that the exact 3-fold screw symmetry, of the chain invoked here, (but modified to a non-integral screw axis with 10 residues in 3 turns later) was found very soon to be valid for the polypeptides poly-L-proline (7) and poly glycine (8), and later / for poly-L-hydroxyproline also, all of which have a unit height of about 3Å .

(Sasisekharan, 1959),

5. Supercoiled structure with $3\frac{1}{3}$ units per repeat (9,10)

---

On careful examination, we were not fully satisfied with this structure, although it was apparently a good first approximation to the correct one. The most important feature that had to be refined was the number of residues per turn in the helix. Fortunately, just two years earlier, Cochran, Crick and Vand (1952) had published a thorough analysis of the theoretically expected diffraction pattern of helical structures, and how the geometrical features relating to the structure, such as unit twist, unit height, and pitch, could be obtained from a measurement of the diffraction pattern. Therefore G.K.Ambadi of our department took a diffraction

Fig.5.  X-ray diffraction pattern of collagen obtained by
keeping the fibre at a suitable angle to the x-ray
beam for making the 3Å spot meridional.  Apart from
substantiating the approximate value 10/3 for the
number of residues per turn (n), it also helped in
getting the  more accurate value of 3.27 for this
quantity obtained later in 1960.

picture of KTT, kept wet and fully stretched, and inclined

to the x-ray beam, so as to record all the reflections down

to 1.5Å. When we analysed this pattern (Fig. 5) according

to the formulae of Cochran et al , we found/that the correct
(9)

value, for the number of residues per turn in the collagen

helix, is close to 3.3, or 10/3 (the nearest ratio of small

integers), rather than 3. Therefore, the individual chains

cannot be helices with an integral number of units per turn.

The schematic picture of one such helix, having about 3.3

residues per turn, is shown in Fig. 6(a). In this, $P_1$, $P_2$, $P_3$

etc., represent the $C^\alpha$ atoms of adjacent peptide units and

these are separated by a twist per residue of 110° about the

helical axis, and a displacement of 3Å at right angles to the

diagram. The number of residues per turn (n) is therefore

360°/110° = 3.27. This value, rather than the value of

10/3 = 3.33, is chosen for Fig. 6 because it is based on/more
a

accurate determination made in 1960. If n = 10/3, then the

unit twist will be 108° instead of 110°; both of these are

to be compared to the value of 120° corresponding to n = 3
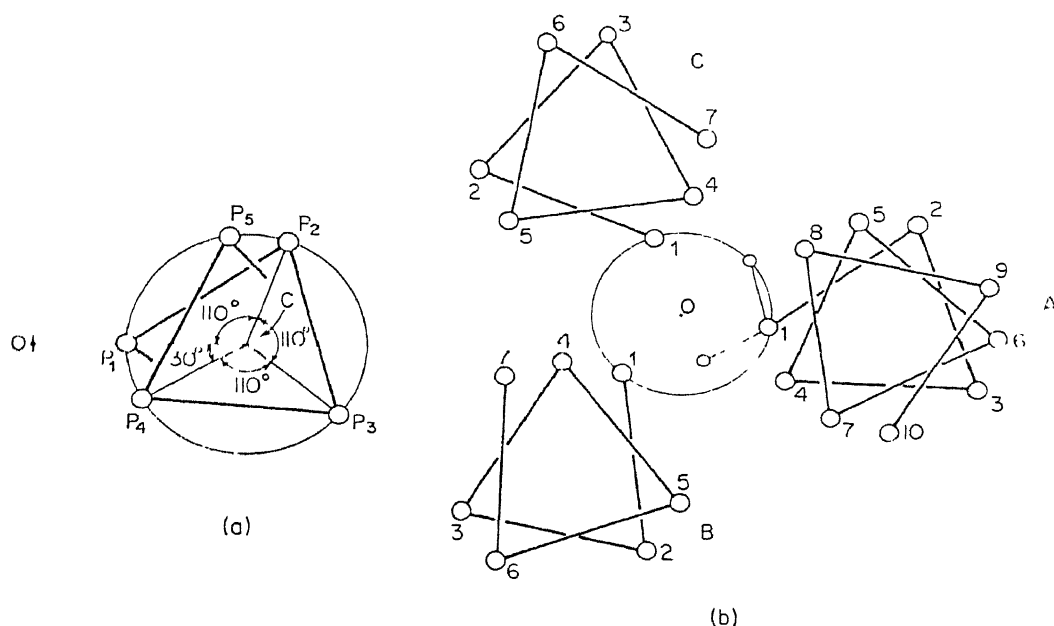
as in Fig. 4(b).

(a)

(b)

Fig.6(a) Projection of the backbone of a non-integral helix
with approximately ten residues in three turns.
(This value corresponds to 108° for the unit twist.
However, the more accurate value of 110° corresponding
to  n = 3.27, which was found later in 1960 in our
laboratory, is used in this figure.)

(b) When three such helices are put together, as in
Fig.4(b), it is seen that the condition that
every third  ∝-carbon atom must be on the inside
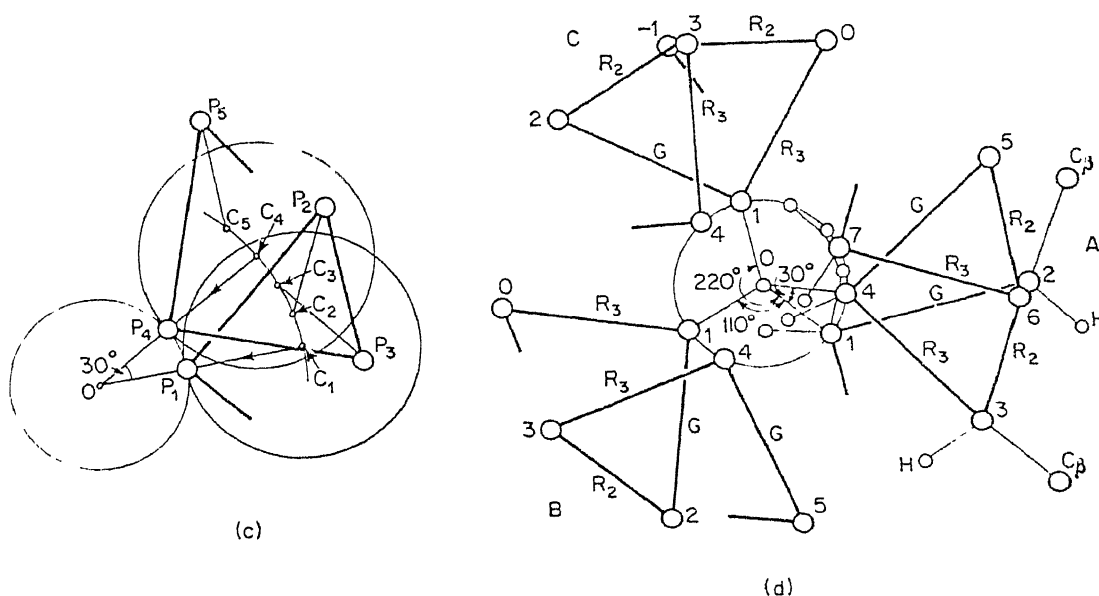of the triple helix, is lost — for example, for the
∝-carbon atoms  4, 7 and 10.

(c)

(d)

Fig.6(c)  This condition can be restored by giving a right-handed
superhelical twist of 30° (360° - 3 x 110°) for the
individual helices as shown in this figure.

(d)  The pattern of the molecular arrangement employing
three such supercoiled left-handed minor helices, all
of which are wound about the central axis with a
right-handed twist. The value of 30° for three
residues of the supercoiling twist will lead to a
major helix pitch of about 108Å, corresponding to
36 residues. (From (1)).

However, when three such helices are put together in

the pattern of Fig. 4(b), as shown in Fig. 6(b), the condition

that every third $\alpha$-carbon atom mut belong to glycine, is lost.

It took Kartha and me some time to figure out how to overcome

this difficulty. Finally, we discovered a way, and this was

to make the three helices intertwine themselves further by

twisting around the common central axis of the structure,

so as to form what is called a "coiled-coil structure".

This feature is shown for a single chain in Fig. 6(c). A

glycine $\alpha$-carbon atom $P_1$, and the next such atom $P_4$, both

now occur at the same distance from the centre O and on the

inside of the major helix. However, they are not one above

the other as in Fig. 4(b), but are relatively displaced by

a right-handed twist of 30°. The origin of this value, 30°,

for the superhelical twist will become clear from Fig. 6(d),

where three left-handed helical chains are put together and

wound about the central axis. In this case, the atom $C_1^\alpha$

of one chain (A), and that of the next one (B) to the left

of it, are separated by 110° about the major helical axis,

and this gives the value of the number of units per turn of

3.27 as required by x-ray data. If we take the $C_1^\alpha$ atoms of

of the chains A, B and C, they are at heights of 0 Å for A,

about 3Å for B and about 6Å for  C.  Then, we come back to

the corresponding  $C^{\alpha}$ atom,  $C_4^{\alpha}$ , of chain  A  at about 9Å .

We have effectively made a total left-handed twist of 3 x 110°,

and therefore we are left with a right-handed twist of 30°

between $C_1^{\alpha}$ and $C_4^{\alpha}$ of the same chain A.

It is obvious from symmetry considerations that this

pattern can be repeated in chains B and C also and the

resultant structure then has a non-crystallographic helical

true symmetry with a unit twist of $-110°$ (the minum refers

to left-handed), about the common central axis, and has a

repeating unit consisting of three residues (—Gly—X—Y—)

in each of the chains.  This is shown in Fig. 6(d).  Consequently

each peptide chain has 36 residues in the pitch of the triple

helix.  (This will be 30 if n = $3\frac{1}{3}$ .)  A side view of the

molecular structure so obtained is shown in Fig. 7.

This structure was also published in Nature (10) in 1955,

and the projection down the fibre axis of this structure is

shown in Fig. 8(a,b).  It will be seen from this that all the

stereochemical properties of the non-coiled-coil structure —

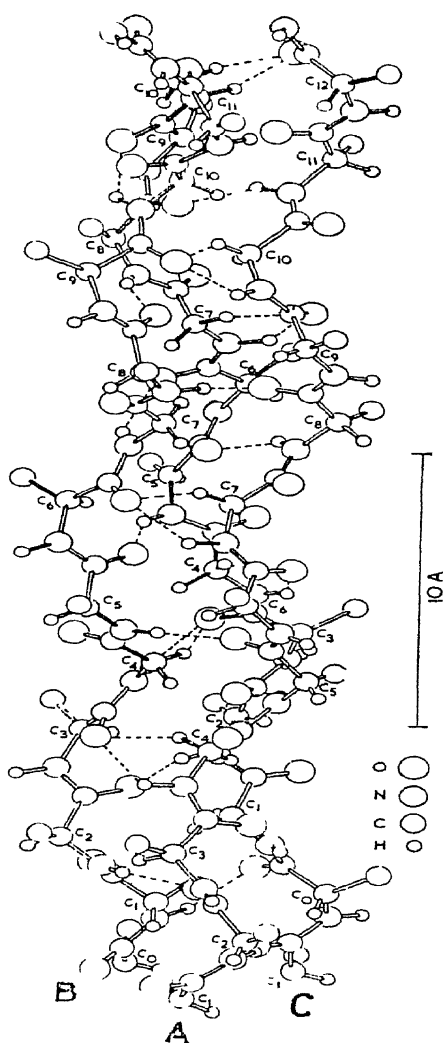namely that the hydrogen bonds are nearly at right angles to

Fig.7. Perspective view of the supercoiled collagen
structure for a height of about 30Å, showing
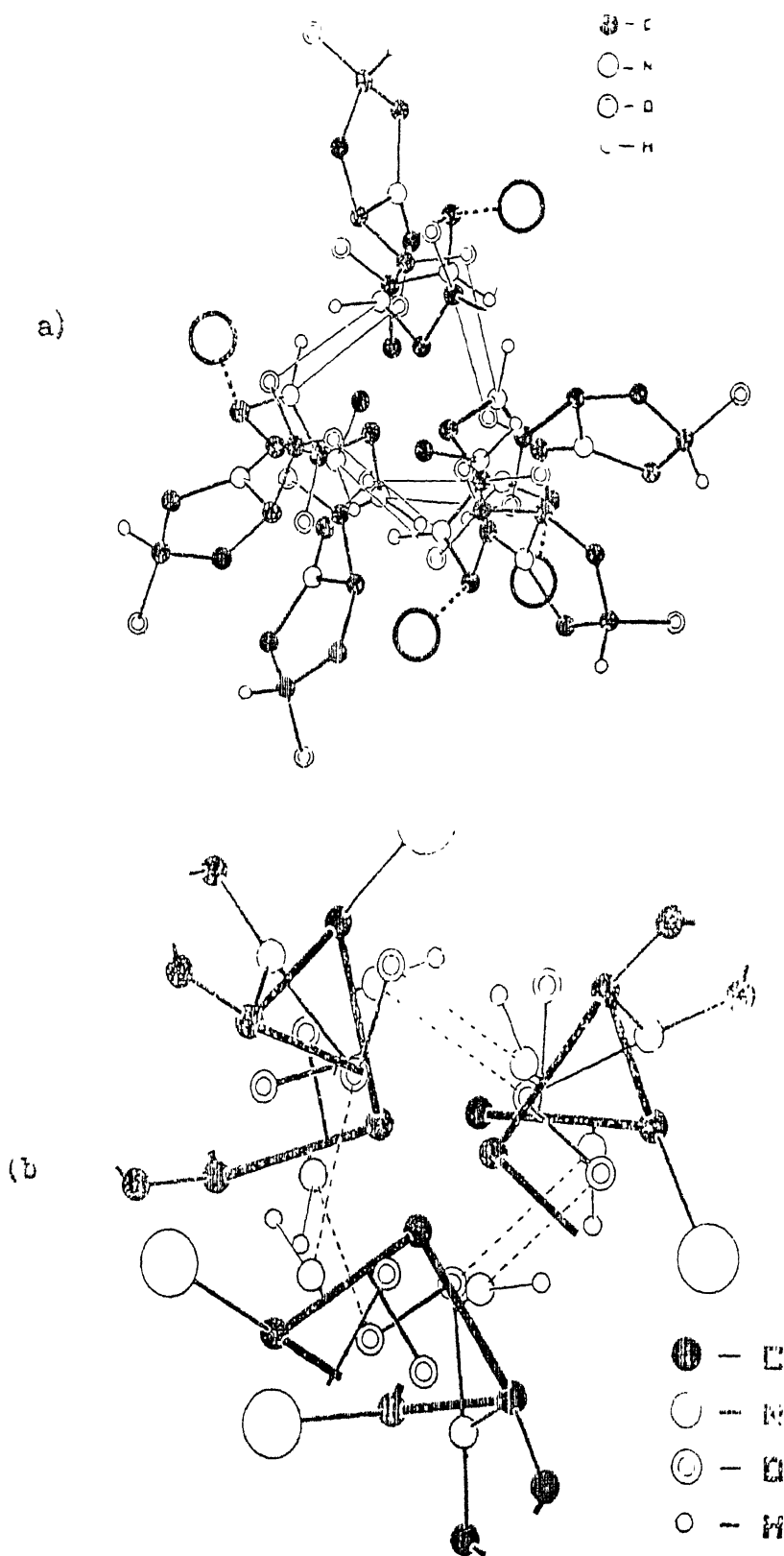the backbone atoms. (From (1)).

Fig.8(a) Projection of the atoms, in the backbone and proline groups, in a height of three residues from 0 to 8.6 Å along the c-axis, drawn analogous to Fig.4(a). Note that all the features pointed out for the non-supercoiled triple helix in Fig.4(a) are preserved in this also.

(b) This diagram brings out the path of the backbone peptide units (which is seen to follow the pattern in Fig. 6(d). In addition to α-carbon atoms, N—H and C—O groups involved in hydrogen bonding are also shown. (From (6)).

the fibre axis, that every residue is a planar _trans_ peptide

unit, that the individual helices are left-handed, and that

Pro and Hyp residues can be readily incorporated, in addition

to the stereochemical requirement that Gly should occur once

in every three residues in each peptide chain — are still

preserved in the supercoiled structure. As will be seen

from Fig. 7, the supercoiling about the central axis of the

protofibril is quite small, being only approximately 90° for

the whole length of the molecule that is shown in the figure.

By this time, the subject of the triple helical structure

of collagen had drawn the attention of two other laboratories

in Britain (Rich and Crick (11) in Cambridge and Cown _et al_

(12) in London) who had also examined the collagen structure

on the basis of their observations on poly-glycine (8) and

poly-L-proline (7) respectively. They obtained essentially

the same conclusions as we did, namely that a detailed fitting

with the spacings, of the layer lines in the observed x-ray

pattern of collagen, requires that the structure must be built

of three individual left-handed helices, with right-handed

supercoiling, as shown in Fig. 8(b). They also used a value

of 10/3 for n, the same as what was adopted by us in 1954-55.

There were minor differences in the details as between

the proposals from the three laboratories (see the next section),

but the basic features were the same in all of these. That

this structure is essentially correct is shown by the optical

diffractometer pattern (Fig. 9) taken/by Dr. Cowan in King's
in 1955

College , London, using a replica of our molecular structure

shown in Fig. 8(a), and employing diffraction of light, instead

of x-rays, for recording it. It will be seen from Fig.9 that

this pattern exhibits all the features of the x-ray pattern,

and contains the equivalents, of the meridional spacing of

3Å (10th layer), and two other strong non-meridional layers

(3rd and 7th), in addition to a weak 4th layer. In fact

the x-ray photograph is almost completely duplicated by this

optical diffractograph. Thus, the essential features of the

helical structure, and its supercoiling, become well established

by the end of 1955.
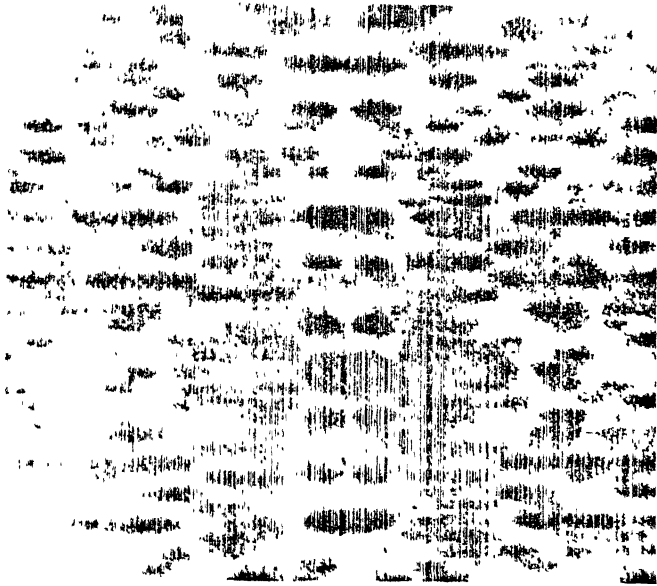
(a)                                    (b)

Fig.9 (a) Optical transform recorded using the supercoiled
       triple helix.

    (b) Diffraction pattern of collagen fibre kept at 75°
        to the x-ray beam. Note the almost exact
        correspondence, in the central region, of the top
        halves of the two figures. (In the x-ray pattern
        the intensity falls off rapidly with the distance
        away from the centre, owing to disorder.) (From (6)).

## 6. The question of one or two hydrogen bonds per three residues

Although there was thus general agreement as to the
helical parameters associated with the collagen triple helix,
there were minor differences in the molecular structure as
proposed by the three laboratories mentioned above. The main
point of difference was, as to whether the number of hydrogen
bonds stabilizing the structure per three residues, was one
or two. Pauling had enunciated the principle that the maximum
number of hydrogen bonds would be made in a stable structure,
and we had therefore tried to introduce both the NH ... O
hydrogen bonds that are possible between adjacent protein
chains, while the third NH was nowhere near the correct
orientation for a hydrogen bond to be formed. Rich and Crick (11)
North
and Cowan, McGavin and/(12) had both concluded that only one
strong hydrogen bond is possible, if the structure is to be
not too closely packed, and, in their structures, the second
NH group that could be involved in a hydrogen bond was left
free. When the preprints of these papers reached us, we
also looked into this question, and it was found that an
one-bonded structure could be built (13). However, this could be

done by rotating the minor helices either in an anti-clockwise direction by +40°, or a clockwise direction by -15°, from the double-bonded structure. Both Rich and Crick and Cowan et al preferred the former (plus) structure, but we found that the latter was superior (Only this minus structure was used for later studies in our laboratory).

During the succeeding years, the structure was looked at in great detail in Madras, and the relative merits of the one-bonded and two-bonded structures, both from considerations of molecular packing, and of their agreement with the x-ray diffraction pattern, were examined in our laboratory. The results were presented in an International Symposium on Collagen arranged by CLRI in 1960 (14). The main conclusion was that the best conformation that could be worked out for both our (minus) one-bonded structure and our two-bonded structure were equally good; but we preferred the two-bonded structure because of the possible stability introduced by the additional hydrogen bond. (In fact, these studies, made mainly in collaboration with V. Sasisekharan and C. Ramakrishnan, brought home to us the fact that there were no hard and fast criteria available in the literature for judging a structure to be good or bad,

and this led us to our studies on protein conformation and on stereochemical criteria for peptide structures in general, but that is another story.)

From the point of view of the x-ray diffraction data, the calculated Fourier transforms of the two structures were very close to one another, and both agreed fairly well with the observed intensity distribution in the layer lines. In fact, independently, Rich and Crick (15) also reported good agreement/ of an one-bonded structure close to our
                    with x-ray data
minus structure.

Attempts were therefore made to resolve the question by means of conformational energy calculations of the two structures (16). Although these gave, as expected, a lower energy for the structure with two hydrogen bonds per three residues, it was found to be only marginally superior, because the extra stabilizing energy was mainly contributed only by the additional hydrogen bond, and the non-bonded energy was, actually, inferior for the two-bonded structure. The situation was puzzling.

In fact, the question of the number of NH ... O hydrogen

bonds per tripeptide in the collagen structure was investigated

experimentally in several laboratories all over the world

during the 1960's, using various physicochemical techniques.

Thus, data from widely different methods, such as deuterium

exchange (Bensusan and Nielsen, 1964 ), tritium exchange

(Englander and Von Hippel, 1967 and McBride / and Harrington 1964

calculation of shrinkage and denaturation temperatures

(Harrington, 1964 ) , all indicated the existence of two NH ... O

hydrogen bonds per three residues, rather than one (for a full

account see (17)).

Although there was all this evidence in favour of the

two-bonded structure, and although, as will be seen from

Fig. 10(a), there is a gap in the cross linking bonds between

the chains in the protofibril of the triple helix in the

one-bonded structure, I had an intuitive feeling that the

structure with two straight interchain NH ... O hydrogen

bonds may not be the final answer, but that there must be some

other way of building the hydrogen bonds for the NH groups in a

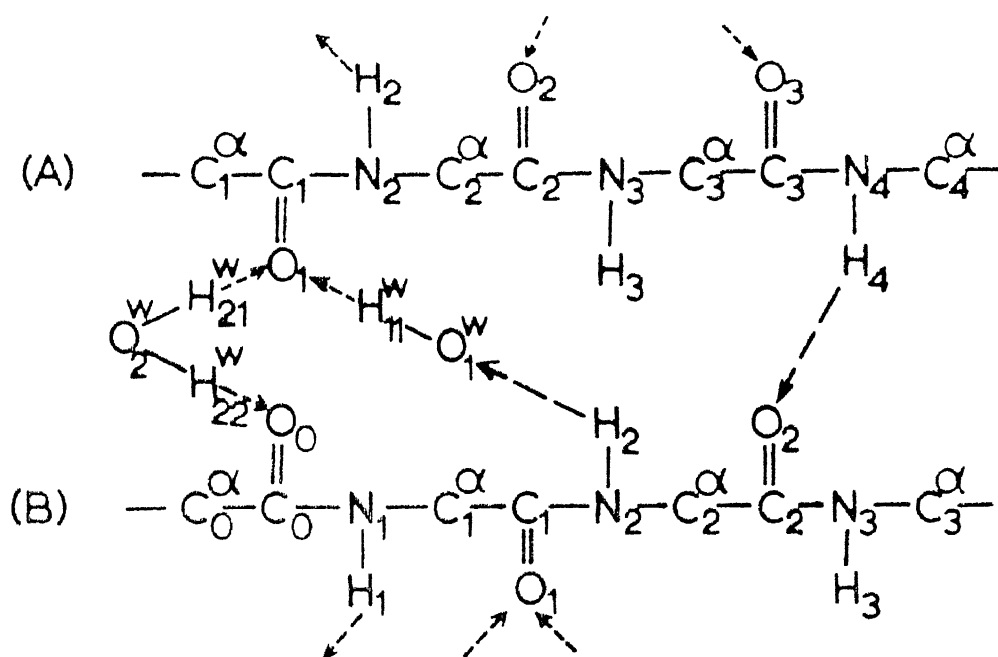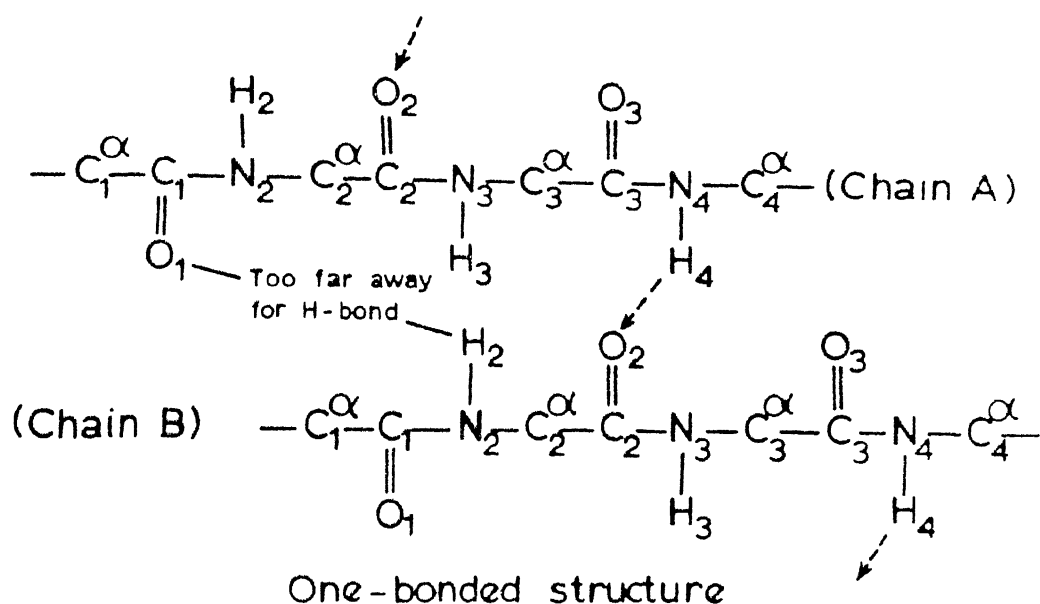structure having the chain configuration close to that of the

One-bonded structure



Water-bridged structure

Fig.10(a): Diagram showing the only set of hydrogen bonds of the type
$N_4H_4 \ldots O_2$ linking the backbones of neighbouring peptide
chains in the collagen triple helix. Note that these are
too few, and are not sufficient to produce good stability
for the structure.

(b): On introducing two water molecules with $O_1^w$ and $O_2^w$, several
OH ... O linkages are produced, in addition to a second
NH ... O hydrogen bond of the type $N_2H_2 \ldots O_1^w$. Note that
practically every NH and every CO is involved in the hydrogen
bond network. (Reproduced from (3)).

one-bonded structure. This idea was supported by the

observation that all the techniques mentioned above indicated

the existence of <u>one strong</u> hydrogen bond and <u>one relatively</u>

<u>weaker</u> hydrogen bond for every three residues — the weakness

and strength being indicated by the ease with which the

hydrogen atoms in the NH groups are displaced by deuterium,

or tritium, as the case may be.

Till about 1967, there was no clear resolution from theory

of this problem. However, a study made in Chicago along with

R.Chandrashekaran gave us a very satisfactory solution (18).

This was that the strong NH ... O hydrogen bond is a direct

one $(N_4 H_4 \ldots O_2)$ between the backbone atoms of neighbouring

chains, while the second weaker one is from the free NH group

$(N_2H_2)$ of Fig.10(a) to the oxygen $O_1^W$ of a water molecule which

again links up with the oxygen of a CO group $(C_1O_1)$ in a

neighbouring chain, <u>via</u> on OH ... O bond (Fig. 10(b)). Then,

as will be seen from Fig. 10(b), it is also possible to

have one more water molecule, $(O_2^W(H_{21}^W , H_{22}^W))$, both of whose

hydrogens bond with CO groups of neighbouring peptide chains.

All this is seen in a wire model of the collagen structure

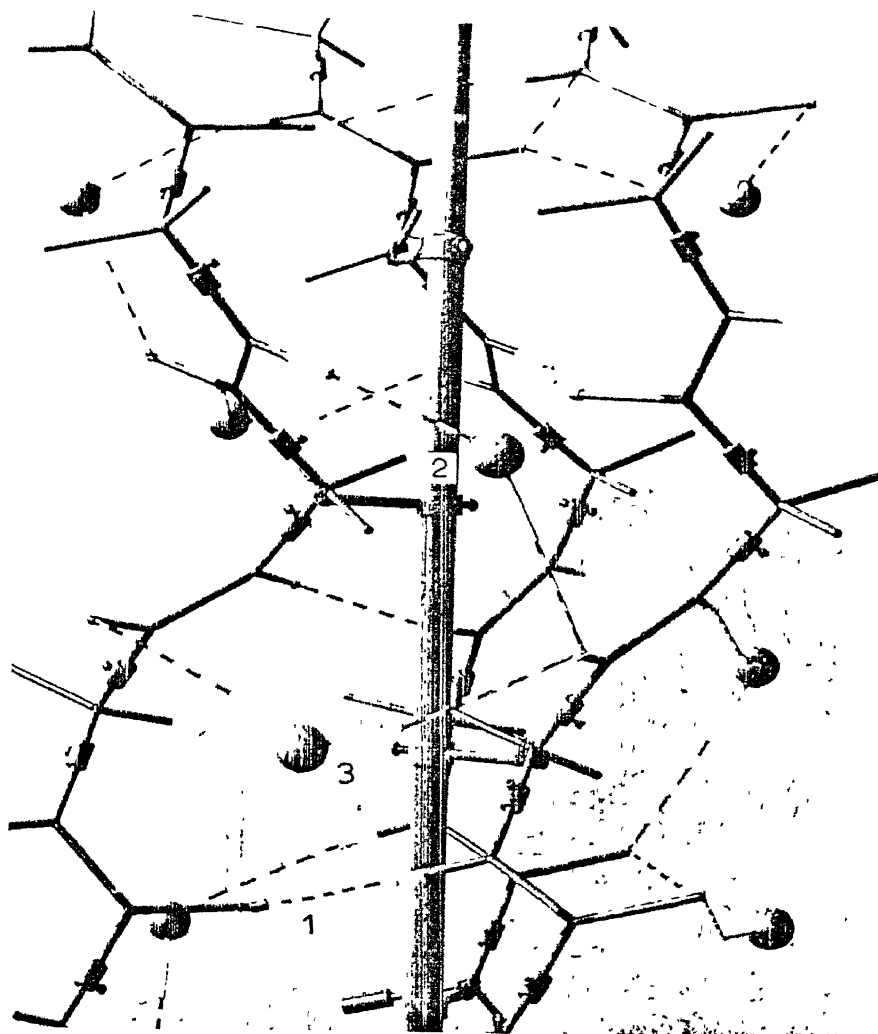with water bridges shown in Fig. 11.

Fig. 11 : Photograph of a wire model showing the hydrogen bond cross links in the water bridge structure. The rods represent the backbone bonds and the balls water oxygens. Hydrogen bonds are indicated by striped wires.

(Reproduced from (18)).

Water molecules are well-known to be firmly bound to

protein structures, but this feature had not been properly

taken into account previously for collagen, although

Cowan et al (12) had suggested it. Soon after, in 1969,

Yonath and Traub (19) also cameout with possible locations

for bound water molecules in the collagen structure based on

their refinement of the structure of poly(Gly—Pro—Pro),

which is relevant for the portions of the collagen triple

helix having this sequence.(see also (20)).

However, this was not the whole story, for the same water

molecule involved in the NH ... O bond in Fig. 10 (b) could

also be shown to serve the purpose of forming additional

hydrogen bonds with the hydroxyl group of hydroxyproline side

chains, producing still greater stability to the triple chain.

This is discussed below.

### 7. Hydroxyproline and its role in the collagen structure

It should be mentioned that vigorous studies, of a

physico-chemical and biochemical nature, had been going on

in various laboratories during the 50's and 60's, and definite

evidence that the collagen molecule is built up of three

chains, and that they form a triple helix, of the nature

indicated above, had become well established by experiment by

1970. There is no time to comment on these studies[*]

However, one important fact relating to collagen which had

not yet been given a proper theoretical basis stood out

prominently — namely the role of hydroxyproline, an amino

acid residue which occurs almost exclusively in collagen

among all proteins. It should be mentioned, in this connection,

that hydroxyproline (Hyp) is not one of the 20 amino acids

that are coded by the genetic code, and it is not incorporated

as such during the metabolic building up of the collagen

chains in living tissues (see the review by Prockop (22) for

further details). It is only incorporated as proline, and,

after the formation of the triple helix, certain   of the

proline residues are hydroxylated by a special enzyme known

as proline hydroxylase. It was also known that these Hyp

residues occur only in the third position/in the repeating
                                        Y

sequence (— Gly — X — Y — ) of the collagen structure,

and that the hydroxylation is of the  4-trans type universally

in mammalian collagen, although other isomers such as

4-cis-Hyp and 3-Hyp, are observed in other systems. Therefore,

* Other articles in the two volumes mentioned in (1) and (2),
  as well as the extensive reviews by Traub and Piez (20),
  and Bornstein and Traub (21), may be referred to for
  further details.

the question occurs prominently — Why are such Hyp residues

manufactured in the collagen structure, and what purpose do

they serve in its physical chemistry or biochemistry? The

problem was in the air in the field of collagen research

for a long time; but an answer to this came from stereochemica.

considerations as was found by us (23) in the early 1970's

in Bangalore.

After I had moved to the Indian Institute of Science,

Bangalore in 1971, we had a visit of Prof. R.S. Bhatnagar

from USA, who had earlier worked on the biochemistry of the

hydroxylation of proline with Prof. D.J. Prockop. During a

discussion that took place over the table, Bhatnagar asked

me pointedly as to what explanation I can give for the

existence of hydroxyproline, and that too, with the particula

4-trans hydroxyl group that occurs in the side chain of Hyp

residues. On inspecting the model, it became pretty clear

to me that one of the water molecules, which was incorporated

by Chandrasekaran and me into the collagen structure, had a

free OH group, and that this was pointing towards the oxygen

latter
in the hydroxyl group of hydroxyproline, if the/is present in
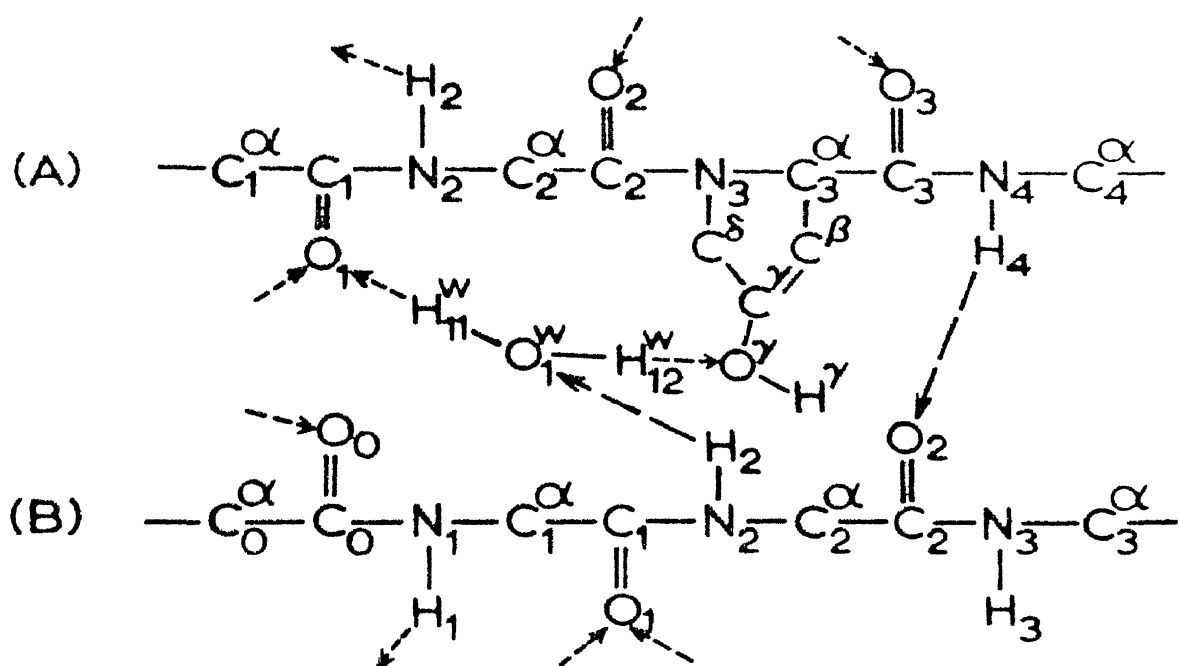
Water - bridged structure (with hydroxyproline)

Fig. 12(a) : More hydrogen bond bridges introduced by hydroxyproline O H group. Note that in addition to those in Fig. 10(b), there is a linkage via hydrogen bonding between $O_1^w$ and $O^\gamma$ and that $O^\gamma$ $H^\gamma$ could serve for possible linkage with another triple helix. (From (3)).
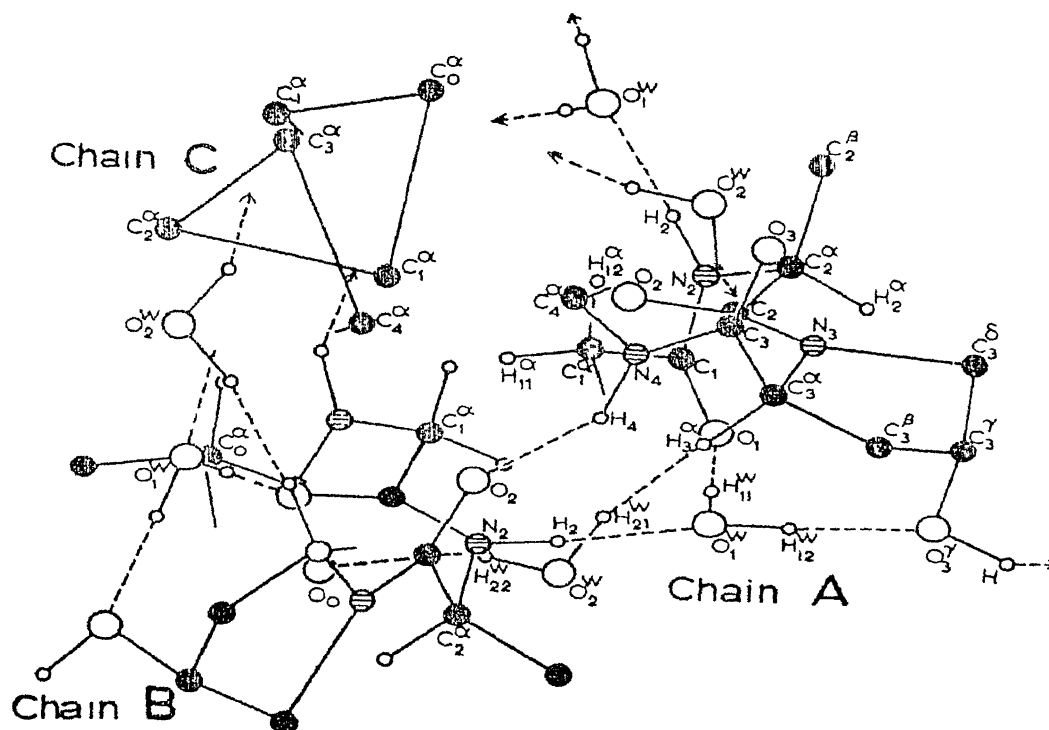
Fig.12(b): The above features are seen in projection
in this figure. Note the occurrence of an
extensive network of hydrogen bonds, linking
neighbouring chains of the triple helical
protofibril. (Reproduced from (3)).

Y

the third position from glycine i.e., the position/previous

to the next glycine. Immediately, I asked Miss Manju Bansal,

who had joined me as a research student, to take the

coordinates of the structure as worked out by Chandrasekaran

and see if it will fit in this extra hydrogen bond involving

the hydroxyproline residue, if slight modifications are made

in the orientation of the surrounding region in the molecule.

It took Manju not more than two weeks to make a thorough

examination of the possibilities, and to arrive at an

acceptable set of coordinates for the collagen triple helix

Y

containing hydroxyproline in the third position/, and/or

X

proline in the second position/, in which this extra

hydrogen bond is also made.

A schematic representation of all these hydrogen bonds,

following the pattern in Figs. 10(a,b) is shown in Fig. 12(a,

and an actual projection down the helical axis of the various

groups involved in the hydrogen bonding network, making use of

the water molecule and hydroxyproline OH group, is shown in

Fig. 12(b). As mentioned before, there is a strong NH ... O

hydrogen bond in the deep inside of the structure, and there is

a second network, consisting of NH ... $O^W$, and $O^WH$ ... $O$ ===

involving the water oxygen $O^W$, linking the backbone. In

addition, there is one more hydrogen bond linking the water

oxygen $O^W$ with $O^Y$ of the Hyp side chain. Further, the

hydroxyl group of the Hyp side chain is also pointed outwards

from the centre of the triple helix, and could serve for

interconnections between neighbouring triple helices via

hydrogen bonds. Thus, hydroxyproline plays an important

part not only in the stability of a single protofibril ,

but also in producing strong linkages between neighbouring

protofibrils.

In this region of research also , experimental biochemists

had been gathering information on the role of hydroxyproline.

(see 22)
For instance, Berg and Prockop /    showed that the melting

temperature of hydroxylated collagen was much higher than

that of the unhydroxylated triple-helical protocollagen

molecules. In fact, the melting temperature of the newly

formed collagen molecule, without any Hyp residues in it,

is as low as 25°C, while fully hydroxylated collagen has a

melting point of 50°C. Since collagen is found mostly in

mammalian systems with body temperatures in the region

of 30° to 40°C, these experiments clearly show that collagen

can be chemically stable in the bodies of all animals only

if it is hydroxylated. These experiments, and several others

made using other techniques such as enzyme degradation etc.,

(see (22)) completely supported the results obtained from

conformational theory that the real role of hydroxyproline

is not chemical, but rather physical, in that its occurrence

helps in the formation of a network of two, or three, hydrogen

bonds for every three residues.

With this, the essentials of the stereochemistry of the

secondary structure of the collagen molecule could be said

to have been firmly established.


### B. Primary sequence with glycine as every third residue

However, the complete resolution of the chemical sequence

of the collagen molecule from the biochemical point of view

became revealed only in the early 70's — thanks to the

project involving
efforts of Karl Piez, who initiated/a large international

for
group with workers in Germany, Britain. and USA,/determining

the primary sequence of the thousand-residue-long collagen

molecule, forming one of the three roughly equivalent chains

in the triple helix (see the article by Piez (24) for a survey).

When this primary structure was solved, it gave full support

to the theoretical assumption, made in the very first structure,

that Gly must occur at every third position, in the helical

regions of the structure. It also indicated that Pro invariably

occurs at the second position X (after Gly), and Hyp always

occurs at the third position/(before Gly), although Pro residues

may remain partially unhydroxylated in the third position also.

Table 1 shows a portion of this primary sequence, in a

number of different species of animals, all of which are seen

to have this property. However, a more complete analysis of

the protocollagen molecule indicated that collagen is primarily

synthesized with an extension, both before the first residue

having the regular sequence — Gly — X — Y — , and after

the termination of this regular sequence. It is believed that

these appendages help to align the three chains in proper

justaposition, so as to initiate the formation of the triple helix.

With this, it can be said that all the mysteries of

the collagen primary structure (amino acid sequence) and

secondary structure (stereochemistry of the triple helix)

Table 1:   Typical amino acid sequences[*] in a part of the helical region of the collagen peptide chain[†]

| Residue | $\alpha 1(I)$ rat | $\alpha 1(I)$ calf | $\alpha 1(I)$ chick | $\alpha 1(II)$ steer | $\alpha 2$ rat | $\alpha 2$ calf |
|---|---|---|---|---|---|---|
| 17 | Gly | Gly | Gly | Gly | Gly | Gly |
| 18 | Pro | Pro | Pro | Val | Pro | Pro |
| 19 | Met | Met | Met | Met | Met | Met |
| 20 | Gly | Gly | Gly | Gly | Gly | Gly |
| 21 | Pro | Pro | Pro | Pro | Leu | Leu |
| 22 | Ser | Ser | Ala | Met | Met | Met |
| 23 | Gly | Gly | Gly | Gly | Gly | Gly |
| 24 | Pro | Pro | Pro | Pro | Pro | Pro |
| 25 | Arg | Arg | Arg | Arg | Arg | Arg |
| 26 | Gly | Gly | Gly | Gly | Gly | Gly |
| 27 | Leu | Leu | Leu | Pro | Pro | Pro |
| 28 | Hyp | Hyp | Hyp | Hyp | Hyp | Hyp |
| 29 | Gly | Gly | Gly | Gly | Gly | Gly |
| 30 | Pro | Pro | Pro | Pro | Ala | Ala |
| 31 | Hyp | Hyp | Hyp | Ala | Val | Ser |
| 32 | Gly | Gly | Gly | Gly | Gly | Gly |
| 33 | Ala | Ala | Ala | Ala | Ala | Ala |
| 34 | Hyp | Hyp | Hyp | Hyp | Hyp | Hyp |
| 35 | Gly | Gly | Gly | Gly | Gly | Gly |
| 36 | Pro | Pro | Pro | Pro | Pro | Pro |
| 37 | Gln | Gln | Gln | Gln | Gln | Gln |
| 38 | Gly | Gly | Gly | Gly | Gly | Gly |
| 39 | Phe | Phe | Phe | Phe | Phe | Phe |
| 40 | Gln | Gln | Gln | Gln | Gln | Gln |
| 41 | Gly | Gly | Gly | Gly | Gly | Gly |
| 42 | Pro | Pro | Pro | Asn | Pro | Pro |
| 43 | Hyp | Hyp | Hyp | Hyp | Ala | Hyp |

[*]Note the occurrence of Gly as every third residue, of Pro always in the second position, and of Hyp in the third position, from Gly (The latter may be partially hydroxylated).

[†]Data from (24)

## The Age of Computerization in Science

## 1. Introduction and preview

Prof. Narasimhaiah and friends,

I am deeply grateful to your President Prof. Narasimhaiah for having invited me to give the Valedictory Address of your Science Forum. I wondered wnat would be the best topic that I could choose and I decided that it should be sometning with which I am closely familiar and which is occupying my current interests. Therefore, I chose this subject, since computation has been the mainstream of our scientific activities during the last 40 years or more, and we could consider examples from our own fields of interest as illustrations for the lecture.

In tnis again, the scope of the talk will be restricted, since I am not a specialist in computer science as such. However, since I am quite familiar with its applications to mathematics, physics, biology, technology, and in recent years to logic, this would be a personal talk based on my experience with computation in these fields. I shall very briefly review tne aspects of the growth of these subjects from the 40's to the 80's as seen by a scientist, wno has used computer

can now improvise games using computers which are even being

utilized to train persons in the practice of games such as chess.


Perhaps it may be worthwhile for me to recall the changes

that have taken place in the utilization of computers in

scientific research. When I entered the research field in

physics in 1942 under Prof. C.V. Raman, I was using only table

calculators which were mechanical and which had to be cranked

for each step of addition even for doing multiplication, and

made sucn a grinding noise. In my doctorate study with Prof.

Raman, I employed these almost incessantly for almost 1 - 2 years

to grind out and test the formulae developed by us regarding

the optical dispersion of crystals. Even when I went to

Cambridge, England in 1949, the position was hardly different.

Electrical calculators had come into use instead of mechanical

ones, but these were only makes that imitated the mechanical

calculators, and the only improvement was that the turning

and shifting were done electrically by pressing a key, and

they were much less noisy. Logarthmic tables still formed the

essential assisting devices for doing calculations. However,

two subjects of crystal structure determination and the study

of biomolecular structure and reactivity, are inextricably

mixed together now-a-days. In both these fields, we in India

have played quite a useful part in the development of the

subject and some thirty years ago the work produced in India

was of the same level as that obtaining in the best laboratories.

This was possible because advanced facilities and computation

had not yet become the dominant part of such investigations.

I shall illustrate how in both these fields continued activity

at the top level has come out of our laboratories so that

even today we are ready to take up work in the advancing

front of knowledge.


Perhaps the most beautiful application of the consequences

of Fourier transforms in crystallography is in the field of

tomography, a subject that came into prominence in the 1970's,

and in which I had the good fortune to write one of the earliest

papers on the application of very specialised techniques based on

Fourier transforms, for obtaining three-dimensional images from

two-dimensional pictures. The theory which we developed, based

of the developed world. But we lost the initiative mainly

because computational facilities did not continue to be at

the same level as the developed countries, and we are now

slowly picking up the thread, thanks to the great stress being

laid on thrust areas of research in recent years.

I shall conclude the lecture by giving you some ideas on

computation as applied to logic, reasoning, and the processing

of knowledge. The topic of artificial intelligence is much

talked of now-a-days, but apparently, even in this intricate

subject, newer approaches can be developed by trying to put

in mathematical language, the essential principles of the

reasoning process, considering it as a physical phenomenon

subject to observation, analysis and systematization. I shall

try to show you that the process of algorithmising "reasoning"

can be done and that it can also be computerized. If there is

time, we shall discuss what this means in the ultimate limit

where we can ask the question, whether computers can replace

men for producing original ideas. If I may forestall what

our conclusion will be, I feel that, while the computers may

increase speed, range, and scope, of thinking processes, it can

never replace the originality and analytical power of the

human mind which form the essential basis of intelligence.

along the length of the rope. In the case of a crystal, its

x-ray diffraction pattern contains many reflections — of the

order of 500 for a simple structure studied in the 50's,

which increased to some 10,000 in the 60's, and 100,000 to

500,000 in the 70's and 80's. In principle, from the intensities

of these reflections, it is possible to calculate the distribution

of electron density in the repeating pattern of the crystals,

and hence the positions of the atoms in the molecule, their

spatial arrangement, the nature of the chemical bonds and so on.

Fig.1 gives the essential principles involved in the mathematical

reconstruction of the molecule from the x-ray diffraction pattern.

---

Fig.1 : Essential formulae of x-ray crystallography

---

In principle, the diffraction pattern, specified by $F(hk\ell)$,

is the Fourier transform of the electron density distribution

$\rho(xyz)$, as shown in Eq.(1), and this equation is reversible

and can be used to give the latter, in terms of the former, as

in Eq.(3). But in practice, there is an unknown factor $\alpha(hk\ell)$

in (2), which has to be determined by special techniques, or

has to be guessed by the crystallographer and then refined and

improved by calculation. In the 1950's, there were no hard

and fast methods for determining the phases in a general case,

but computers were being applied for the arithmetical

manipulations required in Eqs (3a) and (3b), and the increase

in power of the computers has gone side by side with the increase

in size of the molecular structures that were solved by x-ray

crystallography. In the 60's, mathematical theories for the

determination of $\alpha$ from $|I(hk\ \ell)|$ were developed, and the

phase problem became more amenable to computerization. So also,

the process of mounting a crystal and rotating it to various

positions and recording reflection intensities, were also

computerized and automatic diffractometers came into being.

Thus, the labour involved in a crystal structure determination

has been minimized to/almost unbelievable extent during the
　　　　　　　　　　　　　an

last 30 years, by the progressive use of computers at every

level of experiment and theory. In fact, people are talking

of a completely computerized structure determination in which

the crystal is just fixed to the x-ray apparatus and everything

else is computer-implemented and/outcome /a series of diagrams
　　　　　　　　　　　　　　　　　 the　　　 is

indicating positions of the atoms in the crystal.

# Growth of x-ray crystallography in size and complexity

| Years | Nature of molecules (examples) | Number of atoms in a molecule | | Number of reflections N (Maximum) | Computers | Memory |
|-------|----------|---------|----------|----------|-----------|--------|
| | | Maximum | In India | | | |
| 1950's | Amino acids, peptides vitamins | 50 | 20 | 1,000 | Home-made Elliott-803 | 8K |
| 1960's | Vitamin B12 Proteins — first examples | 1,000 | 50 | 10,000 | IBM-7090 CDC-3600 | 64K |
| 1970's | Enzymes Immunoglobulin Transfer RNA | 10,000 | 100 | 100,000 | IBM-370 DEC-10 | 512K |
| 1980's | Viruses 50 S Ribosome particle | 100,000 | 10,000 | 500,000 | CRAY-1 CRAY-2 | Unlimited |

②

I have been talking of molecules and their sizes.

Fig.4 gives an idea of the relative sizes of molecules, and

---

Fig.4 : Molecular size comparison .

---

the diagram is self-explanatory. The molecules shown at the

top are simple organic compounds — for example, lysine is an

amino acid, and heme is a planar compound of great importance in

the blood protein hemoglobin which is shown below. In the

latter, the thick rod-like structures shown, actually consist

of a chain of atoms . composed of /amino acid residues and
                                a sequence of

they are folded together in this convoluted form to make the

biologically active blood protein hemoglobin, which transports

oxygen from the lungs to the tissues. The location of the heme

molecule is also shown by the black discs in the figure. The

increase in size of the protein molecule, in relation to small

organic molecules, is clear from the figure. This jump in

size took place in the early 60's and this has gone up by

another factor of 10-50 during the last twenty years, and even

"living molecules" like viruses have been tackled in the last

few years.

## 3. Biomolecular structures

Leaving aside crystallography as such, we may ask the

question ― "Why does the molecular chain in/hemoglobin , large molecules, such as

take such a convoluted structure, and how can this be understood

from chemistry and physics, and how can this be correlated with

biological activity?" Since this talk is not purely about

biochemistry or biophysics, I shall only make brief comments

about this. The essence of the process of protein folding is

contained in Fig.5. As will be seen from this, a protein molecule

---

Fig.5 :   Folding of polypeptide chains

---

is a long chain of repeating peptide units, each of which is one of

20 different types, but which are closely similar, and differ

only in what are known as side groups (R). The chain is flexible,

and therefore it folds back on itself and gets stabilized by

"chemical bonds"such as   ―S―S― linkages, or "physical bonds"
                      interactions and
such as electrostatic/hydrogen bonds, occurring between different

parts of the chain,which hold them together in the form of a

loop or similar structures. It is natural that the whole

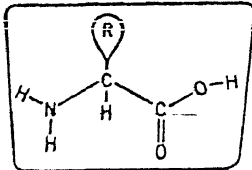molecule will take up a compact globular configuration having nearly
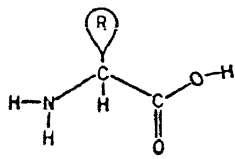
UN IONIZED AMINO ACID
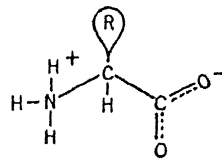
ZWITTERION

*The amino acids exist in solution as doubly charged zwitterions*

THE AMINO ACID MONOMER
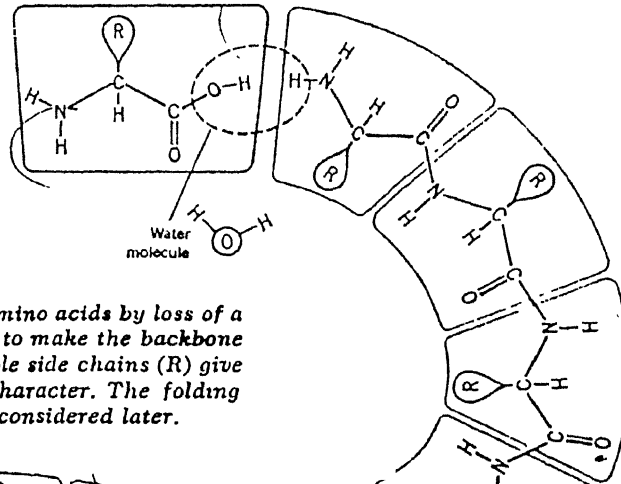
Water
molecule

*The peptide bond formed between amino acids by loss of a*
*water molecule joins identical units to make the backbone*
*of the polypeptide chain. The variable side chains (R) give*
*each protein chain its distinctive character. The folding*
*of the chain in three dimensions is considered later.*

Hydrogen bond

Disulfide
bridge

*Polypeptide chains*
*are cross-connected*
*by hydrogen bonds*
*and disulfide bridges.*

(5)
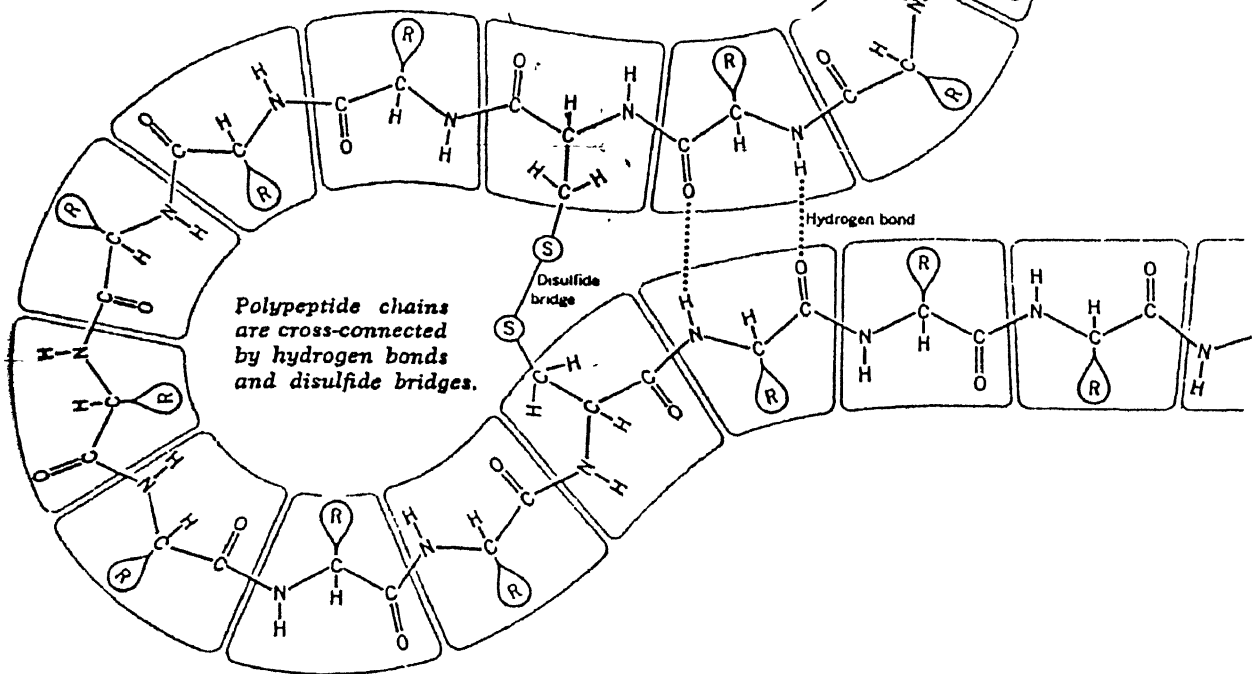
the two strands of the double helix, or the three st ands

of the triple helix, as the case may be. Fig. 7(b) shows tne

twist of the three strands around one another in the collagen

triple helix. This twist is regular because the peptide

---

Fig.7 (a) Peptide chain configuration in myoglobin which is
          closely related to hemoglobin.

      (b) Twisted rope configuration of the collagen triple
          helix.

---

sequence in collagen is based on repeating / tripeptide (— Gly—X—Y—)

and half of X and Y are also based on the / same amino acid residue

proline. On the other hand, the sequence is irregular and

different in myoglobin, / and this favours the formation of single

$\alpha$-helical segments, which are then folded up to form a globular

structure as in Fig. 7(a).

Obviously, in checking the validity of such structures,

and of the wide range of possibilities that may occur, computers

have played a very important part. However, the early studies in

the 50's were mostly done without computers, and in fact a

pioneering study of the principles governing these, which were

published from Madras in 1963, was completely done with

mechanical desk calculators, which took almost a year to complete

(see Fig.8). Now-a-days, ^these / will take a fraction of a second

on the best computers. I shall not mention much about these

theoretical ideas except to indicate that the two principal

considerations that were put forward by us at that time have

played a dominant part in all techniques involving molecular

packing, and the fitting of molecular chains to electron density

diagrams obtained from crystal structure determination, ever since.

These are the use of dihedral angles (rotation angles), such as

$(\phi, \psi)$ of Fig. 8(a), and the idea that non-bonded atoms cannot come

closer than certain specified/ ^limiting distances, which was used for

preparing the diagram in Fig. 8(b). These have been so well

---

Fig.8 (a) Rotation angles $\phi$ and $\psi$ specifying the configuration
of a pair of peptide units.

(b) The $(\phi, \psi)$-plot of the observed configurations
in myoglobin (marked as small circles.)

---

computerized that, now-a-days one can just feed in electron

density and the chemical nature of the molecule and ask the

computer to shift, turn and twist the molecule until the best

fit with the electron density is obtained, subject to the

restrictions imposed by the limiting distances mentioned above.

More about this later.

## 4. Crystallography and Biology with special reference to drug design

The precise techniques / that have been developed for
energy minimization, and for the

working out of the best possible configurations of molecules,

have resulted in a body of knowledge which has vastly expanded

our knowledge / of the physical chemistry, molecular biophysics and

biochemistry, / of the biological systems. I shall illustrate these by a few examples.

We had seen that the heme molecule in hemoglobin is located

in between the different protein chains and the knowledge

obtained by crystallography about the particular amino acid

residues near the location of the heme had led to a clear

understanding of the mechanism of the function and reactivity

of / the protein for oxygen transport in blood. Simultaneously, it has / also become

clear that many biological conditions associated with / the malfunction

of hemoglobin can be attributed to changes in one or two

amino acid residues produced by mutation. Fig.9 shows the

location of the mutated residues for a large number of

---

Fig.9 Mutation sites in pathalogical hemoglobins.

---

pathalogical hemoglobins. It will be seen that, though the amino

in a linear peptide chain
sequence/may not show it,      the x-ray crystallographic

three-dimensional structure shows that they often occur

close to the region where the heme is held, or in the gaps

between neighbouring molecules. Apart from this general

feature, the knowledge obtained from crystallography has

profoundly affected biochemical investigations for the
                              of the blood,
understanding of such  diseases/and has led even to the design

of suitable chemical drugs for alleviating the symptoms.

Turning to    more fundamental problems in biology,

such as the action of enzymes, the chemists had formulated

the theory that enzymes, which regulate all biological reactions,

did so by forming a close molecular fit with the substrate

molecule concerned. This was a conjuncture in the past,

although a very reasonable one and much substantiated by all

available knowledge. However, a confirming evidence in this

matter has come only from crystallographic studies. One of

the earliest proteins to be solved, namely lysozyme which is

shown in Fig. 10, displays this feature remarkably well.

Fig. 10 : Space-filling molecular models of free lysozyme (left)
        and when bound to its substrate (right).

The first picture is that of the protein which is seen to have

a cleft on its left, and the second picture is of the enzyme-

substrate complex. It will be seen that the substrate and

enzyme form a perfect fit and the theoretical postulation of

molecular fit has been demonstrated in all detail by three-

dimensional x-ray crystallography.

I will therefore restrict myself to the subject of

molecular fit, and give a few examples where crystallography,

in combination with conformational theorey, has been applied

to obtain information regarding this. Fig.11 is a study made

by Prof. Viswamitra of the Indian Institute of Science

---

Fig.11 :  Space-filling models of (a) the antibiotic TANDEM
          and  (b) its complex with DNA.

---

in which he determined the molecular structure of an antibiotic

TANDEM, shown in the top half of Fig. 11, which acts by binding

with DNA. He could then show that this fits very well to the

DNA double helix, as shown in the lower half. This idea has

been employed in a number of studies for finding the basis of

drug-molecule interactions. Such crystallographic studies

have even been extended to virus molecules and the interaction

of antivirul agents with the virus. Fig. 12 is an example of a

very recent study made by Prof. Rossman in Purdue University, USA.

---

Fig. 12 : Crystallographically determined mode of binding of
          antiviral agent WIN to virus protein VP1.

---

Not only does it show that the viral protein has a cavity which

can be fitted by the antiviral agent WIN, so that the normal

biochemical reactions involved in the reproduction of the virus

molecule are inhibited, but it is also an example of the extent

to which crystallography has progressed. The virus itself

required 200,000 reflections for its structure determination.

After that, the mode of attachment of the drug was determined

by using what is known as a difference-Fourier synthesis,

in which the difference in intensities between x-ray diffraction

patterns, of the pure virus and the virus-drug complex, was

used to find out the location of the drug molecule within the

virus. Fig. 13 shows the electron density distribution of the

drug molecule as reconstructed by a computer. It will be seen

---

Fig. 13. Front page of Science magazine containing an
         illustration of the molecule of an antiviral agent
         fitted to the computer-displayed electron density cage.

---

that it effectively indicates a cage within which the molecule

must be situated, and the exact location of the molecule and

its conformation can also be varied and fitted in, by employing

the computer for on-line operation. I had the good fortune

to advise some of the workers in this field of computer-fitting

of molecules in USA in 1977-78, and this approach is

universally employed both in crystallography and conformational

theory now-a-days, and "computer graphics terminals" are even

commercially available for this purpose.

## 5. Conformational theory for molecular fit

Coming back to theoretical calculations, Fig. 14 is an

example where the crystal structure determination of an

---

Fig. 14: Daunomycin fitted to DNA by energy minimization on
         a computer

---

antibiotic (daunomycin) was the starting point of an attempt

to fit the molecule to the DNA double helix. The work was

done in Prof. Pullman's laboratory in Paris, and as the figure

will show, the fitting can be done readily by opening out the

double helix in the region where the drug makes its entry.

I wish to emphasize that this has not been done by constructing

a model, and then turning and twisting the molecular structure

so as to obtain the fit, but purely via the use of a computer.

On the other hand, the fitting of/inhibitor, $\beta$-P-P-P, to the

an

active site of the enzyme thermolysin, which was carried out

by Prof. Rao in Bangalore, was done, not automatically on a

computer, but using the computer without on-line interaction.

The investigator starts with a structure close to a reasonable

fit, and then modifies it step by step, by feeding in and

taking out the data from the computer in the form of drawings.

This study showed that the inhibitor enters the enzyme in one

---

Fig.15: Mode of attachment of the inhibitor $\beta$-P-P-P which
imitates the substrate of the enzyme thermolysin.

---

configuration and then both its conformation, and that of the

enzyme in the local region, are modified so as to obtain the

best fit. The reason why I am showing this picture is to indicate

that some of the latest ideas of this type have been computerized

right in our country and interesting consequences have been

deduced. However, even with a computer at one's disposal, this

is a long process requiring weeks or months. If only an interactive

graphic display system, as has been used in all the advanced

countries, were made available, the time could be reduced to

probably a week or two. In fact, as already mentioned, such

systems are available commercially in USA and Japan, and

I believe that it is time that such up-to-date equipment be

made available to the capable scientists of our country.

Emergence of the new field QSAR

Such studies on molecular fit have opened many new

avenues in the field of rational drug design,because drugs

invariably either promote,or inhibit,particular biological

reactions,either in the organism producing infection,or in the

host for producing resistance. In fact, a new field entitled

Quantitative Structure Activity Relationship (QSAR) has emerged

during the last few years, and strategies are being developed

for actively pursuing this for drug design. Fig. 16 shows the

title/of a very interesting paper, on protein crystallography
      page

and computer graphics,published in the German journal Angewandte

---

Fig. 16: Title page and flow chart about QSAR reproduced from
         a recent article.

---

Chemie (Applied Chemistry), which emphasizes this aspect of

biomolecular physics. As will be seen from one of the flow

charts in this report, computer graphics and theoretical

$\overset{\text{a}}{\text{calculations play}}$ as important/part as the determination of

three-dimensional structure by x-ray crystallography. Even

regular symposia are held on this subject of QSAR, as will be

seen from Fig. 16.


## 6. Computerized Tomography

Turning to the subject of tomography, my interests in

this field dates from 1970 when some very interesting work was

done along with Dr. Lakshminarayanan, at the University of Chicago

on the application of Fourier transforms and the novel technique

$\overset{\text{the title page of}}{\text{of convolutions for this purpose. Fig. 17 shows/our first paper}}$

on the subject entitled "Reconstruction of substance from shadow"

---

Fig. 17:  Fourier transform equations formulated in polar
coordinates for axial tomography.

---

As you will see from this, the Fourier transform Eq.(1), which

is basic to this subject, is the same as that used in x-ray

crystallography. However, although crystallographers had applied

this technique for finding out the three-dimensional structure

of virus particles from two-dimensional electron microscopic

images in different directions, we suggested that this

reconstruction is much faster by employing polar coordinates

for the Fourier transform as in Eqs. (2a,b), using the

principle of axial tomography. In these, $g(\ell ; \theta)$ is the

linear shadowgraphs at $\theta$, and $f(r, \phi)$ is the two-dimensional

image that is reconstructed for a two-dimensional section at

right angles to the axis of rotation.

This was followed by another paper published in the

Proceedings of the National Academy of Sciences, USA, whose

title page is shown in Fig. 18. In this, the way in which a

series of two-dimensional images can be converted into a

three-dimensional reconstruction by axial tomography is shown

in the diagram. In addition, a new modification of the Fourier

method, which we called as the "convolution method" was presented
this paper, and its
in / basic mathematics is shown in Eqs. (3), (4) and (5).

---

Fig. 18:    Principle of the convolution method of reconstruction
            in axial tomography.

---

In this, the reconstructed two-dimensional image $f(r , \phi)$ is
                                     density                at $\theta$,
a simple integral of the/functions $g'(\ell ; \theta)$, which represent

the
/modified shadowgraphs at $\theta$ , produced as a result

of the convolution process shown in Eqs (4a,b). The detailed

way in which the modified function is calculated is by the

procedure given in Eqs (5a,b). The mathematics is not quite

relevant to this lecture, but what is important is that,

as a consequence, the reconstruction can be performed almost

a hundred times faster than the Fourier method and, as will be

seen from Fig. 19, it is also about ten times more accurate.

---

Fig. 19: Tables 1 and 3 reproduced from the PNAS paper
showing the distinct advantage in speed and
accuracy of the convolution method.

---

Tables 1 and 3 of Fig. 19 are the comparative data of the

accuracy and computing time required for the different methods

of reconstruction , namely Cartesian Fourier transform (FTC),

Polar Fourier transform (FTP) and Convolution (CON). The

most interesting feature is that the larger the amount of

data that is available from measurements for reconstruction,

the faster is the convolution method relative to the Fourier

method. This technique is applied widely in CAT-scan

instruments particularly for medical radiography.

| $\theta_0(°)$ | $a$ | ℜ(in %) for | | |
|---|---|---|---|---|
| | | CON | FTP | FTC |
| 30 | 0.2 | 1.5 | 9.4 | 7.7 |
| 30 | 0.1 | 0.6 | 7.6 | 5.3 |
| 15 | 0.2 | 1.2 | 8.0 | 7.6 |
| 15 | 0.1 | 0.3 | 5.3 | 5.2 |
| 15 | 0.1 | 0.3 | 2.9* | 3.9* |

* $R_0$ = 0.125; for other values, $R_0$ = 0.25.

TABLE 3. *Computing times (t) required using the different methods for a typical object used in the study*

| $\theta_0(°)$ | $a$ | $t$ in seconds | | |
|---|---|---|---|---|
| | | CON | FTP | FTC |
| 30 | 0.2 | 2.0 | 30.2 | 130 |
| 30 | 0.1 | 2.6 | 48.6 | 480 |
| 15 | 0.2 | 3.6 | 41.5 | 130 |
| 15 | 0.1 | 4.8 | 83.0 | 480 |
| 15 | 0.1 | 4.8 | 158.5* | 487* |

* $R_0$ = 0.125 for these values; $R_0$ = 0.25 for others.

(19)

BRUKER MEDIZINTECHNIK BMT1100

INVERSION RECOVERY



BRUKER MEDIZINTECHNIK BMT1100

SPIN ECHO IMAGING

from these two physical methods, other imaging techniques

have come up, such as positron emission tomography, which are

also finding increasing application, particularly in the

early detection of cancer. We will not spend more time on

this subject except to say that this field has been one, in

which, not only were new mathematical approaches very effective,

but the computational aspects were the ones that decided the

practical applicability of the technique. In fact, a complete

scientific journal entitled "The Journal of Computed Tomography"

is being published, having come up to the tenth volume in 1986,

which shows the outstandingly significant role of computation

in modern medical research and practice. I believe that the

convolution method has still greater potentialities by the

inclusion of what is known as parallel processing in computation,

which is being rapidly developed in recent years.


## 7. Computerization of Logic

We shall now discuss the last topic of my talk, namely

computers for logic. This is a topic in which I have been

working for the last ten years and much of it is still in the

process of development. However, I thought I should include

but which at the same time is directly interpretable on the

computer _via_ Boolean algebra.  I hope I will have time to

indicate at least in a broad way its nature and its advantage.

Essentially, all computers have a central processing unit

(CPU) in which the fundamental arithmetical operations of

addition and multiplication are carried out, but not all

instructions that have to be given to computers, for automatic

control of machines and so on, are arithmetical, or in the form

in which they can be converted into ordinary algebra.  For

instance, it will be necessary to compare one thing with

another, and if the two are equal, it will / out one operation
                                          carry

while if the two are different, another operation is to be

performed.  This is the simplest of what may be called a

"logical operation".  So also, we may be given a list of

related entities /such as  a is related to b, c and d,  b is

related to c, e and f,  c is related to  a, d and f, and so on.

Then we have to answer the question — which are the entities

that are related by the given relation to b, c and d ?  In

standard computer practice, these can be worked out, but the

procedures depend upon what are  termed "goal-seeking searches",

## BOOLEAN ALGEBRA

| a ⊕ b = c | | |
|---|---|---|
| a\b | 1 | 0 |
| 1 | 1 | 1 |
| 0 | 1 | 0 |

| a ⊗ b = c | | |
|---|---|---|
| a\b | 1 | 0 |
| 1 | 1 | 0 |
| 0 | 0 | 0 |

| $a^c$ = b | |
|---|---|
| a | b |
| 1 | 0 |
| 0 | 1 |

Boolean algebra
is closed with
respect to sum,
product and
complement.

## LOGICAL TRUTH VALUES

| a OR b = c | | |
|---|---|---|
| a\b | T | F |
| T | T | T |
| F | T | F |

| a AND b = c | | |
|---|---|---|
| a\b | T | F |
| T | T | F |
| F | F | F |

| NOT a = b | |
|---|---|
| a | b |
| T | F |
| F | T |

There are only
two truth values
T, F, in
classical logic.

## SWITCHING CIRCUITS



OR GATE        AND GATE        Inverter(NOT)

| a\b | ON | OFF |
|---|---|---|
| ON | ON | ON |
| OFF | ON | OFF |

| a\b | ON | OFF |
|---|---|---|
| ON | ON | OFF |
| OFF | OFF | OFF |

| a | b |
|---|---|
| ON | OFF |
| OFF | ON |

The principles
of these switching
circuits are the
basis of all
computers.

The three sets of tables are isomorphic to one another.

gives rise to the operation of subtraction by reversing the

equation for it, as in (1), and similarly, multiplication leads

to division by reversal of Eq.(2) as given in Fig. 24, both

of which are included in standard arithmetic. It is supposed

that these "reverse" operations have no analogs in Boolean

algebra which has only the properties of a "Boolean ring"

in the parlance of modern algebra. However, this is not the

case. On can, in fact, reverse the operations of Boolean sum

---

Fig. 24: BA-2 algebra is essentially needed for the
representation of reverse relations in classical logic.

---

and Boolean product, by using tne corresponding analogs of

the above two equations (1) and (2) for subtraction and division,

and thus obtain interesting consequences. These are shown

in Fig. 24, and it will be seen that both types of reversal

are possible for Boolean arithmetic, but that they lead to two

more new "truth values", D and X, in addition to T and F of

classical logic.

Thus, BA-1 algebra is insufficient for representing fully

all consequences in classical logic, in that T and F are not

enough to represent the extended algebra of the logical

truth tables, and the reversal of the operations OR and AND

leads to four different "truth values". Without going into

details, I may say that it is possible to build a matrix algebra,

in which the four states T, F, D, X are represented by 2-element

Boolean vectors as shown in Eq.(3), which we have named BA-2.

The use of "or" and "and" between T and F for the states D and X

is justified by the descriptions given in Eq.(3). There is no

time to give details.

This is not the end of the story, but the beginning of

a whole new approach to logic. Perhaps the most significant

feature of this approach is the recognition and symbolization

of such reverse operators, which lead to what we have called

"unary" relations. Thus, the truth tables for OR and AND
as in Eqs.(1a) and (2a) of Fig.25,
can be taken over as 2x2 matrices,/and the answer to a query

such as, "If a AND b is given to be true, and a is given as

T (or F), what is b in each case?" can be written in the

language of Boolean vector-matrix formalism (BVMF) as in Eqs.(1b,c)
these
and (2b,c). As will be seen from / , all the <u>essential</u> conclusions

---

Fig. 25: Unary relations in BA-2 algebra and in the general
theory of relations.

made in Fig. 24 about the consequences of the reversed

operators are reproduced in the examples (1b,c) and (2b,c)

by this formalism.

There is no time to talk about various applications

of this vector-matrix formalism, using 2-vectors and 3-vectors,

for the branches of logic known as propositional calculus and

quantifier calculus. I shall, however, indicate the great

convenience of the vector-matrix formalism in what is known as

predicate calculus, which is most relevant to practical applications

of logic. The basic formula that connects two sets of objects

can be represented in this formalism by a matrix $R_{ij}$ as

shown in Eq.(4). Thus, suppose $A_1$, $A_2$ ... $A_m$ are a number

of persons and $B_1$, $B_2$ ... $B_n$ are the possible cities from

which they come. Then, if $A_1$ belongs to the city $B_2$, we say

that the element $R_{12}$ of the relation "coming from city" is 1,

and 0 otherwise. Thus, we obtain an array of $R_{ij}$ (i = 1 to m,

j = 1 to n), each of which is a Boolean number, having only the

values 1, 0. Then, if m = 5, and we have a subset of A's

consisting of $A_1$, $A_3$, $A_5$, then we can represent them by a

5-element Boolean vector $(a_1, a_2, a_3, a_4, a_5) = (1\ 0\ 1\ 0\ 1)$.

$$s \xrightarrow[\underset{\sim}{P}]{\text{Professor of}} \underset{\sim}{p} \qquad\qquad \underset{\sim}{p} \xrightarrow[\underset{\sim}{S}]{\text{Student of}} \underset{\sim}{s}$$

$$
\text{GM1} = \;
\begin{array}{c|cccc}
 & p_1 & p_2 & p_3 & p_4 \\
\hline
s_1 & 1 & 0 & 1 & 0 \\
s_2 & 0 & 1 & 1 & 0 \\
s_3 & 1 & 0 & 0 & 1 \\
s_4 & 0 & 0 & 0 & 1 \\
s_5 & 1 & 1 & 1 & 0 \\
\end{array}
$$

$$
\begin{aligned}
\text{GM3} &= \\
&= \text{GT(GM1)}
\end{aligned}
\quad
\begin{pmatrix}
1 & 0 & 1 & 0 & 1 \\
0 & 1 & 0 & 0 & 1 \\
1 & 1 & 0 & 0 & 1 \\
0 & 0 & 1 & 1 & 0 \\
\end{pmatrix}
$$

$$s \xrightarrow[\underset{\sim}{P^c}]{\text{Not professor of}} \underset{\sim}{p} \qquad\qquad \underset{\sim}{p} \xrightarrow[\underset{\sim}{S^c}]{\text{Not student of}} \underset{\sim}{s}$$

$$
\begin{aligned}
\text{GM2} &= \\
&= \text{GC(GM1)}
\end{aligned}
\quad
\begin{pmatrix}
0 & 1 & 0 & 1 \\
1 & 0 & 0 & 1 \\
0 & 1 & 1 & 0 \\
1 & 1 & 1 & 0 \\
0 & 0 & 0 & 1 \\
\end{pmatrix}
$$

$$
\begin{aligned}
\text{GM4} &= \\
&= \text{GTC(GM1)}
\end{aligned}
\quad
\begin{pmatrix}
0 & 1 & 0 & 1 & 0 \\
1 & 0 & 1 & 1 & 0 \\
0 & 0 & 1 & 1 & 0 \\
1 & 1 & 0 & 0 & 1 \\
\end{pmatrix}
$$

---

1. Who are all the professors who teach the students $s_1$ and $s_2$ ?

| | |
|---|---|
| $s_1, s_2$ gives GVS = (1 1 0 0 0) | Vector representing given students |
| GUNPDT(GVS,GM1) = GVP | Relation "student to professor" GM1 applied (1a) |
| GVP = (1 1 1 0) —Ans: $p_1, p_2, p_3$ | Vector representing required professors. (1b) |

2. Which professors among $p_2$ and $p_4$ teach at least one of the students $s_1$, $s_3$ ?

Input GVP1 = (0 1 0 1), GVS1 = (1 0 1 0 0)

GUNPDT(GVS1,GM1) = GVP, GVP = (1 0 1 1)  :  $p_1$, $p_3$, $p_4$ teach either $s_1$ or $s_2$ (2a)

GVIDYA(GVP,GVP1) = GVP2 , GVP2 = (0 0 0 1) : Of these, only $p_4$ is contained in the input set $p_2$, $p_4$ (2b)

26

we give some examples in Fig. 27.

The statements there are fairly self-explanatory in view

of what has been said earlier. In these problems also, we

---

Fig. 27: Examples of logical "tongue-twisters" solved using
BVMF and MATLOG

---

deal with a set of five students and four professors, but

consider two relations that can exist between them, namely

$P$ = "professor of", and $E$ = "examiner of", whose matrices are

given in (1), as GMP and GME. (Note that this GMP is different

from that in Fig. 26.) Now we ask two questions as given in

(a) and (b) of Fig. 27. It requires some effort to differentiate

between the two in common language, but in BVMF, the two are

given by quite different formulae, as in (3) and (4) of Fig. 27.

Thereafter, we can leave it to the computer, and interestingly,

the output vectors GVPA and GVPB, for the questions in (a) and (b),

are quite different as given in (5). These mean that the answer

to the first problem is that there is only one professor $p_4$

who satisfies both the conditions of this problem, while the

answer to the second problem is that three professors $p_1$, $p_2$, $p_4$

satisfy the given conditions. I am giving this example to

indicate how even such tricky problems can be formulated in

almost a single line in BVMF, and can be computerized via MATLOG.

C7116 46

As already mentioned in the beginning, this does not

at all mean that computers can replace human thought.

The computer can only work out the answer to a query regarding

the validity of a conclusion in the forms : "yes, "no",

"does not follow from the assumption" or "contradictory to

known facts"; but it cannot suggest what questions should be

asked to clinch the issue, if the deduction is found to be

inconclusive.  Much can be said on this inter-relationship

between semantics and logic, but, in essence, we can say that

the latter, which is the grammar of reasoning, can be

computerized, while the former continues to be still in the

realm of human endeavour.

these
(see Fig.8). Now-a-days, / will take a fraction of a second

on the best computers. I shall not mention much about these

theoretical ideas except to indicate that the two principal

considerations that were put forward by us at that time have

played a dominant part in all techniques involving molecular

packing, and the fitting of molecular chains to electron density

diagrams obtained from crystal structure determination, ever since.

These are the use of dihedral angles (rotation angles), such as

$(\phi, \psi)$ of Fig. 8(a), and the idea that non-bonded atoms cannot come

limiting
closer than certain specified/distances, which was used for

preparing the diagram in Fig. 8(b). These have been so well

---

Fig.8 (a) Rotation angles $\phi$ and $\psi$ specifying the configuration
        of a pair of peptide units.

     (b) The $(\phi, \psi)$-plot of the observed configurations
        in myoglobin (marked as small circles.)

---

computerized that, now-a-days one can just feed in electron

density and the chemical nature of the molecule and ask the

computer to shift, turn and twist the molecule until the best

fit with the electron density is obtained, subject to the

restrictions imposed by the limiting distances mentioned above.

More about this later.

**4. Crystallography and Biology with special reference to drug design**

The precise techniques / energy minimization, and that have been developed for for the

working out of the best possible configurations of molecules,

have resulted in a        body of knowledge which has vastly expanded

of the
our knowledge / physical chemistry, molecular biophysics and

of the biological systems.
biochemistry, / I shall illustrate these by a few examples.

We had seen that the heme molecule in hemoglobin is located

in between the different protein chains and the knowledge

obtained by crystallography about the particular amino acid

residues near the location of the heme had led to a clear

understanding of the mechanism of the function and reactivity

the protein for
of / oxygen transport in blood. Simultaneously, it has / become
                                                        also

                                                        the
clear that many biological conditions associated with / malfunction

of hemoglobin can be attributed to changes in one or two

amino acid residues produced by mutation. Fig.9 shows the

location of the mutated residues for a large number of

---

Fig.9 Mutation sites in pathalogical hemoglobins.

---

pathalogical hemoglobins. It will be seen that, though the amino

.18.

in a linear peptide chain
sequence/may not show it,      the x-ray crystallographic

three-dimensional structure shows that they often occur

close to the region where the heme is held, or in the gaps

between neighbouring molecules. Apart from this general

feature, the knowledge obtained from crystallography has

profoundly affected biochemical investigations for the

of the blood,
understanding of such  diseases/and has led even to the design

of suitable chemical drugs for alleviating the symptoms.

Turning to     more fundamental problems in biology,

such as the action of enzymes, the chemists had formulated

the theory that enzymes, which regulate all biological reactions,

did so by forming a close molecular fit with the substrate

molecule concerned. This was a conjuncture in the past,

although a very reasonable one and much substantiated by all

available knowledge. However, a confirming evidence in this

matter has come only from crystallographic studies. One of

the earliest proteins to be solved, namely lysozyme which is

shown in Fig. 10, displays this feature remarkably well.

Fig. 10 : Space-filling molecular models of free lysozyme (left)
and when bound to its substrate (right).

The first picture is that of the protein which is seen to have

a cleft on its left, and the second picture is of the enzyme-

substrate complex. It will be seen that the substrate and

enzyme form a perfect fit and the theoretical postulation of

molecular fit has been demonstrated in all detail by three-

dimensional x-ray crystallography.

I will therefore restrict myself to the subject of

molecular fit, and give a few examples where crystallography,

in combination with conformational theorey, has been applied

to obtain information regarding this. Fig.11 is a study made

by Prof. Viswamitra of the Indian Institute of Science

---

Fig.11 : Space-filling models of (a) the antibiotic TANDEM
and (b) its complex with DNA.

---

in which he determined the molecular structure of an antibiotic

TANDEM, shown in the top half of Fig. 11, which acts by binding

with DNA. He could then show that this fits very well to the

DNA double helix, as shown in the lower half. This idea has

been employed in a number of studies for finding the basis of

drug-molecule interactions. Such crystallographic studies

have even been extended to virus molecules and the interaction

of antivirul agents with the virus. Fig. 12 is an example of a

very recent study made by Prof. Rossman in Purdue University, USA.

---

Fig. 12 : Crystallographically determined mode of binding of
          antiviral agent WIN to virus protein VP1.

---

Not only does it show that the viral protein has a cavity which

can be fitted by the antiviral agent WIN, so that the normal

biochemical reactions involved in the reproduction of the virus

molecule are inhibited, but it is also an example of the extent

to which crystallography has progressed. The virus itself

required 200,000 reflections for its structure determination.

After that, the mode of attachment of the drug was determined

by using what is known as a difference-Fourier synthesis,

in which the difference in intensities between x-ray diffraction

patterns, of the pure virus and the virus-drug complex, was

used to find out the location of the drug molecule within the

virus. Fig. 13 shows the electron density distribution of the

drug molecule as reconstructed by a computer. It will be seen

---

Fig. 13. Front page of Science magazine containing an
         illustration of the molecule of an antiviral agent
         fitted to the computer-displayed electron density cage.

---

that it effectively indicates a cage within which the molecule

must be situated, and the exact location of the molecule and

its conformation can also be varied and fitted in, by employing

the computer for on-line operation. I had the good fortune

to advise some of the workers in this field of computer-fitting

of molecules in USA in 1977-78, and this approach is

universally employed both in crystallography and conformational

theory now-a-days, and "computer graphics terminals" are even

commercially available for this purpose.

## 5. Conformational theory for molecular fit

Coming back to theoretical calculations, Fig. 14 is an

example where the crystal structure determination of an

---

Fig. 14: Daunomycin fitted to DNA by energy minimization on
          a computer

---

antibiotic (daunomycin) was the starting point of an attempt

to fit the molecule to the DNA double helix. The work was

done in Prof. Pullman's laboratory in Paris, and as the figure

will show, the fitting can be done readily by opening out the

double helix in the region where the drug makes its entry.

I wish to emphasize that this has not been done by constructing

a model, and then turning and twisting the molecular structure

so as to obtain the fit, but purely _via_ the use of a computer.

                                an
On the other hand, the fitting of/inhibitor, $\beta$-P-P-P, to the

active site of the enzyme thermolysin, which was carried out

by Prof. Rao in Bangalore, was done, not automatically on a

computer, but using the computer without on-line interaction.

The investigator starts with a structure close to a reasonable

fit, and then modifies it step by step, by feeding in and

taking out the data from the computer in the form of drawings.

This study showed that the inhibitor enters the enzyme in one

---

Fig.15: Mode of attachment of the inhibitor $\beta$-P-P-P which
imitates the substrate of the enzyme thermolysin.

---

configuration and then both its conformation, and that of the

enzyme in the local region, are modified so as to obtain the

best fit. The reason why I am showing this picture is to indicate

that some of the latest ideas of this type have been computerized

right in our country and interesting consequences have been

deduced. However, even with a computer at one's disposal, this

is a long process requiring weeks or months. If only an interactive

graphic display system, as has been used in all the advanced

countries, were made available, the time could be reduced to

probably a week or two. In fact, as already mentioned, such

systems are available commercially in USA and Japan, and

I believe that it is time that such up-to-date equipment be

made available to the capable scientists of our country.


Emergence of the new field QSAR

Such studies on molecular fit have opened many new

avenues in the field of rational drug design, because drugs

invariably either promote, or inhibit, particular biological

reactions, either in the organism producing infection, or in the

host for producing resistance. In fact, a new field entitled

Quantitative Structure Activity Relationship (QSAR) has emerged

during the last few years, and strategies are being developed

for actively pursuing this for drug design. Fig. 16 shows the
      page
title/of a very interesting paper, on protein crystallography

and computer graphics, published in the German journal Angewandte

---

Fig. 16: Title page and flow chart about QSAR reproduced from
         a recent article.

---

Chemie (Applied Chemistry), which emphasizes this aspect of

biomolecular physics. As will be seen from one of the flow

charts in this report, computer graphics and theoretical

calculations play as important/part as the determination of
a

three-dimensional structure by x-ray crystallography. Even

regular symposia are held on this subject of QSAR, as will be

seen from Fig. 16.


## 6. Computerized Tomography

Turning to the subject of tomography, my interests in

this field dates from 1970 when some very interesting work was

done along with Dr. Lakshminarayanan, at the University of Chicago

on the application of Fourier transforms and the novel technique

the title page of
of convolutions for this purpose. Fig. 17 shows/our first paper

on the subject entitled "Reconstruction of substance from shadow"

---

Fig. 17:   Fourier transform equations formulated in polar
coordinates for axial tomography.

---

As you will see from this, the Fourier transform Eq.(1), which

is basic to this subject, is the same as that used in x-ray

crystallography. However, although crystallographers had applied

this technique for finding out the three-dimensional structure

of virus particles from two-dimensional electron microscopic

images in different directions, we suggested that this

reconstruction is much faster by employing polar coordinates

for the Fourier transform as in Eqs. (2a,b), using the

principle of axial tomography. In these, $g(\ell ; \theta)$ is the

linear shadowgraphs at $\theta$, and $f(r, \phi)$ is the two-dimensional

image that is reconstructed for a two-dimensional section at

right angles to the axis of rotation.

This was followed by another paper published in the

Proceedings of the National Academy of Sciences, USA, whose

title page is shown in Fig. 18. In this, the way in which a

series of two-dimensional images can be converted into a

three-dimensional reconstruction by axial tomography is shown

in the diagram. In addition, a new modification of the Fourier

method, which we called as the "convolution method" was presented
this paper, and its
in / basic mathematics is shown in Eqs. (3), (4) and (5).

---

Fig. 18:   Principle of the convolution method of reconstruction
           in axial tomography.

---

In this, the reconstructed two-dimensional image $f(r , \phi)$ is
                         density                    at $\theta$,
a simple integral of the/functions $g'(\ell ; \theta)$, which represent

the
/modified shadowgraphs at $\theta$ , produced as a result

of the convolution process shown in Eqs (4a,b). The detailed

way in which the modified function is calculated is by the

procedure given in Eqs (5a,b). The mathematics is not quite

relevant to this lecture, but what is important is that,

as a consequence, the reconstruction can be performed almost

a hundred times faster than the Fourier method and, as will be

seen from Fig. 19, it is also about ten times more accurate.

Fig. 19: Tables 1 and 3 reproduced from the PNAS paper
showing the distinct advantage in speed and
accuracy of the convolution method.

Tables 1 and 3 of Fig. 19 are the comparative data of the

accuracy and computing time required for the different methods

of reconstruction , namely Cartesian Fourier transform (FTC),

Polar Fourier transform (FTP) and Convolution (CON). The

most interesting feature is that the larger the amount of

data that is available from measurements for reconstruction,

the faster is the convolution method relative to the Fourier

method. This technique is applied widely in CAT-scan

instruments particularly for medical radiography.

**TABLE 1.** *Values of the mean relative errors ($\mathcal{R}$) for the circular disk for $r < 0.8$ using the CON, FTP, and FTC methods*
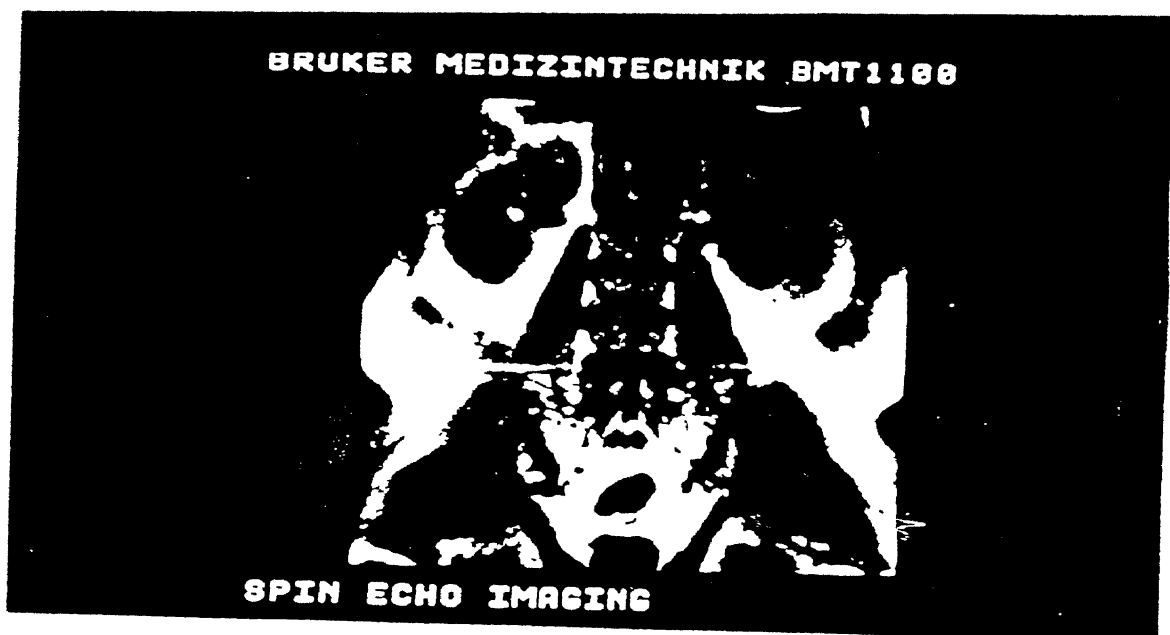
| $\theta_0(°)$ | $a$ | $\mathcal{R}$(in %) for CON | FTP | FTC |
|---|---|---|---|---|
| 30 | 0.2 | 1.5 | 9.4 | 7.7 |
| 30 | 0.1 | 0.6 | 7.6 | 5.3 |
| 15 | 0.2 | 1.2 | 8.0 | 7.6 |
| 15 | 0.1 | 0.3 | 5.3 | 5.2 |
| 15 | 0.1 | 0.3 | 2.9* | 3.9* |

* $R_0 = 0.125$; for other values, $R_0 = 0.25$.

**TABLE 3.** *Computing times ($t$) required using the different methods for a typical object used in the study*

| $\theta_0(°)$ | $a$ | $t$ in seconds CON | FTP | FTC |
|---|---|---|---|---|
| 30 | 0.2 | 2.0 | 30.2 | 130 |
| 30 | 0.1 | 2.6 | 48.6 | 480 |
| 15 | 0.2 | 3.6 | 41.5 | 130 |
| 15 | 0.1 | 4.8 | 83.0 | 480 |
| 15 | 0.1 | 4.8 | 158.5* | 487* |

* $R_0 = 0.125$ for these values; $R_0 = 0.25$ for others.

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SPIN ECHO IMAGING

(21)

from these two physical methods, other imaging techniques

have come up, such as positron emission tomography, which are

also finding increasing application, particularly in the

early detection of cancer. We will not spend more time on

this subject except to say that this field has been one, in

which,not only were new mathematical approaches very effective,

but the computational aspects were the ones that decided the

practical applicability of the technique. In fact, a complete

scientific journal entitled "The Journal of Computed Tomography"

is being published, having come up to the tenth volume in 1986,

which shows the outstandingly significant role of computation

in modern medical research and practice. I believe that the

convolution method has still greater potentialities by the

inclusion of what is known as parallel processing in computation,

which is being rapidly developed in recent years.


## 7. Computerization of Logic

We shall now discuss the last topic of my talk, namely

computers for logic. This is a topic in which I have been

working for the last ten years and much of it is still in the

process of development. However, I thought I should include

but which at the same time is directly interpretable on the

computer _via_ Boolean algebra.  I hope I will have time to

indicate at least in a broad way its nature and its advantage.

Essentially, all computers have a central processing unit

(CPU) in which the fundamental arithmetical operations of

addition and multiplication are carried out, but not all

instructions that have to be given to computers, for automatic

control of machines and so on, are arithmetical, or in the form

in which they can be converted into ordinary algebra.  For

instance, it will be necessary to compare one thing with

                                           carry
another, and if the two are equal, it will / out one operation

while if the two are different, another operation is to be

performed.  This is the simplest of what may be called a

"logical operation".  So also, we may be given a list of

related entities /such as  a is related to b, c and d,  b is

related to c, e and f,  c is related to  a, d and f, and so on.

Then we have to answer the question — which are the entities

that are related by the given relation to b, c and d ?  In

standard computer practice, these can be worked out, but the

procedures depend upon what are  termed "goal-seeking searches",

## BOOLEAN ALGEBRA

$a \oplus b = c$

| a\b | 1 | 0 |
|-----|---|---|
| 1   | 1 | 1 |
| 0   | 1 | 0 |

$a \otimes b = c$

| a\b | 1 | 0 |
|-----|---|---|
| 1   | 1 | 0 |
| 0   | 0 | 0 |

$a^c = b$

| a | b |
|---|---|
| 1 | 0 |
| 0 | 1 |

Boolean algebra
is closed with
respect to sum,
product and
complement.
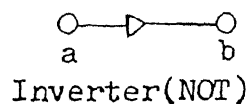
## LOGICAL TRUTH VALUES

a OR b = c

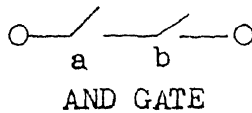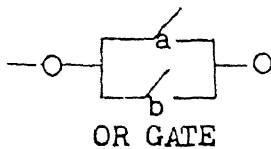| a\b | T | F |
|-----|---|---|
| T   | T | T |
| F   | T | F |

a AND b = c

| a\b | T | F |
|-----|---|---|
| T   | T | F |
| F   | F | F |

NOT a = b

| a | b |
|---|---|
| T | F |
| F | T |

There are only
two truth values
T, F, in
classical logic.

## SWITCHING CIRCUITS

OR GATE

AND GATE

Inverter(NOT)

| a\b | ON | OFF |
|-----|----|-----|
| ON  | ON | ON  |
| OFF | ON | OFF |

| a\b | ON  | OFF |
|-----|-----|-----|
| ON  | ON  | OFF |
| OFF | OFF | OFF |

| a   | b   |
|-----|-----|
| ON  | OFF |
| OFF | ON  |

The principles
of these switching
circuits are the
basis of all
computers.

The three sets of tables are isomorphic to one another.

gives rise to the operation of subtraction by reversing the

equation for it, as in (1), and similarly, multiplication leads

to division by reversal of Eq.(2) as given in Fig. 24, both

of which are included in standard arithmetic. It is supposed

that these "reverse" operations have no analogs   in Boolean

algebra which has only the properties of a "Boolean ring"

in the parlance of modern algebra. However, this is not the

case. On can, in fact, reverse the operations of Boolean sum

---

Fig. 24: BA-2 algebra is essentially needed for the
representation of reverse relations in classical logic.

---

and Boolean product, by using the corresponding analogs of

the above two equations (1) and (2) for subtraction and division,

and thus obtain interesting consequences. These are shown

in Fig. 24, and it will be seen that both types of reversal

are possible for Boolean arithmetic, but that they lead to two

more new "truth values", D and X, in addition to T and F of

classical logic.

Thus, BA-1 algebra is insufficient for representing fully

all consequences in classical logic, in that T and F are not

enough to represent the extended algebra of the logical

truth tables, and the reversal of the operations OR and AND

leads to four different "truth values". Without going into

details, I may say that it is possible to build a matrix algebra,

in which the four states T, F, D, X are represented by 2-element

Boolean vectors as shown in Eq.(3), which we have named BA-2.

The use of "or" and "and" between T and F for the states D and X

is justified by the descriptions given in Eq.(3). There is no

time to give details.

This is not the end of the story, but the beginning of

a whole new approach to logic. Perhaps the most significant

feature of this approach is the recognition and symbolization

of such reverse operators, which lead to what we have called

"unary" relations. Thus, the truth tables for OR and AND

as in Eqs.(1a) and (2a) of Fig.25,
can be taken over as 2x2 matrices,/and the answer to a query

such as, "If a AND b is given to be true, and a is given as

T (or F), what is b in each case?" can be written in the

language of Boolean vector-matrix formalism (BVMF) as in Eqs.(1b,c)

these
and (2b,c). As will be seen from / , all the essential conclusions

Fig. 25: Unary relations in BA-2 algebra and in the general
theory of relations.

made in Fig. 24 about the consequences of the reversed

operators are reproduced in the examples (1b,c) and (2b,c)

by this formalism.

There is no time to talk about various applications

of this vector-matrix formalism, using 2-vectors and 3-vectors,

for the branches of logic known as propositional calculus and

quantifier calculus. I shall, however, indicate the great

convenience of the vector-matrix formalism in what is known as

predicate calculus, which is most relevant to practical applications

of logic. The basic formula that connects two sets of objects

can be represented in this formalism by a matrix $R_{ij}$ as

shown in Eq.(4). Thus, suppose $A_1$, $A_2$ ... $A_m$ are a number

of persons and $B_1$, $B_2$ ... $B_n$ are the possible cities from

which they come. Then, if $A_1$ belongs to the city $B_2$, we say

that the element $R_{12}$ of the relation "coming from city"is 1,

and 0 otherwise. Thus, we obtain an array of $R_{ij}$ (i = 1 to m,

j = 1 to n), each of which is a Boolean number, having only the

values 1, 0. Then, if m = 5, and we have a subset of A's

consisting of $A_1$, $A_3$, $A_5$, then we can represent them by a

5-element Boolean vector $(a_1, a_2, a_3, a_4, a_5)$ = (1 0 1 0 1).

$$s \xrightarrow[\underset{\sim}{p}]{\text{Professor of}} p$$

$$p \xrightarrow[\underset{\sim}{S}]{\text{Student of}} s$$

$$
GM1 = \begin{array}{c|cccc}
 & p_1 & p_2 & p_3 & p_4 \\
\hline
s_1 & 1 & 0 & 1 & 0 \\
s_2 & 0 & 1 & 1 & 0 \\
s_3 & 1 & 0 & 0 & 1 \\
s_4 & 0 & 0 & 0 & 1 \\
s_5 & 1 & 1 & 1 & 0 \\
\end{array}
$$

$$
GM3 = GT(GM1) = \begin{pmatrix}
1 & 0 & 1 & 0 & 1 \\
0 & 1 & 0 & 0 & 1 \\
1 & 1 & 0 & 0 & 1 \\
0 & 0 & 1 & 1 & 0 \\
\end{pmatrix}
$$

$$s \xrightarrow[\underset{\sim}{p^c}]{\text{Not professor of}} p$$

$$p \xrightarrow[\underset{\sim}{S^c}]{\text{Not student of}} s$$

$$
GM2 = GC(GM1) = \begin{pmatrix}
0 & 1 & 0 & 1 \\
1 & 0 & 0 & 1 \\
0 & 1 & 1 & 0 \\
1 & 1 & 1 & 0 \\
0 & 0 & 0 & 1 \\
\end{pmatrix}
$$

$$
GM4 = GTC(GM1) = \begin{pmatrix}
0 & 1 & 0 & 1 & 0 \\
1 & 0 & 1 & 1 & 0 \\
0 & 0 & 1 & 1 & 0 \\
1 & 1 & 0 & 0 & 1 \\
\end{pmatrix}
$$

1. Who are all the professors who teach the students $s_1$ and $s_2$ ?

$s_1, s_2$ gives GVS = (1 1 0 0 0)   Vector representing given students

GUNPDT(GVS,GM1) = GVP   Relation "student to professor" GM1 applied          (1a)

GVP = (1 1 1 0) —Ans: $p_1, p_2, p_3$   Vector representing required professors.          (1b)

2. Which professors among $p_2$ and $p_4$ teach at least one of the students $s_1$, $s_3$ ?

Input GVP1 = (0 1 0 1),  GVS1 = (1 0 1 0 0)

GUNPDT(GVS1,GM1) = GVP,  GVP = (1 0 1 1)  :  $p_1$, $p_3$, $p_4$ teach either $s_1$ or $s_2$          (2a)

GVIDYA(GVP,GVP1) = GVP2 , GVP2 = (0 0 0 1) : Of these, only $p_4$ is contained in the input set $p_2$, $p_4$          (2b)

$\textcircled{2 6}$

we give some examples in Fig. 27.

The statements there are fairly self-explanatory in view

of what has been said earlier. In these problems also, we

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Fig. 27: Examples of logical "tongue-twisters" solved using
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deal with a set of five students and four professors, but

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$P$ = "professor of", and $E$ = "examiner of", whose matrices are

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between semantics and logic, but, in essence, we can say that

the latter, which is the grammar of reasoning, can be

computerized, while the former continues to be still in the

realm of human endeavour.

**TABLE 1.** *Values of the mean relative errors (ℜ) for the circular disk for r < 0.8 using the CON, FTP, and FTC methods*
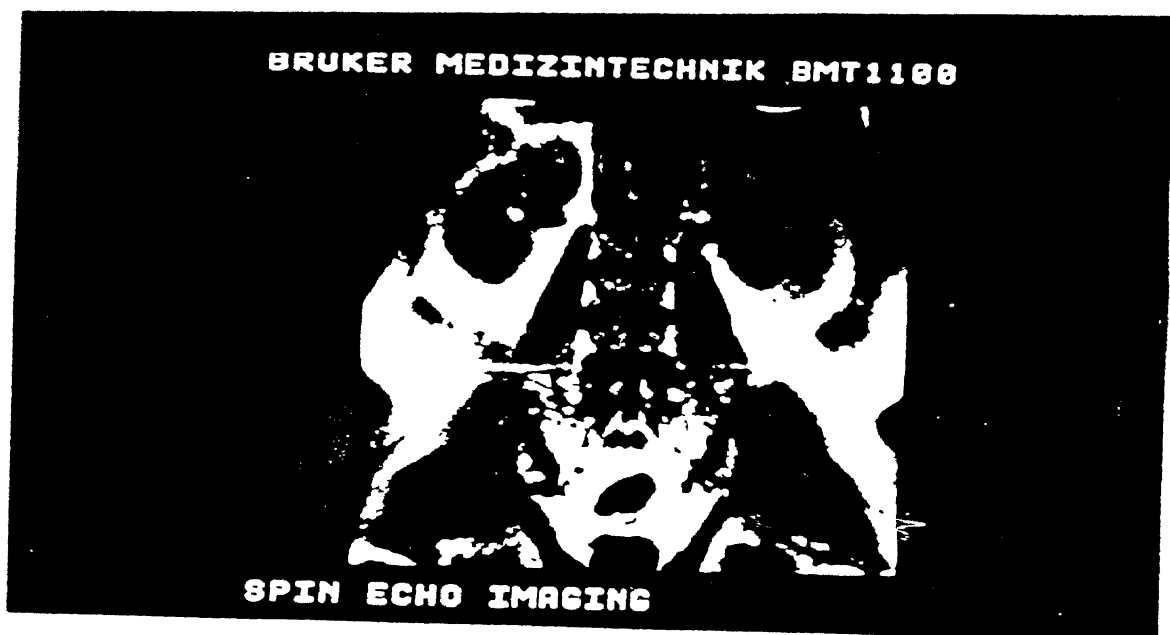
| $\theta_i(°)$ | $a$ | ℜ(in %) for CON | ℜ(in %) for FTP | ℜ(in %) for FTC |
|---|---|---|---|---|
| 30 | 0.2 | 1.5 | 9.4 | 7.7 |
| 30 | 0.1 | 0.6 | 7.6 | 5.3 |
| 15 | 0.2 | 1.2 | 8.0 | 7.6 |
| 15 | 0.1 | 0.3 | 5.3 | 5.2 |
| 15 | 0.1 | 0.3 | 2.9* | 3.9* |

\* $R_i$ = 0.125; for other values, $R_i$ = 0.25.

**TABLE 3.** *Computing times (t) required using the different methods for a typical object used in the study*

| $\theta_i(°)$ | $a$ | t in seconds CON | t in seconds FTP | t in seconds FTC |
|---|---|---|---|---|
| 30 | 0.2 | 2.0 | 30.2 | 130 |
| 30 | 0.1 | 2.6 | 48.6 | 480 |
| 15 | 0.2 | 3.6 | 41.5 | 130 |
| 15 | 0.1 | 4.8 | 83.0 | 480 |
| 15 | 0.1 | 4.8 | 158.5* | 487* |

\* $R_i$ = 0.125 for these values; $R_i$ = 0.25 for others.

⑲

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INVERSION RECOVERY



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SPIN ECHO IMAGING

(21)

from these two physical methods, other imaging techniques

have come up, such as positron emission tomography, which are

also finding increasing application, particularly in the

early detection of cancer. We will not spend more time on

this subject except to say that this field has been one, in

which, not only were new mathematical approaches very effective,

but the computational aspects were the ones that decided the

practical applicability of the technique. In fact, a complete

scientific journal entitled "The Journal of Computed Tomography"

is being published, having come up to the tenth volume in 1986,

which shows the outstandingly significant role of computation

in modern medical research and practice. I believe that the

convolution method has still greater potentialities by the

inclusion of what is known as parallel processing in computation,

which is being rapidly developed in recent years.

## 7. Computerization of Logic

We shall now discuss the last topic of my talk, namely

computers for logic. This is a topic in which I have been

working for the last ten years and much of it is still in the

process of development. However, I thought I should include

but which at the same time is directly interpretable on the

computer _via_ Boolean algebra.  I hope I will have time to

indicate at least in a broad way its nature and its advantage.

Essentially, all computers have a central processing unit

(CPU) in which the fundamental arithmetical operations of

addition and multiplication are carried out, but not all

instructions that have to be given to computers, for automatic

control of machines and so on, are arithmetical, or in the form

in which they can be converted into ordinary algebra.  For

instance, it will be necessary to compare one thing with

carry
another, and if the two are equal, it will / out one operation

while if the two are different, another operation is to be

performed.  This is the simplest of what may be called a

"logical operation".  So also, we may be given a list of

related entities /such as  a is related to b, c and d,  b is

related to c, e and f,  c is related to  a, d and f, and so on.

Then we have to answer the question — which are the entities

that are related by the given relation to b, c and d ?  In

standard computer practice, these can be worked out, but the

procedures depend upon what are  termed "goal-seeking searches",

## BOOLEAN ALGEBRA

| a ⊕ b = c | | |
|---|---|---|
| a\b | 1 | 0 |
| 1 | 1 | 1 |
| 0 | 1 | 0 |

| a ⊙ b = c | | |
|---|---|---|
| a\b | 1 | 0 |
| 1 | 1 | 0 |
| 0 | 0 | 0 |

$a^c = b$

| a | b |
|---|---|
| 1 | 0 |
| 0 | 1 |

Boolean algebra is closed with respect to sum, product and complement.

## LOGICAL TRUTH VALUES

| a OR b = c | | |
|---|---|---|
| a\b | T | F |
| T | T | T |
| F | T | F |

| a AND b = c | | |
|---|---|---|
| a\b | T | F |
| T | T | F |
| F | F | F |

NOT a = b

| a | b |
|---|---|
| T | F |
| F | T |

There are only two truth values T, F, in classical logic.

## SWITCHING CIRCUITS



OR GATE          AND GATE          Inverter(NOT)

| a\b | ON | OFF |
|---|---|---|
| ON | ON | ON |
| OFF | ON | OFF |

| a\b | ON | OFF |
|---|---|---|
| ON | ON | OFF |
| OFF | OFF | OFF |

| a | b |
|---|---|
| ON | OFF |
| OFF | ON |

The principles of these switching circuits are the basis of all computers.

The three sets of tables are isomorphic to one another.

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gives rise to the operation of subtraction by reversing the

equation for it, as in (1), and similarly, multiplication leads

to division by reversal of Eq.(2), as given in Fig. 24, both

of which are included in standard arithmetic. It is supposed

that these "reverse" operations have no analogs   in Boolean

algebra which has only the properties of a "Boolean ring"

in the parlance of modern algebra. However, this is not the

case. On can, in fact, reverse the operations of Boolean sum

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Fig. 24: BA-2 algebra is essentially needed for the
          representation of reverse relations in classical logic.

---

and Boolean product, by using the corresponding analogs of

the above two equations (1) and (2) for subtraction and division,

and thus obtain interesting consequences. These are shown

in Fig. 24, and it will be seen that both types of reversal

are possible for Boolean arithmetic, but that they lead to two

more new "truth values", D and X, in addition to T and F of

classical logic.

    Thus, BA-1 algebra is insufficient for representing fully

all consequences in classical logic, in that T and F are not

enough to represent the extended algebra of the logical

truth tables, and the reversal of the operations OR and AND

leads to four different "truth values". Without going into

details, I may say that it is possible to build a matrix algebra,

in which the four states T, F, D, X are represented by 2-element

Boolean vectors as shown in Eq.(3), which we have named BA-2.

The use of "or" and "and" between T and F for the states D and X

is justified by the descriptions given in Eq.(3). There is no

time to give details.

This is not the end of the story, but the beginning of

a whole new approach to logic. Perhaps the most significant

feature of this approach is the recognition and symbolization

of such reverse operators, which lead to what we have called

"unary" relations. Thus, the truth tables for OR and AND.

can be taken over as 2x2 matrices,/and the answer to a query

as in Eqs.(1a) and (2a) of Fig.25,

such as, "If a AND b is given to be true, and a is given as

T (or F), what is b in each case?" can be written in the

language of Boolean vector-matrix formalism (BVMF) as in Eqs.(1b,c)

and (2b,c). As will be seen from / , all the _essential_ conclusions

these

Fig. 25: Unary relations in BA-2 algebra and in the general
         theory of relations.

made in Fig. 24 about the consequences of the reversed

operators are reproduced in the examples (1b,c) and (2b,c)

by this formalism.

There is no time to talk about various applications

of this vector-matrix formalism, using 2-vectors and 3-vectors,

for the branches of logic known as propositional calculus and

quantifier calculus. I shall, however, indicate the great

convenience of the vector-matrix formalism in what is known as

predicate calculus, which is most relevant to practical applications

of logic. The basic formula that connects two sets of objects

can be represented in this formalism by a matrix $R_{ij}$ as

shown in Eq.(4). Thus, suppose $A_1$, $A_2$ ... $A_m$ are a number

of persons and $B_1$, $B_2$ ... $B_n$ are the possible cities from

which they come. Then, if $A_1$ belongs to the city $B_2$, we say

that the element $R_{12}$ of the relation "coming from city" is 1,

and 0 otherwise. Thus, we obtain an array of $R_{ij}$ (i = 1 to m,

j = 1 to n), each of which is a Boolean number, having only the

values 1, 0. Then, if m = 5, and we have a subset of A's

consisting of $A_1$, $A_3$, $A_5$, then we can represent them by a

5-element Boolean vector $(a_1, a_2, a_3, a_4, a_5) = (1\ 0\ 1\ 0\ 1)$.

$\xrightarrow[\underset{\sim}{P}]{\text{Professor of}}$ applied to $\underset{\sim}{s} \to \underset{\sim}{p}$

$\underset{\sim}{p} \xrightarrow[\underset{\sim}{S}]{\text{Student of}} \underset{\sim}{s}$

|  | $p_1$ | $p_2$ | $p_3$ | $p_4$ |
|---|---|---|---|---|
| $s_1$ | 1 | 0 | 1 | 0 |
| $s_2$ | 0 | 1 | 1 | 0 |
| GM1 = $s_3$ | 1 | 0 | 0 | 1 |
| $s_4$ | 0 | 0 | 0 | 1 |
| $s_5$ | 1 | 1 | 1 | 0 |

$$GM3 = GT(GM1) = \begin{pmatrix} 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 \end{pmatrix}$$

$\underset{\sim}{s} \xrightarrow[\underset{\sim}{P^c}]{\text{Not professor of}} \underset{\sim}{p}$

$$GM2 = GC(GM1) = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

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$$GM4 = GTC(GM1) = \begin{pmatrix} 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 & 1 \end{pmatrix}$$

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